

## Lecture 17: September 8, 2017

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## 17.1 Recap

We defined the Arrovian social welfare function (ASWF) to be a mapping from the set of all preference profiles of  $n$  agents to a single preference profile. Hence it is a function  $F : \mathcal{R}^n \rightarrow \mathcal{R}$ , where  $\mathcal{R}$  is the set of all possible orderings over  $|A|$  candidates. Two desirable properties that were listed were

**Definition 17.1** A Social Welfare Function  $F$  satisfies weak Pareto (WP) if

$$\forall a, b \in A, [aP_i b, \forall i \in N] \implies [a\hat{F}(R)b].$$

**Definition 17.2** A Social Welfare Function  $F$  satisfies strong Pareto if

$$\forall a, b \in A, [aR_i b, \forall i \in N, \exists j, aP_j b] \implies [a\hat{F}(R)b].$$

Clearly, strong Pareto  $\implies$  weak Pareto. The other desirable property in the ASWF setup is *independence of irrelevant alternatives*.

## 17.2 Independence of Irrelevant Alternatives

This property is the crux of Arrow's impossibility result. It is a property that connects two different preference profiles.

Two preferences of player  $i$ , say  $R_i$  and  $R'_i \in \mathcal{R}$  are said to *agree* over  $\{a, b\}$  if for agent  $i$

- $aP_i b \Leftrightarrow aP'_i b$
- $bP_i a \Leftrightarrow bP'_i a$
- $aI_i b \Leftrightarrow aI'_i b$

We denote this using the notation  $R_i|_{a,b} = R'_i|_{a,b}$ . Two preference profiles  $R, R'$  agree if for every  $i \in N$ ,  $R_i|_{a,b} = R'_i|_{a,b}$  and is denoted by

$$R|_{a,b} = R'|_{a,b}.$$

**Definition 17.3 (Independence of Irrelevant Alternatives)** An ASWF  $F$  satisfies independence of irrelevant alternatives (IIA) if for all  $a, b \in A$

$$[R|_{a,b} = R'|_{a,b}] \implies [F(R)|_{a,b} = F(R')|_{a,b}].$$

**Illustration** Consider an ASWF  $F$ , where given the position of the ranking for every agent, some scores are assigned to the candidates. Formally, say the score vector is  $(s_1, s_2, s_3, \dots, s_m)$ ,  $s_i \geq s_{i+1}, i = 1, 2, \dots, m - 1, s_i \geq 0, \forall i \in N$ . Finally all scores of a particular candidate are added and the final ranking is based on the decreasing order of these scores. This is one special class of ASWF.

Some well-known scoring rules are described below:

- **Plurality:** In this case we assign top score, i.e 1 to  $s_1$  and 0 to all others, So  $s_1 = 1, \text{and } s_2 = s_3 = \dots = s_m = 0$ .

**Question:** Does plurality satisfy IIA?

Consider two preference profiles  $R$  and  $R'$ . The preferences of 4 voters are as follows.

|                 |                 |
|-----------------|-----------------|
| $R$             | $R'$            |
| $a \ a \ c \ d$ | $d \ c \ b \ b$ |
| $b \ c \ b \ c$ | $a \ a \ c \ a$ |
| $c \ b \ a \ b$ | $b \ b \ a \ d$ |
| $d \ d \ d \ a$ | $c \ d \ d \ c$ |

Plurality gives a social ordering between  $a$  and  $b$  as:

$$a \hat{F}^{Plu}(R)b, \text{ and } b \hat{F}^{Plu}(R')a.$$

However, we see that the ordering of  $a$  and  $b$  remains same for every agent in  $R$  and  $R'$ . IIA would require that the social ordering remain unchanged, which does not happen for plurality. Thus we conclude that plurality does not satisfy IIA.

