

CS 335: Notes on Runtime Environments

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- Assembly listings use ATnT syntax
- q is for 64 bits, l is for 32 bits, w is 16 bits, b is 8 bits. S is for 32-bit floating point and l is for 64-bit floating point.
- %rbp is the frame pointer on x86_64

Understanding Stack Manipulation during Function Calls

Source Code

```
void proc(long a1, long* a1p, int a2, int* a2p, short a3, short* a3p, char a4, char* a4p) {
    *a1p += a1;
    *a2p += a2;
    *a3p += a3;
    *a4p += a4;
}
```

Compilation Command and Assembly

```
gcc -S -m64 -fno-asynchronous-unwind-tables -fno-exceptions proc-call.c
```

proc:

```
endbr64 ; ignore, related to CET
; start of prologue
pushq %rbp ; save caller's %rbp in callee stack, callee-saved
movq %rsp, %rbp ; %rbp is now the base of the new stack frame
; %rsp will change, so we can use constant offsets from %rbp while generating code for
; local variables and function parameters
; end of prologue
movq %rdi, -8(%rbp) ; save a1 (8 bytes) on the stack
movq %rsi, -16(%rbp) ; save a1p (8 bytes) on the stack
movl %edx, -20(%rbp) ; save a2 (4 bytes) on the stack
movq %rcx, -32(%rbp) ; save a2p (8 bytes) on the stack
; stack pointer is aligned to 16 bytes before making a call
movl %r8d, %eax ; move a3 (lower 32-bits of %r8) to %eax
movq %r9, -40(%rbp) ; save a3p (8 bytes) on the stack
movl 16(%rbp), %edx ; move a4 (4 bytes) to %edx
movw %ax, -24(%rbp) ; move a3 (lower 2 bytes of %eax) to stack
movl %edx, %eax ; move a4 to %eax
movb %al, -44(%rbp) ; save a4 (4 bytes) on the stack
```

```

; *a1p += a1;
movq  -16(%rbp), %rax ; move a1p to %rax
movq  (%rax), %rdx ; move *a1p to %rdx
movq  -8(%rbp), %rax ; %rax contains a1
addq  %rax, %rdx ; add a1 to *a1p, store in %rdx
movq  -16(%rbp), %rax ; move a1p to %rax
movq  %rdx, (%rax) ; store a1+*a1p to the location a1p
; *a2p += a2;
movq  -32(%rbp), %rax ; move a2p to %rax
movl  (%rax), %edx ; move *a2p to %edx
movl  -20(%rbp), %eax ; move a2 to %eax
addl  %eax, %edx ; add a2 and *a2p, store in %edx
movq  -32(%rbp), %rax ; move a2p to %rax
movl  %edx, (%rax) ; store a2+*a2p to the location a2p
; *a3p += a3;
movq  -40(%rbp), %rax ; move a3p to %rax
movzwl (%rax), %eax ; move *a3p (low order 2 bytes) to %eax
movl  %eax, %edx ; move *a3p to %edx
movzwl -24(%rbp), %eax ; move a3 to %eax
addl  %edx, %eax ; add a3 and *a3p, store in %eax
movl  %eax, %edx ; move a3+*a3p to %edx
movq  -40(%rbp), %rax ; move a4p to %rax
movw  %dx, (%rax) ; store a3+*a3p (2 bytes) to the location a3p
; *a4p += a4;
movq  24(%rbp), %rax ; move a4p to %rax
movzbl (%rax), %eax ; move *a4p (low order 1 byte) to %eax
movl  %eax, %edx ; move *a4p to %edx
movzbl -44(%rbp), %eax ; move a4 to %eax
addl  %edx, %eax ; add *a4p and a4, store in %eax
movl  %eax, %edx ; move a4+*a4p to %edx
movq  24(%rbp), %rax ; move a4p to %rax
movb  %dl, (%rax) ; store a4+*a4p (1 byte) to the location a4p
; start of epilogue
nop
popq  %rbp ; restore the callee-saved register
; end of epilogue
ret

```

Compilation Command and Assembly

```
gcc -O2 -S -m64 -fno-asynchronous-unwind-tables -fno-exceptions proc-call.c
```

-O2 has the effect of including -fomit-frame-pointer.

proc:

```

endbr64 ; ignore, related to CET
movq  16(%rsp), %rax ; save *a4p to %rax
addq  %rdi, (%rsi) ; *a1p+=a1;
addl  %edx, (%rcx) ; *a2p += a2;
movl  8(%rsp), %edx ; move a4 to %dl (1 byte)
addw  %r8w, (%r9) ; *a3p += a3; (2 bytes)

```

```

    addb    %dl, (%rax) ; *a4p += a4; (1 byte)
    ret

```

Use of Callee-Saved Registers

Source Code

```

long proc2(long);

long proc1(long x, long y) {
    long u = proc2(y);
    long v = proc2(x);
    return u + v;
}

```

Compilation Command and Assembly

```
gcc -S -m64 -fno-asynchronous-unwind-tables -fno-exceptions callee-saved-regs.c
```

```

proc1:
    ; x is in %rdi, y is in %rsi
    ; start of prologue
    pushq   %rbp ; save caller's %rbp in callee stack, callee-saved register
    movq    %rsp, %rbp ; %rbp is now the base of the new stack frame
    ; %rsp will change, so we can use constant offsets from %rbp while generating code for
    ; local variables and function parameters
    ; end of prologue
    subq    $32, %rsp ; allocate space on stack
    movq    %rdi, -24(%rbp) ; save x on stack
    movq    %rsi, -32(%rbp) ; save y on stack
    movq    -32(%rbp), %rax ; move y to %rax
    movq    %rax, %rdi ; move y to %rdi, prepare first parameter
    call    proc2@PLT
    movq    %rax, -16(%rbp) ; save return value u to stack
    movq    -24(%rbp), %rax ; move x to %rax
    movq    %rax, %rdi ; move x to %rdi, prepare first parameter
    call    proc2@PLT
    movq    %rax, -8(%rbp) ; save return value v to stack
    movq    -16(%rbp), %rdx ; pop u
    movq    -8(%rbp), %rax ; pop v
    addq    %rdx, %rax ; store u+v in %rax
    leave  ; epilogue
    ret

```

Compilation Command and Assembly

```
gcc -O2 -S -m64 -fno-asynchronous-unwind-tables -fno-exceptions callee-saved-regs.c
```

```

proc1:
    pushq   %rbp ; save caller's %rbp in callee stack, callee-saved register
    movq    %rdi, %rbp ; save x on stack
    movq    %rsi, %rdi ; move y to %rdi
    pushq   %rbx ; save caller's %rbp in callee stack, callee-saved register

```

```

; cannot use 128-byte red zone below stack pointer because proc1 is not a leaf function,
; so proc2 can clobber the red zone
subq   $8, %rsp ; reserve space on stack for locals
call   proc2@PLT
movq   %rbp, %rdi ; move x to %rdi
movq   %rax, %rbx ; move u to %rbx
call   proc2@PLT
addq   $8, %rsp ; deallocate space from stack
addq   %rbx, %rax ; store u+v in %rax
popq   %rbx ; restore the callee-saved register
popq   %rbp ; restore the callee-saved register
ret

```

References

- Randal E. Bryant and David R. O'Hallaron. Computer Systems: A Programmer's Perspective, 3rd edition.
- Where the top of the stack is on x86
- Stack frame layout on x86-64
- Position Independent Code (PIC) in shared libraries
- C++ Internals Memory layout
- x86 Disassembly/Calling Conventions
- Calling conventions for different C++ compilers and operating systems