CS738: Advanced Compiler Optimizations

Constant Propagation

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Agenda

- Using data flow analysis to identify “constant expressions” in a program
- Identify similarity/differences with bit-vector data flow analyses discussed earlier
- Other properties of constant propagation

Constant Propagation

- CP: Replace expressions that evaluate to same constant “c” every time they are executed, by the value “c”

DF Framework for CP

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Special Values for CP

- *NAC*: not a constant
  - If variable is inferred not to be a constant
  - Multiple (different valued) defs, non-const defs, assigned an "un-interpreted" value, ...
- *Undef*: No definition of the variable is seen yet - nothing known!

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NAC vs *Undef*

- *NAC* $\Rightarrow$ *too many* definitions seen for a variable v to declare v is **NOT** a constant
- *Undef* $\Rightarrow$ *too few* definitions seen to declare anything about the variable
- $\top$ is *Undef*; $\bot$ is *NAC*

CP Meet $\bigwedge$

- Recall the requirement

$$\top \bigwedge x = x$$

$$\bot \bigwedge x = \bot$$

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$c_1 \bigwedge c_2 = NAC$ when $c_1 \neq c_2$

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$Undef \bigwedge c = c$

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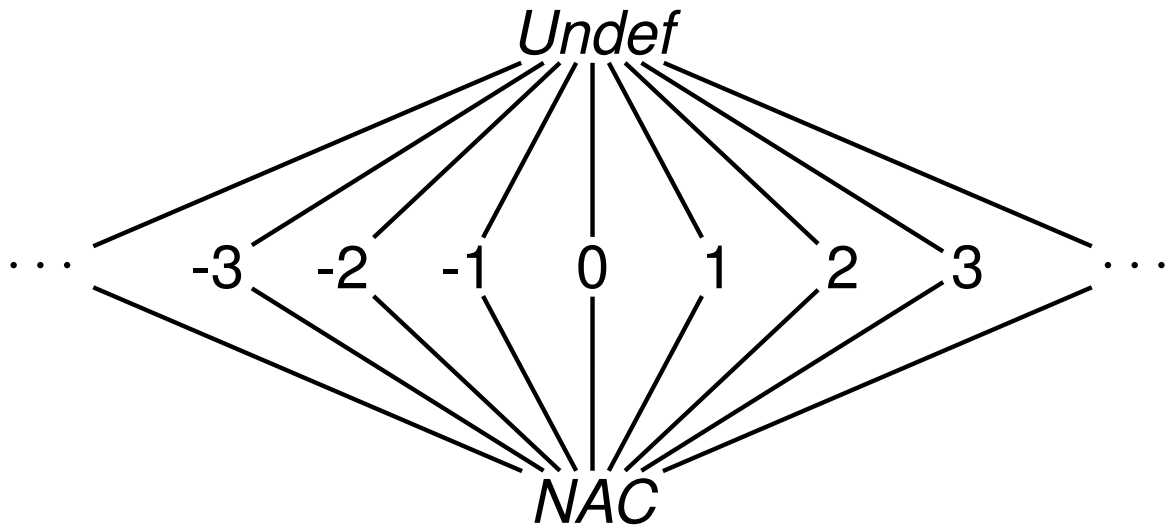
$c_1 \bigwedge c_2 = NAC$ when $c_1 \neq c_2$

$c \bigwedge c = c$

CP Semilattice for an integer variable



- Infinite domain, but finite height

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- CP Semilattice = Product of such lattices for all variables (of all types)
- Each semilattice has a **finite** height

Computing GEN

- Informal representation

Statement	GEN
$x = c$ // const	
$x = y + z$	
$x =$ *complicated expr*	

Computing GEN

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Computing GEN

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Statement	GEN
$x = c$ // const	$\{x \to c\}$
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$x =$ *complicated expr*	

Computing GEN

- Informal representation

Statement	GEN
$x = c$ // const	$\{x \to c\}$
$x = y + z$	if $\{y \to c_1, z \to c_2\}$ in IN then $\{x \to c_1 + c_2\}$ else if $\{y \to NAC\}$ in IN then $\{x \to NAC\}$
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$x = \mathit{complicated\ expr}$	$\{x \to \mathit{NAC}\}$

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- Case analysis on transfer function f
- $\mathit{NAC} \leq c \leq \mathit{Undef}$
- $x = c$ has constant transfer function.
- $x = \mathit{complicated\ expr}$ also has constant transfer function
- See the next slide for $x = y + z$ (and similar statements)

Monotonicity of CP: $x = y + z$

- Fix z to be one of *Undef*, c_2, *NAC*

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- Do this for all z choices.

