

Lecture 16: September 6, 2017

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16.1 Recap

In the last class we continued our discussion on *mechanism design*. We first discussed *social choice functions (SCF)*, utility functions for SCFs and their types, defined *mechanism design (MD)* and its types. Then we went on to discuss *weakly dominant strategy* w.r.t. mechanism and conditions under which a SCF is implemented in dominant strategy. It was followed by discussion on *strategyproofness* or *dominant strategy incentive compatibility (DSIC)* and *revelation principle*.

Before going forward let's revisit the definition of DSIC as it will be used later to illustrate a subtle point w.r.t. Bayesian incentive compatibility. The definition is as follows:

Definition 16.1 A direct mechanism $\langle \Theta, f \rangle$ is strategyproof or dominant strategy incentive compatible (DSIC) if

$$u_i(f(\theta_i, \tilde{\theta}_{-i}), \theta_i) \geq u_i(f(\theta'_i, \tilde{\theta}_{-i}), \theta_i), \quad \forall \theta_i, \theta'_i \in \Theta_i, \forall \tilde{\theta}_{-i} \in \Theta_{-i}, \forall i \in N. \quad (16.1)$$

This definition states that irrespective of whether others are reporting their types to the central authority truthfully or not, for any agent i , reporting her type truthfully is a weakly dominant strategy.

16.2 Relating Mechanism Design to Bayesian Games

Suppose the types are generated from a common prior P and type of any player i , θ_i is revealed only to the respective players. We can consider the mechanism design scenario in a Bayesian game setup, where the Bayesian game is as follows.

$$\langle N, (A_i)_{i \in N}, (\Theta_i)_{i \in N}, P, (\Gamma_\theta)_{\theta \in \Theta} \rangle.$$

- N = Set of Players $\{1, 2, \dots, n\}$
- M_i = Message space of Player i (Corresponding to Action set in Bayesian Games)
- Θ_i = Set of type for player i
- P = Common prior over $\times_{i \in N} \Theta_i$
- $\Gamma_\theta = \langle N, (A_i)_{i \in N}, (u_i(\theta))_{i \in N} \rangle$

Consider the function $m_i : \Theta_i \rightarrow M_i$. We can observe that given any type, the function m_i maps to a member of the message space. It is equivalent to pure strategies of players in a normal Bayesian game. Intuitively every player sends some message to central authority to maximize its utility which are equivalent to strategy a player might have taken in the case of Bayesian games.

16.3 SCF implemented in Bayesian Equilibrium

Definition 16.2 A mechanism $\langle M, g \rangle$ implements SCF f in Bayesian equilibrium if the following two conditions hold.

1. $\exists(m_1, m_2, \dots, m_n)$ s.t. $m_i(\theta_i)$ maximizes the ex-interim utility of agent i , $\forall \theta_i \in \Theta_i, \forall i \in N$, i.e.,

$$\mathbb{E}_{\theta_{-i}|\theta_i}[u_i(g(m_i(\theta_i), m_{-i}(\theta_{-i})), \theta_i)] \geq \mathbb{E}_{\theta_{-i}|\theta_i}[u_i(g(m'_i, m_{-i}(\theta_{-i})), \theta_i)], \quad \forall \theta_i \in \Theta_i, \forall i \in N \quad (16.2)$$

2. $g(m_i(\theta_i), m_{-i}(\theta_{-i})) = f(\theta_i, \theta_{-i})$

Observation: If SCF f is implementable in dominant strategies then f will be implementable in Bayesian equilibrium.

16.4 Bayesian Incentive Compatibility (BIC)

Definition 16.3 A direct mechanism $\langle \Theta, f \rangle$ is Bayesian incentive compatible if

$$\mathbb{E}_{\theta_{-i}|\theta_i}[u_i(f(\theta_i, \theta_{-i}), \theta_i)] \geq \mathbb{E}_{\theta_{-i}|\theta_i}[u_i(f(\theta'_i, \theta_{-i}), \theta_i)], \quad \forall \theta_i, \theta'_i \in \Theta_i, \forall i \in N. \quad (16.3)$$

Observation. From Equations 16.1 and 16.3, we can observe that the condition for a direct mechanism to be DSIC is required to hold for all $\theta_{-i} \in \Theta_{-i}$ but for it to be BIC this condition is not required because we are doing weighted average over all possible θ_{-i} in some sense.

16.5 Revelation Principle for BI SCFs

Theorem 16.4 If a Social Choice Function f is implementable in Bayesian Equilibrium then f is Bayesian Incentive Compatible.

The proof of the above theorem is similar to the proof of Revelation Principle for DSI SCFs and hence left as an exercise.

16.6 Arrovian Social Welfare Function (SWF)

Even before considering strategyproofness, how can we aggregate the individual preferences into a social preference?

To further understand the implication of this question consider the following example. Suppose three friends want to watch a movie every day for a week. Seven movies are chosen – they are the alternatives. But everyone wants to watch their favorite movie first and does not want to watch an already watched movie. Hence, everyone has a different preference order over the movies. However the three friends want to watch the movies *together*. So what is the best order or watching the movies together? In other words, how can we have a social ordering of these alternatives? The question leads us to study of *Arrovian Social Welfare Functions* (ASWF).

Definition 16.5 *Arrovian Social Welfare Function* takes as input individual preferences of different agents outputs a social preference for all the agents.

Notation

- Set of alternatives $A = \{a_1, \dots, a_m\}$.
- Agents $N = \{1, 2, \dots, n\}$.
- $aR_i b$: Alternative a is atleast as good as b for agent i .
- Set of all possible ordering \mathcal{R} .

Properties of R_i (ordering)

- *Completeness*: For every $a, b \in A$ either $aR_i b$ or $bR_i a$
- *Reflexivity*: $\forall a \in A, aR_i a$
- *Transitivity*: If $aR_i b, bR_i c$ then $aR_i c$

Lets divide the relation R_i into two parts P_i (asymmetric part) and I_i (symmetric part) i.e.,

- $aP_i b$: Alternative a is strictly better than b for i .
- $aI_i b$: Alternative a is indifferent to b for i .

Definition 16.6 An ordering R_i is linear if $aR_i b$ and $bR_i a$, then $a = b$, i.e., indifference is not allowed.

16.6.1 Arrovian Social Welfare Function Representation

Using the notations we have discussed have Arrovian Social Welfare Function F is represented as $F : \mathcal{R}^n \rightarrow \mathcal{R}$. Since $F(R)$ is an ordering over the alternatives like R_i we can split it into two parts $\hat{F}(R)$ (asymmetric part) and $\bar{F}(R)$ (symmetric part) where $R = (R_1, \dots, R_n)$.

16.6.2 Desirable Properties for Arrovian SWF

Definition 16.7 A Social Welfare Function F satisfies weak Pareto (WP) if

$$\forall a, b \in A, [aP_i b, \forall i \in N] \implies [a\hat{F}(R)b].$$

Definition 16.8 A Social Welfare Function F satisfies strong Pareto if

$$\forall a, b \in A, [aR_i b, \forall i \in N, \exists j, aP_j b] \implies [a\hat{F}(R)b].$$

Clearly, strong Pareto \implies weak Pareto.