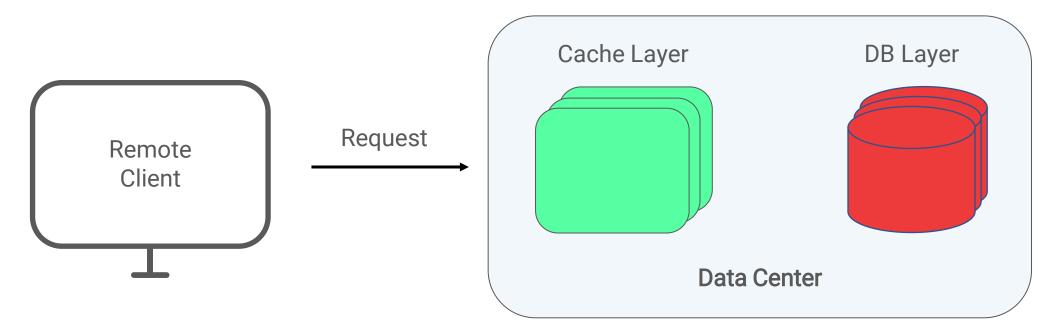


Kangaroo: Caching Billions of Tiny Objects on Flash

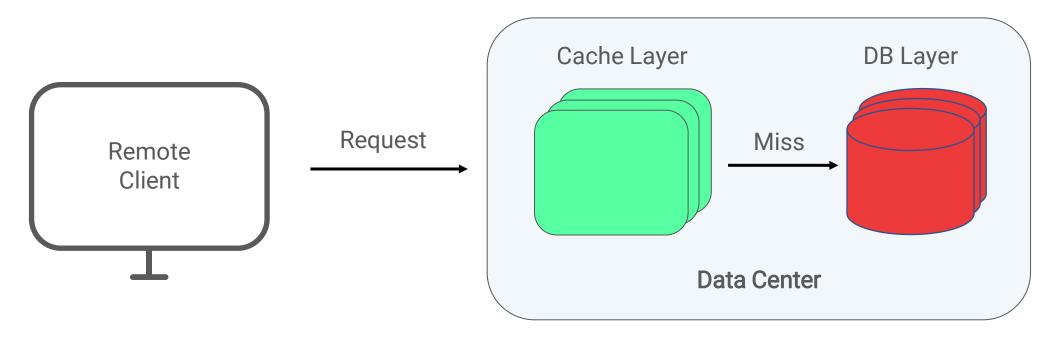
Sara McAllister, Benjamin Berg, Julian Tutuncu-Macias, Juncheng Yang, Sathya Gunasekar, Jimmy Lu, Daniel S. Berger, Nathan Beckmann, Gregory R. Ganger

