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An MILP Encoding for Efficient Verification of Quantized DNNs

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Introduction



Figure 4.1: Latency-vs-accuracy tradeoff of float vs. integer-only MobileNets on ImageNet using Snapdragon 835 big cores.

B. Jacob, S. Kligys, B. Chen, M. Zhu, M. Tang, A. Howard, H. Adam, and D. Kalenichenko, "Quantization and Training of Neural Networks for Efficient Integer-Arithmetic-Only Inference," in IEEE Conference on Computer Vision and Pattern Recognition (CVPR), 2018

DM	Туре	Precision	Recall	LITTLE (ms)	big (ms)
100%	floats	68%	76%	711	337
	8 bits	66%	75%	372	154
50%	floats	65%	70%	233	106
	8 bits	62%	70%	134	56
25%	floats	56%	64%	100	44
	8 bits	54%	63%	67	28

Table 4.5: Face detection accuracy of floating point and integeronly quantized models. The reported precision / recall is averaged over different precision / recall values where an IOU of xbetween the groundtruth and predicted windows is considered a correct detection, for x in $\{0.5, 0.55, \ldots, 0.95\}$. Latency (ms) of floating point and quantized models are reported on Qualcomm Snapdragon 835 using a single LITTLE and big core, respectively.

Adversarial Attacks



M. Giacobbe et al, "How Many Bits Does it Take to Quantize Your Neural Network?" TACAS, 2020



 $+.007 \times$

x "panda" 57.7% confidence

Goodfellow et al, Explaining and Harnessing Adversarial Examples, ICLR 2015





sign $(\nabla_{\boldsymbol{x}} J(\boldsymbol{\theta}, \boldsymbol{x}, y))$ "nematode" 8.2% confidence



 x_1 [2.09375,3]



 x_2 [0.5,1]



 x_2

[0.5, 1]



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• $Q4.4 \cdot Q4.4 = Q8.8 \rightarrow Q8.8 \Rightarrow Q8.8 \Rightarrow 4 = Q8.4 \rightarrow Q4.4 = \min(255, \max(-256, Q8.4))$



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$$y_1 = 33/2^4 = 2.0625$$

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 $\eta - \text{offset} \leq \overline{\zeta} \rightarrow 13 \leq \overline{\zeta}$ $\overline{\zeta} \le \eta \to \overline{\zeta} \le 13.75$

MILP Encoding of QNN



Floating-point DNN → Fixed-point DNN





• Encode input bounds

$$33 \leq \overline{\mathbf{A}_{1,1}} \leq 48$$
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• Dot product with weights, add bias

$$\overline{pr_{2,1}} = [16, -16]^{\mathsf{T}} \cdot \overline{\mathbf{X}^{2,1}} + 0$$
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• Apply ReLU

$$\overline{z_{2,1}} = \max(0, \overline{q_{2,1}})$$
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Baranowski et al. "An SMT Theory of Fixed-Point Arithmetic," IJCAR, 2020 Henzinger et al. "Scalable Verification of Quantized Neural Networks," AAAI, 2021

MILP vs BVSMT - MNIST





MILP vs BVSMT - CoAv, TwinStream, ACAS Xu





MILP vs BV2SMT - MNIST & Fashion MNIST

Benchmark	# Props	Time (s) (Mean Median)		# Timeouts	
		MILP	BV2	MILP	BV2
MNIST-C	400	5.53 5.4	90 5	0	82
FASHION-C	400	5.73 5.46	4914	0	206

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- Code & data available at https://github.com/iitkcpslab/QNNV