- **Borda:** The scoring rule in this case is:  $s_1 = m - 1, s_2 = m - 2, \dots, s_{m-1} = 1, s_m = 0$ .
- **Veto:** The scoring rule is:  $s_1 = s_2 = \dots = s_{m-1} = 1, s_m = 0$ . We can check by suitable examples that neither Borda nor veto satisfies IIA.
- **Dictatorial:** A voting rule is dictatorial if it always selects the preference ordering of a distinguished agent, whom we call the *dictator*. Thus it is trivial that a dictatorial voting rule satisfies IIA.

We are now going to present a classic result in social choice.

**Theorem 17.4 (Arrow 1950)** For  $|A| \geq 3$ , if an ASWF  $F$  satisfies weak Pareto and IIA then it must be dictatorial.

**Proof:** The proof of the following two lemmas will lead us to eventually prove Arrow's theorem. Informally we state the basic statements of the lemmas as follows.

1. *Field Expansion Lemma:* if a group  $G \subseteq N, G \neq \emptyset$  is *decisive* over  $a, b$ , then it is decisive over all pairs of alternatives. Informally, a decisive group is a group such that if every agent in that group agrees on a ranking between a pair of alternatives, that ranking is reflected in the social ranking. Therefore, with this lemma, it is enough to call a group *decisive* since it implies that it is decisive over all pairs of alternatives.
2. *Group Contraction Lemma:* if a group  $G$  is decisive, there exists a strict subset of  $G$  that is also decisive.

First we define decisiveness formally.

**Definition 17.5** Given  $F : \mathcal{R}^n \rightarrow \mathcal{R}$ . Let  $G \subseteq N, G \neq \emptyset$ .

1.  $G$  is almost decisive over  $a, b$  if

$$[aP_i b, \forall i \in G, \text{ and } bP_j a, \forall j \notin G] \implies [a\hat{F}(R)b]$$

2.  $G$  is called decisive over  $a, b$  if

$$[aP_i b, \forall i \in G] \implies [a\hat{F}(R)b]$$

We will use the notation  $\bar{D}_G(a, b)$  to denote that  $G$  is almost decisive over  $a, b$  and  $D_G(a, b)$  to denote that  $G$  is decisive over  $a, b$ . Clearly,  $D_G(a, b) \implies \bar{D}_G(a, b)$ .

**Lemma 17.6 (Field Expansion)** Let  $F$  satisfies weak Pareto and IIA then  $\forall a, b, x, y, a \neq b, x \neq y$ , we have

$$\bar{D}_G(a, b) \implies D_G(x, y).$$

**Proof:** We consider the following set of exhaustive cases to prove this lemma.

1.  $\bar{D}_G(a, b) \implies D_G(a, y)$  where  $y \neq a, b$
2.  $\bar{D}_G(a, b) \implies D_G(x, b)$  where  $x \neq a, b$
3.  $\bar{D}_G(a, b) \implies D_G(x, y)$  where  $x \neq a, b$  and  $y \neq a, b$
4.  $\bar{D}_G(a, b) \implies D_G(x, a)$  where  $x \neq a, b$
5.  $\bar{D}_G(a, b) \implies D_G(b, y)$  where  $y \neq a, b$
6.  $\bar{D}_G(a, b) \implies D_G(b, a)$
7.  $\bar{D}_G(a, b) \implies D_G(a, b)$

**Case 1:** Given:  $\bar{D}_G(a, b)$ , we need to show  $D_G(a, y)$ . Pick arbitrary  $R$  such that,

$$aP_i y, \forall i \in G, \text{ need to show } a\hat{F}(R)y.$$

Construct  $R'$  as follows.

$$\begin{array}{cc} G & N \setminus G \\ a \succ b \succ y & b \succ a \text{ and } b \succ y \end{array}$$