Shiv Bhushan Tripathi

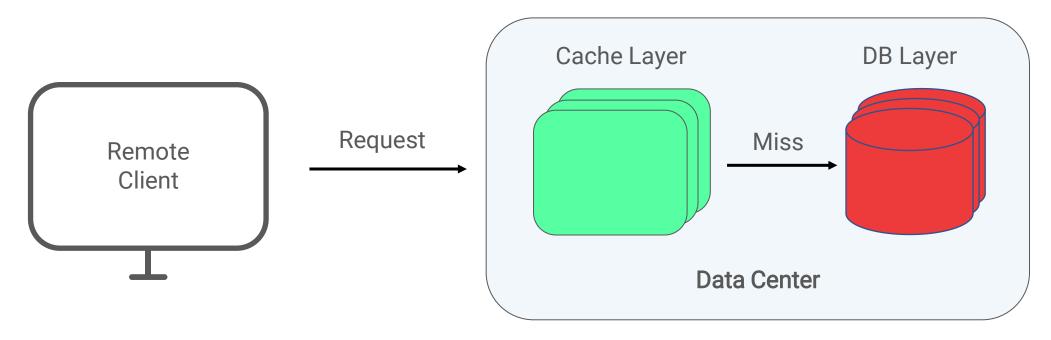
CDOS Talk 16th July 2022



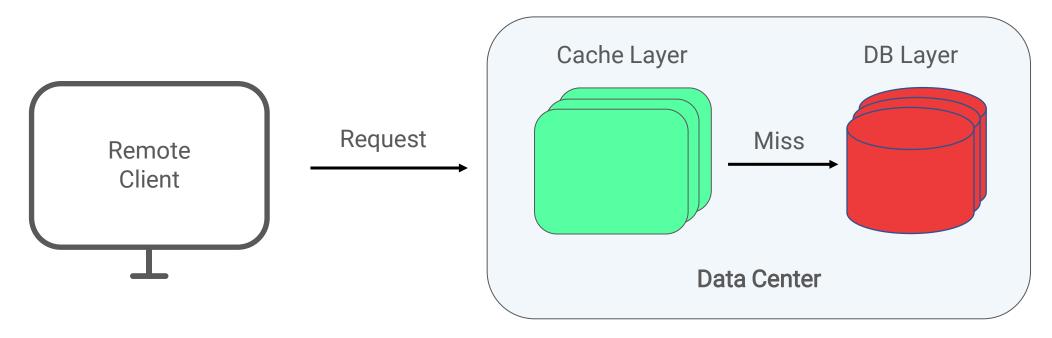
• Added Cache layer



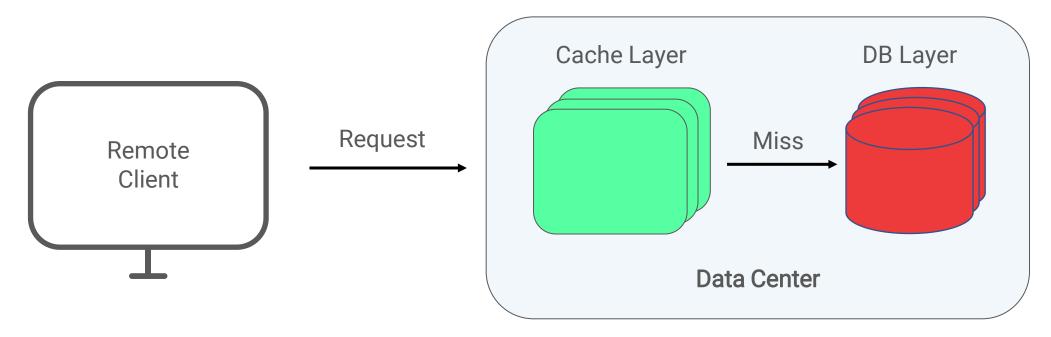
Added Cache layer



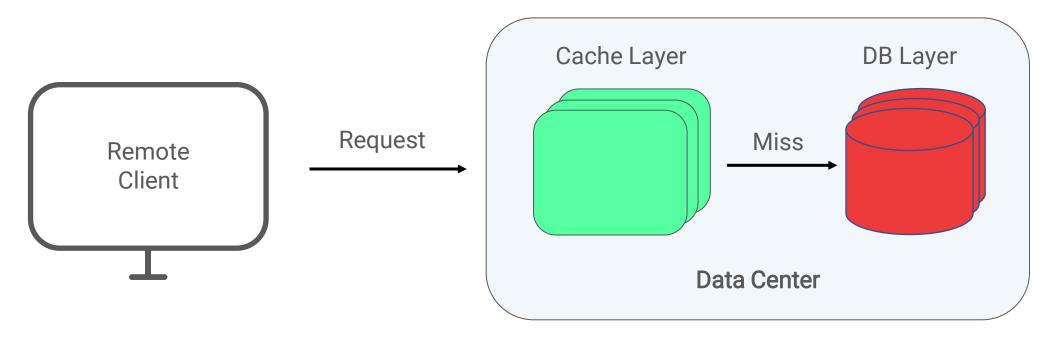
• Added Cache layer: faster response



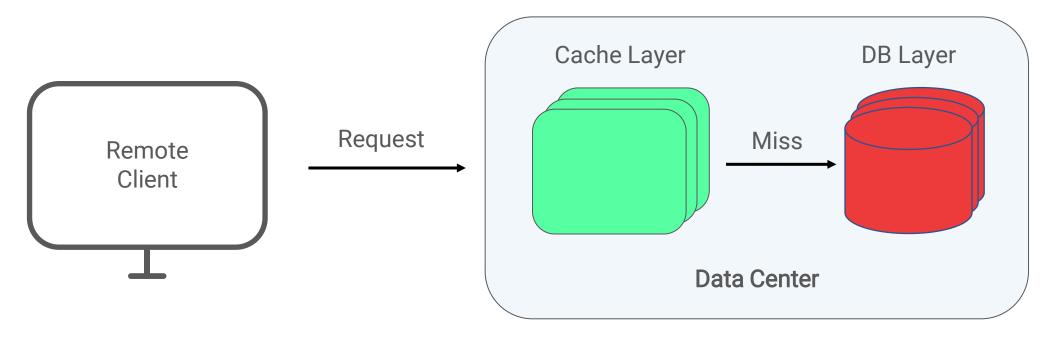
• Added Cache layer: faster response, less traffic to back end DB layer



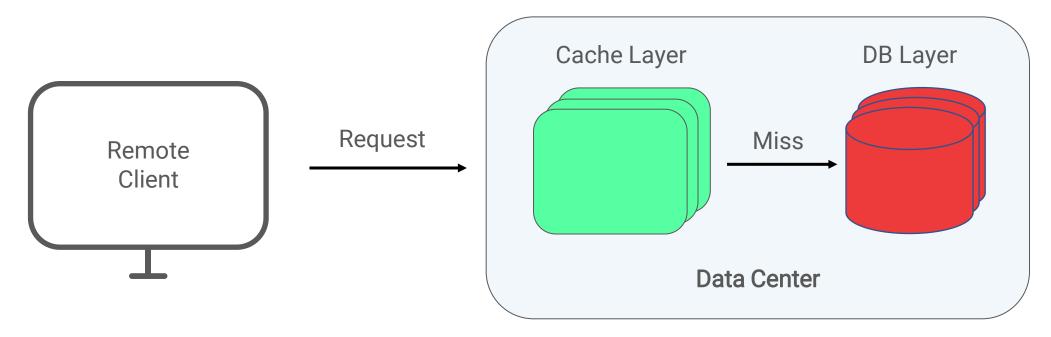
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- Caches should be high performant and low cost