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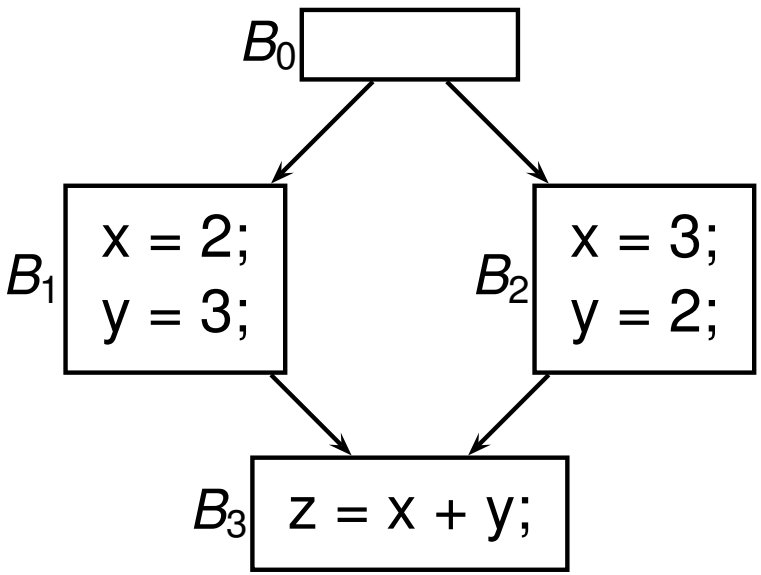
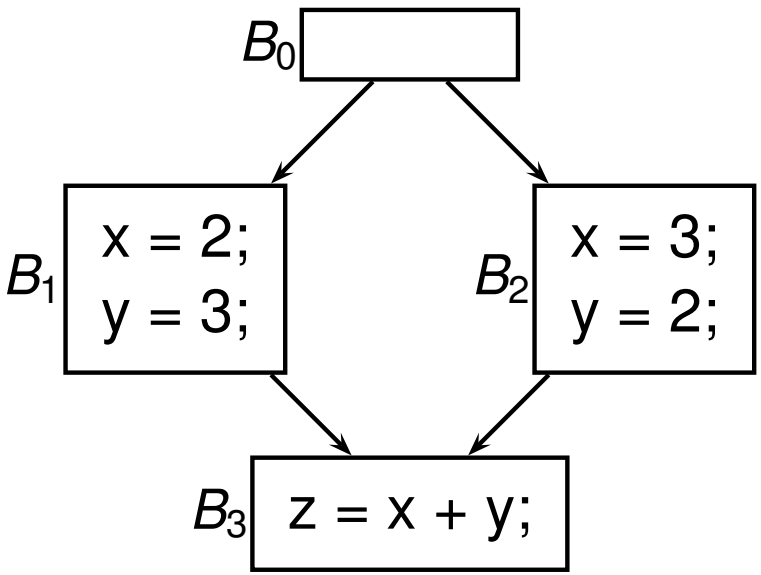
- Vary y over *Undef*, c_1, *NAC*
- Confirm that x does not "*increase*"
- Do this for all z choices.
- Similarly, fix y and vary z.