Where  $a \succ b$  denotes  $a$  is more preferred than  $b$ . For the agents in  $N \setminus G$ , we ensure that the ranking of  $a$  and  $y$  remain identical to the ranking of these two alternatives in  $R$ . Therefore we have

$$R|_{a,y} = R'|_{a,y}.$$

Now since  $aR'_i b, \forall i \in G$  and  $bR'_j a, \forall j \notin G$ , by definition of  $\bar{D}_G(a, b)$  we conclude that  $a\hat{F}(R')b$ . Since  $b$  is preferred over  $y$  by all agents in  $N$ , WP implies that  $b\hat{F}(R')y$ . Using transitivity of  $F(R')$ , we have,  $a\hat{F}(R')y$ . Since the relative ranking of  $a$  and  $y$  in  $R$  and  $R'$  are same, using IIA we get  $a\hat{F}(R)y$ .

**Case 2:** Given:  $\bar{D}_G(a, b)$ , we need to show  $D_G(x, b)$ . Pick arbitrary  $R$  such that,

$$xP_i b, \forall i \in G, \text{ need to show } x\hat{F}(R)b.$$

Construct  $R'$  as follows.

$$\begin{array}{cc} G & N \setminus G \\ x \succ a \succ b & x \succ a \text{ and } b \succ a \end{array}$$

For the agents in  $N \setminus G$ , we ensure that the ranking of  $x$  and  $b$  remain identical to the ranking of these two alternatives in  $R$ . Therefore we have

$$R|_{x,b} = R'|_{x,b}.$$

Now since  $aR'_i b, \forall i \in G$  and  $bR'_j a, \forall j \notin G$ , by definition of  $\bar{D}_G(a, b)$  we conclude that  $a\hat{F}(R')b$ . Since  $x$  is preferred over  $a$  by all agents in  $N$ , WP implies that  $x\hat{F}(R')a$ . Using transitivity of  $F(R')$ , we have,  $x\hat{F}(R')b$ . Since the relative ranking of  $a$  and  $y$  in  $R$  and  $R'$  are same, using IIA we get  $x\hat{F}(R)b$ .

**Case 3:**

$$\begin{aligned} \bar{D}_G(a, b) &\implies D_G(a, y), y \neq a, b && \text{(by Case 1)} \\ &\implies \bar{D}_G(a, y) && \text{(by definition)} \\ &\implies \bar{D}_G(x, y), \text{ as } x \neq a, y && \text{(by Case 2)} \end{aligned}$$

**Case 4:**

$$\begin{aligned} \bar{D}_G(a, b) &\implies D_G(x, b), x \neq a, b && \text{(by Case 2)} \\ &\implies \bar{D}_G(x, b) && \text{(by definition)} \\ &\implies \bar{D}_G(x, a), \text{ as } a \neq b, x && \text{(by Case 1)} \end{aligned}$$

**Case 5:**

$$\begin{aligned} \bar{D}_G(a, b) &\implies D_G(a, y), y \neq a, b && \text{(by Case 1)} \\ &\implies \bar{D}_G(a, y) && \text{(by definition)} \\ &\implies \bar{D}_G(b, y), \text{ as } b \neq a, y && \text{(by Case 2)} \end{aligned}$$

**Case 6:**

$$\begin{aligned} \bar{D}_G(a, b) &\implies D_G(x, b), x \neq a, b && \text{(by Case 2)} \\ &\implies \bar{D}_G(x, b) && \text{(by definition)} \\ &\implies \bar{D}_G(a, b), \text{ as } a \neq b, x && \text{(by Case 2)} \end{aligned}$$

**Case 7:**

$$\begin{aligned} \bar{D}_G(a, b) &\implies D_G(b, y), y \neq a, b && \text{(by Case 5)} \\ &\implies \bar{D}_G(b, y) && \text{(by definition)} \\ &\implies \bar{D}_G(b, a), \text{ as } a \neq b, y && \text{(by Case 1)} \end{aligned}$$

In the next class we will prove the *Group Contraction Lemma* to complete our proof of *Arrow's Theorem*. ■