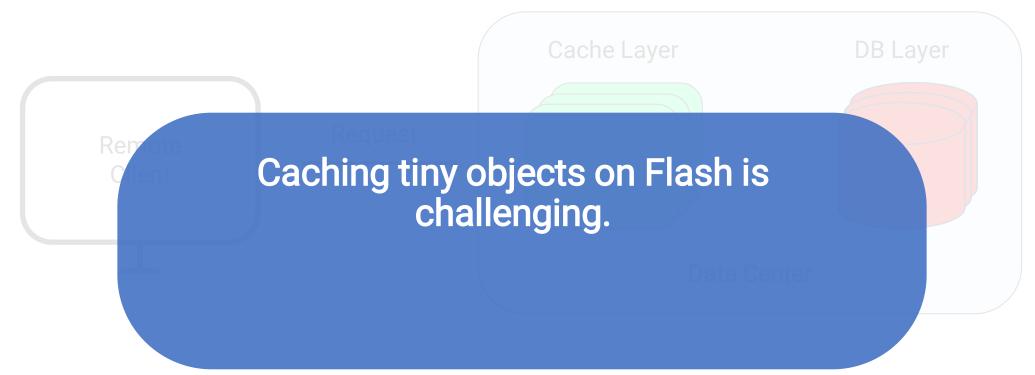
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- Flash as Cache: cheaper than DRAM and NVMs

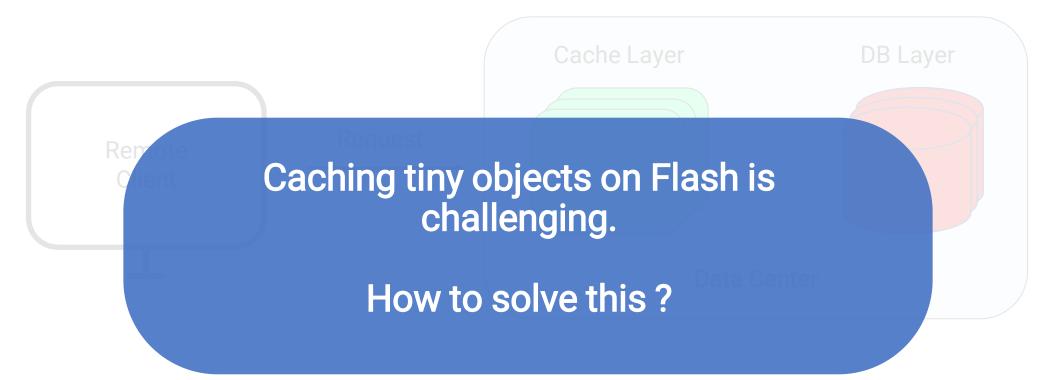


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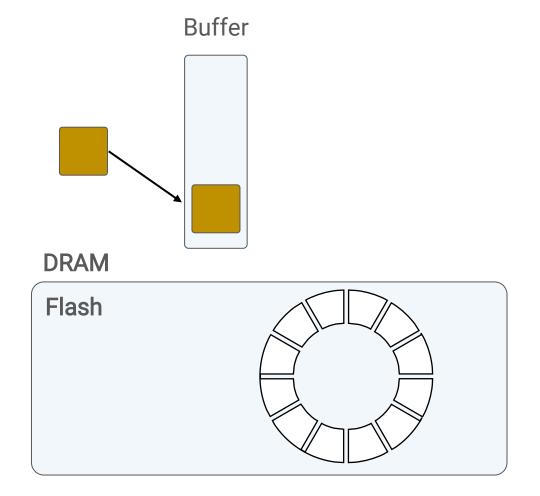
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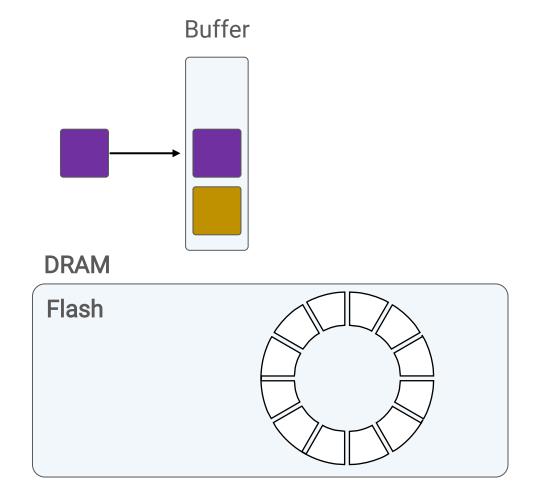
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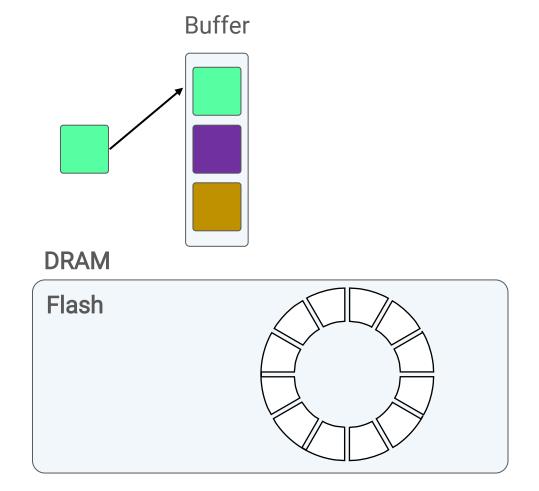
 \circ $\,$ Tiny objects are buffered in DRAM $\,$