Nondistributivity of CP



- All paths:

Nondistributivity of CP



- All paths:
  - $B_0 \rightarrow B_1 \rightarrow B_3$

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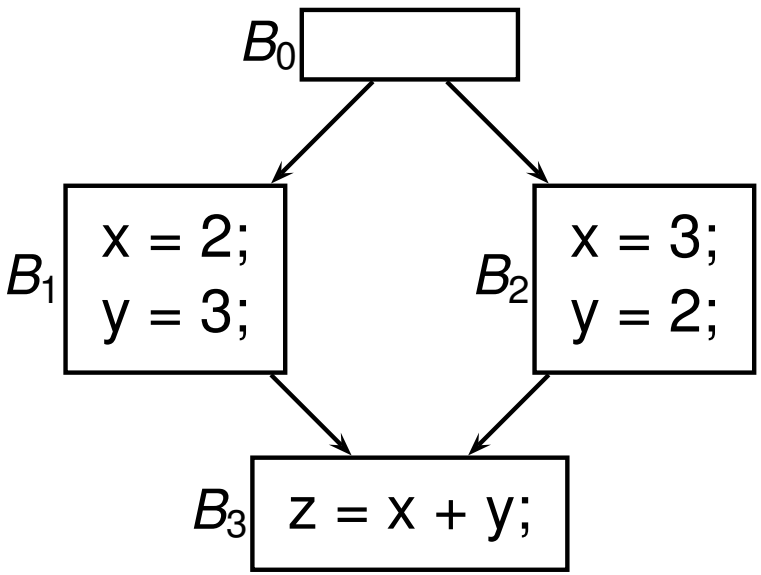

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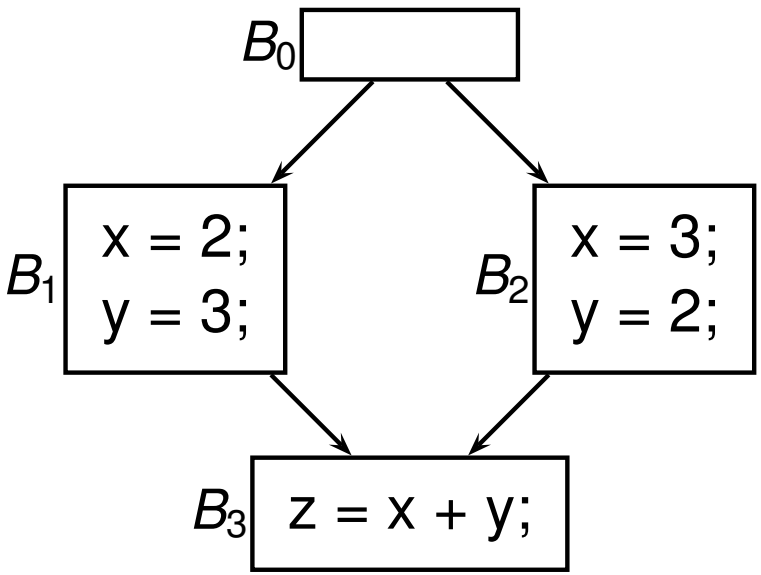

- All paths:
  - $B_0 \rightarrow B_1 \rightarrow B_3$
  - $B_0 \rightarrow B_2 \rightarrow B_3$
- Value of z is 5 along both the paths.

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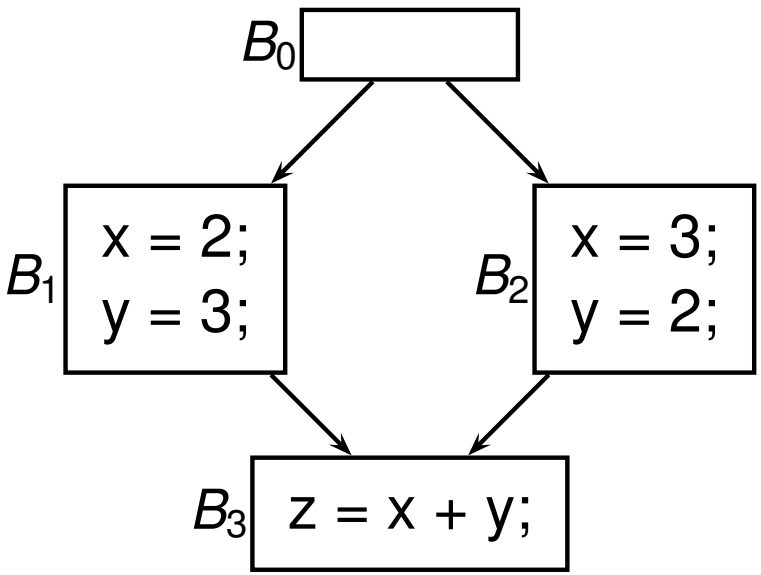