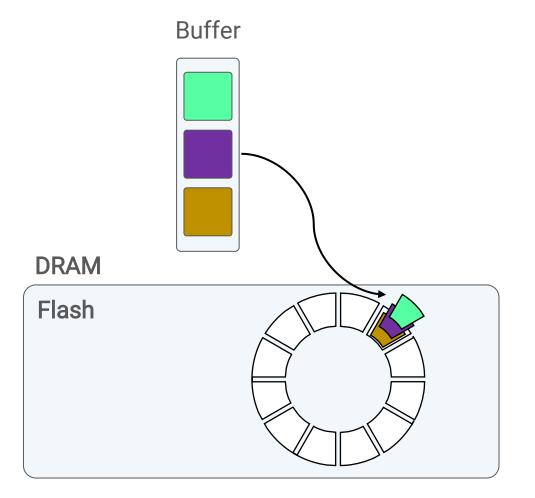
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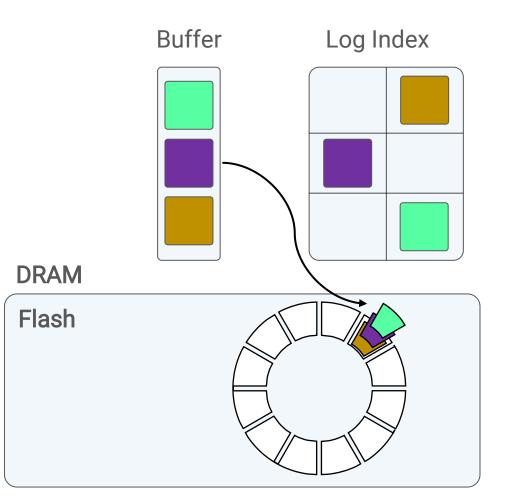
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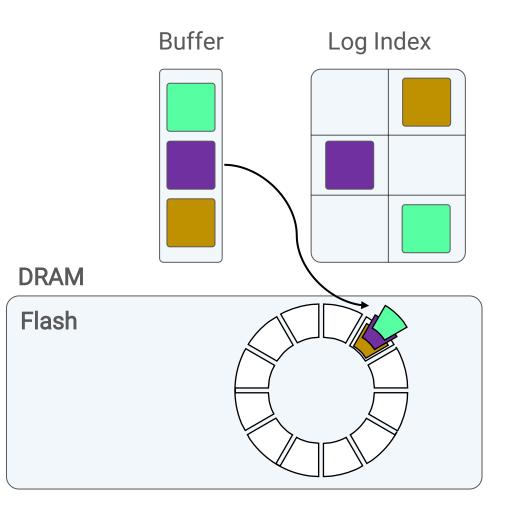
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- A log index in DRAM to track each object



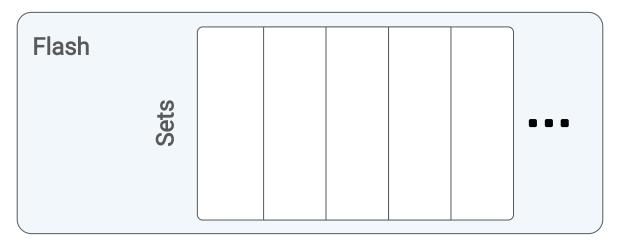
- \circ $\,$ Tiny objects are buffered in DRAM $\,$
- \circ $\,$ Moved to flash, when buffer is full
- A log index in DRAM to track each object
- Reduces flash writes but huge amount of DRAM



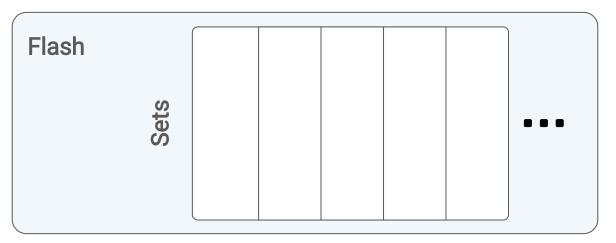
Flash as Cache: Set Associative

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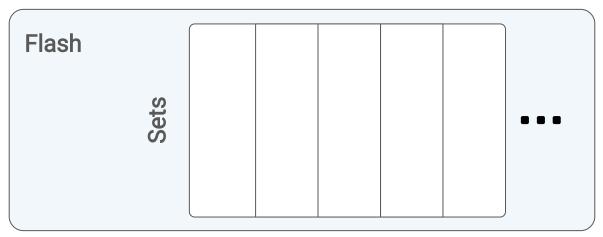
DRAM

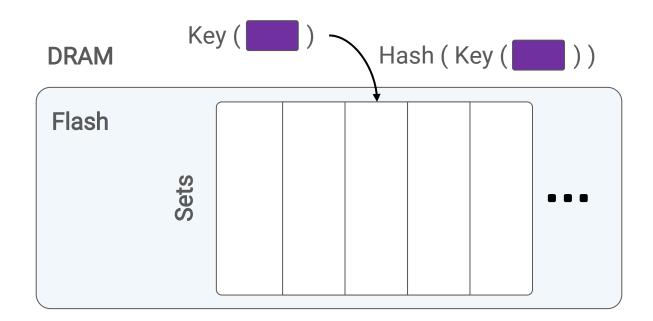


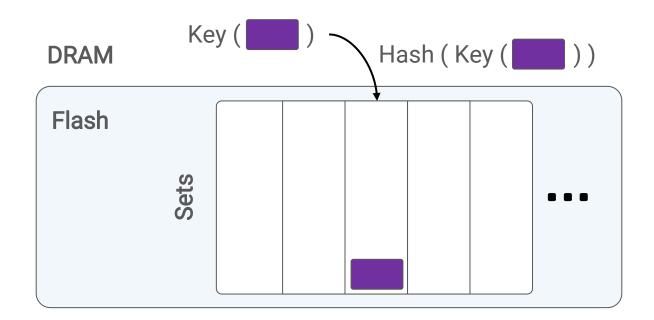
DRAM

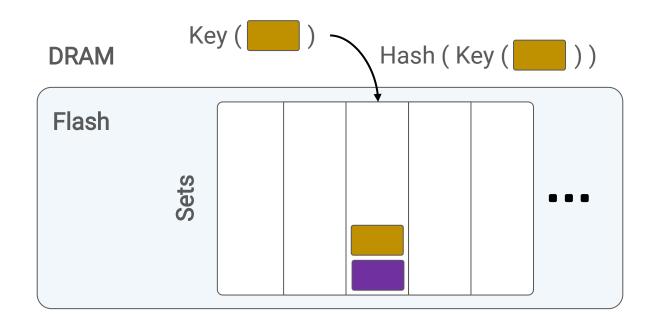


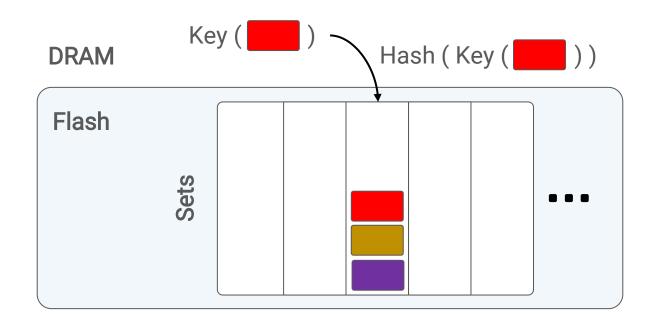


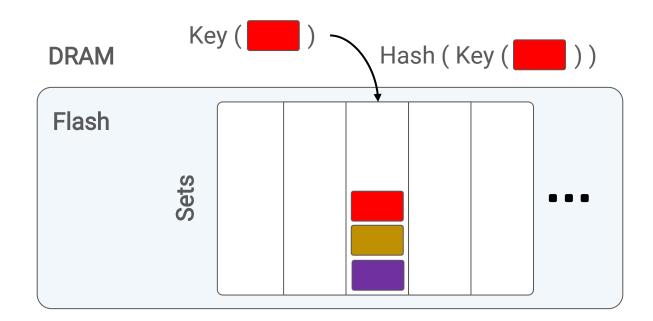




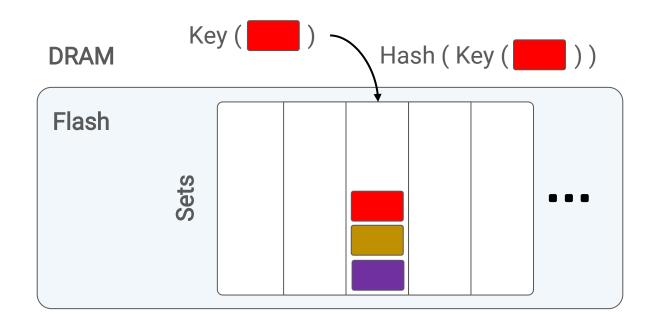








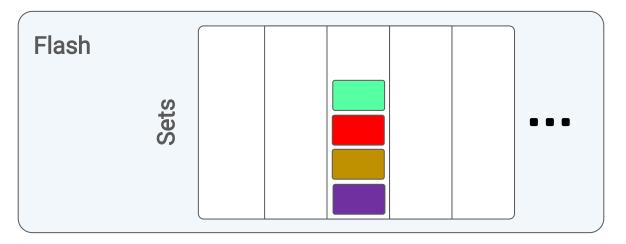
- Objects are inserted into sets using Hashing
- Write amplification is too high for tiny objects