- All paths:
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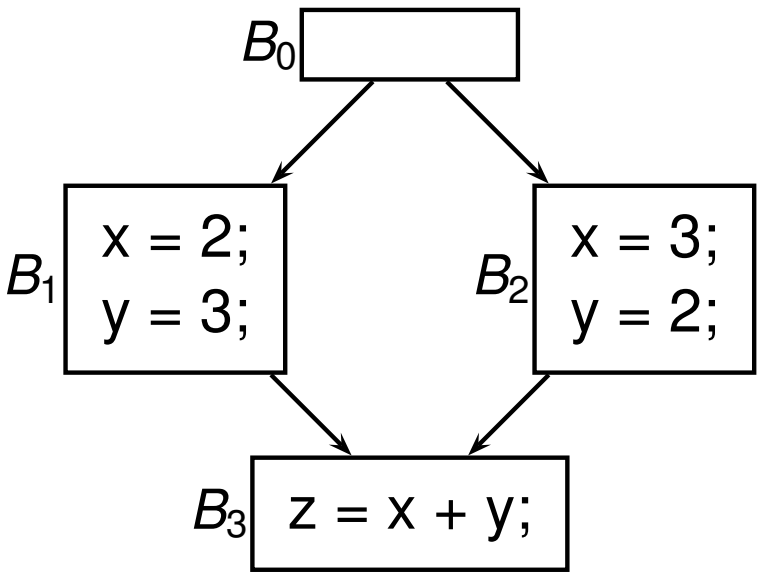
  - $B_0 \rightarrow B_2 \rightarrow B_3$
- Value of z is 5 along both the paths.
- MOP value for z is 5.

Nondistributivity of CP



- All paths:
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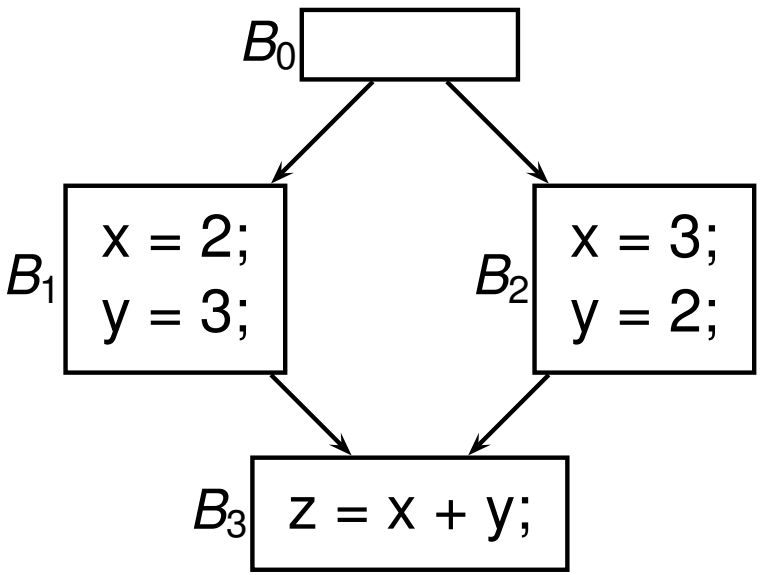
- MOP value for z is 5.
- MFP value for z is *NAC*. (Exercise)

Nondistributivity of CP



- All paths:
  - $B_0 \rightarrow B_1 \rightarrow B_3$
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- Value of z is 5 along both the paths.
- MOP value for z is 5.
- MFP value for z is *NAC*. (Exercise)
- MFP value $\neq$ MOP value (MFP $<$ MOP)