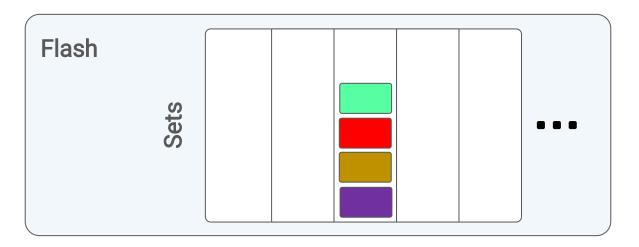
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Write Amplification = 4096 / 100 = 40x

DRAM

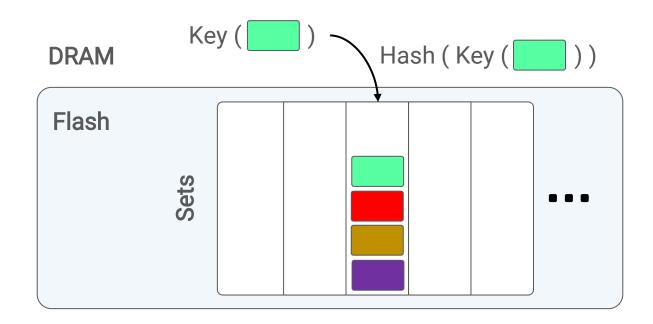




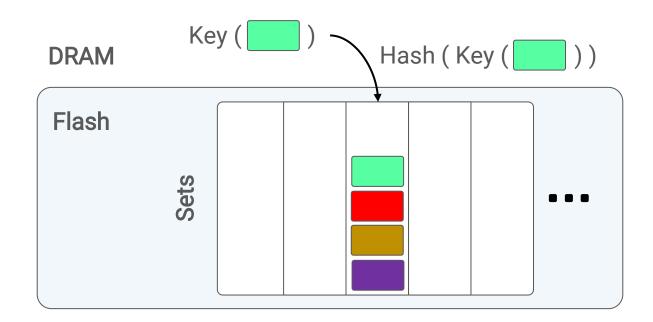


• Objects are searched based on key

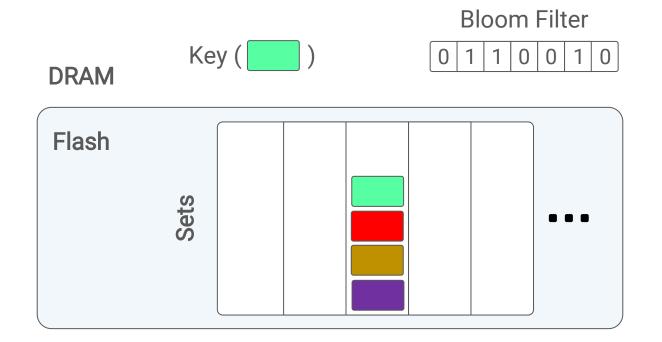
DRAM



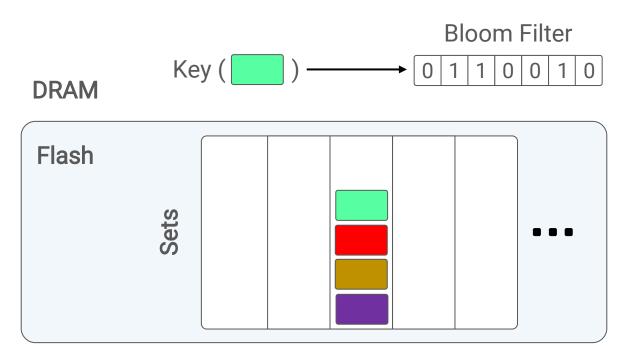
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- Objects are searched based on key
- Have to read all tiny objects in the set (flash read)

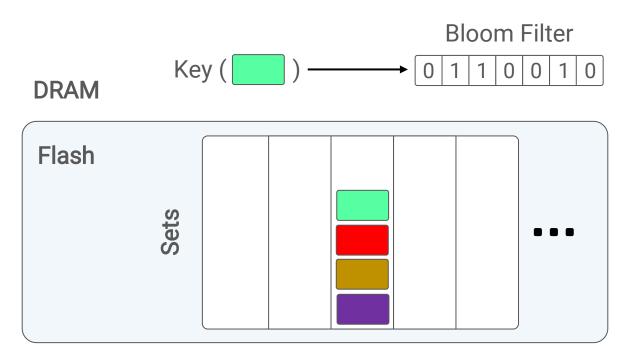


- Objects are searched based on key
- Have to read all tiny objects in the set (flash read): Add bloom filter



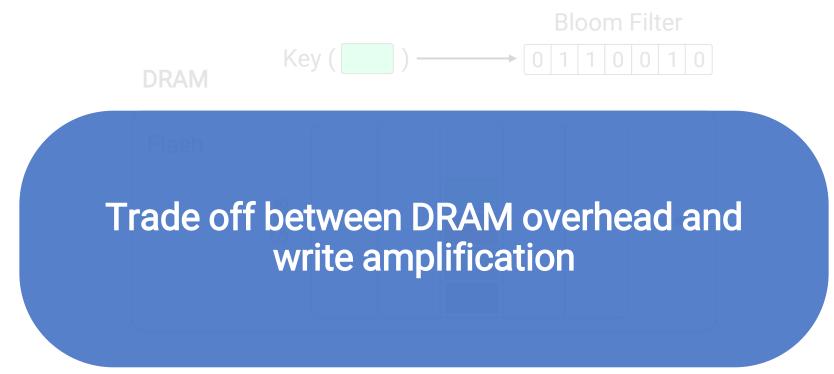
- Objects are searched based on key
- Have to read all tiny objects in the set (flash read): Add bloom filter

Flash as Cache: Set Associative (Search)



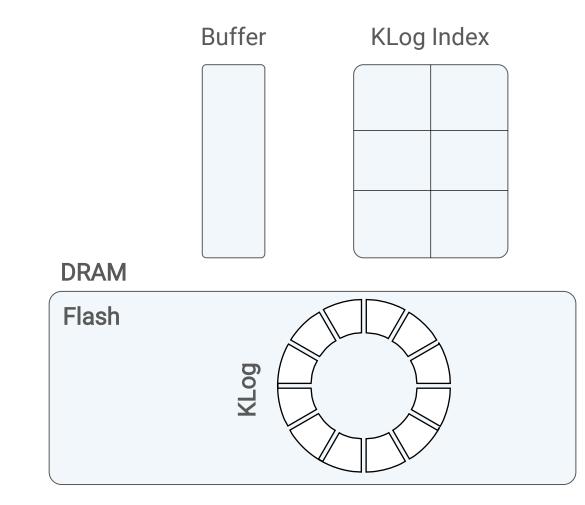
- Objects are searched based on key
- Have to read all tiny objects in the set (flash read): Add bloom filter
- Less DRAM overhead but have huge write amplification

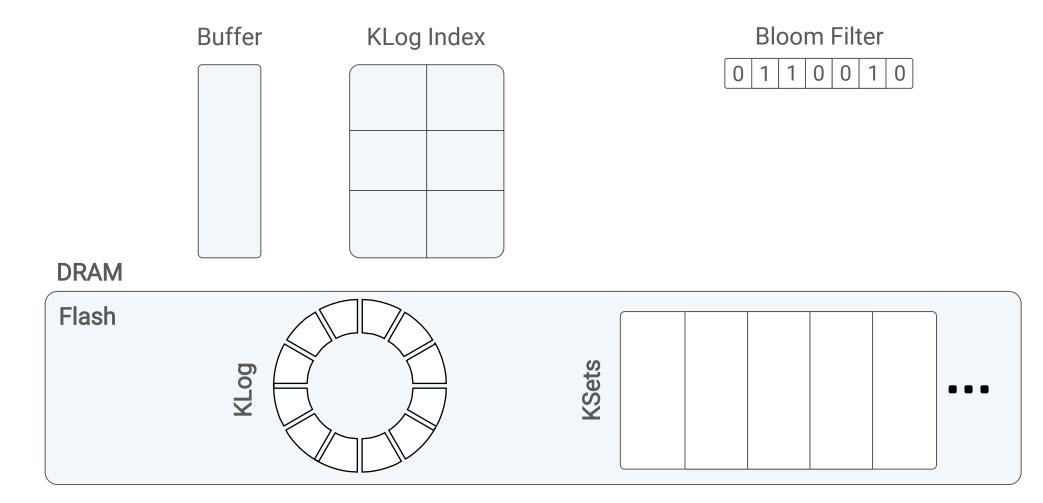
Flash as Cache



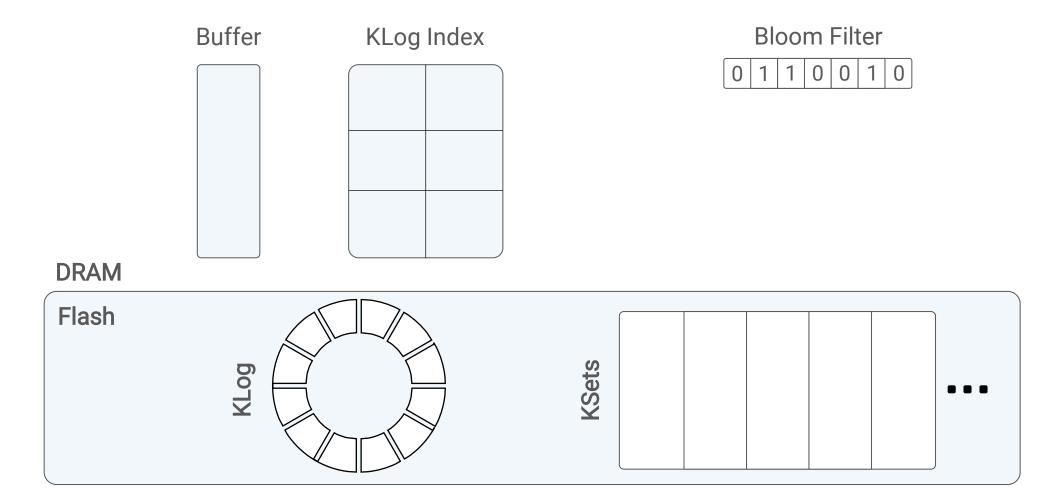
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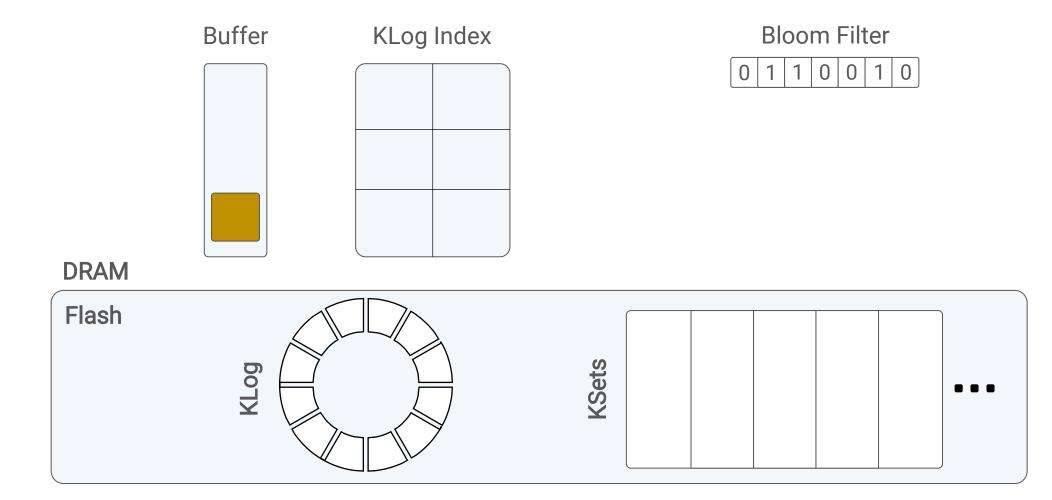


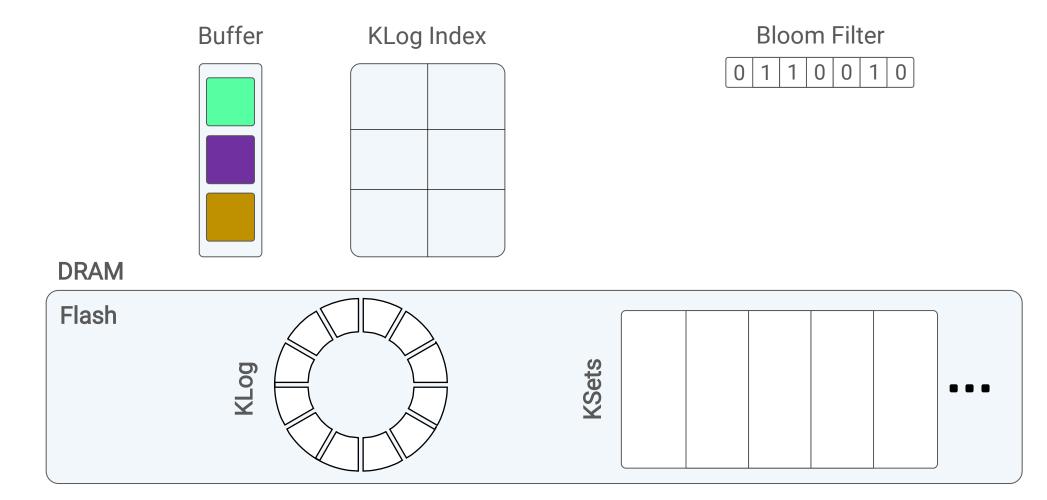


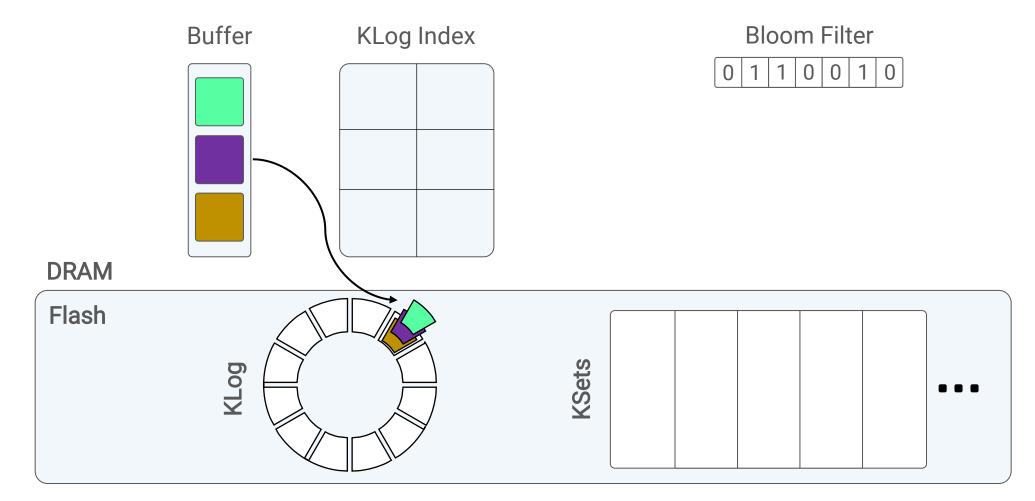
• Have advantage of both the design

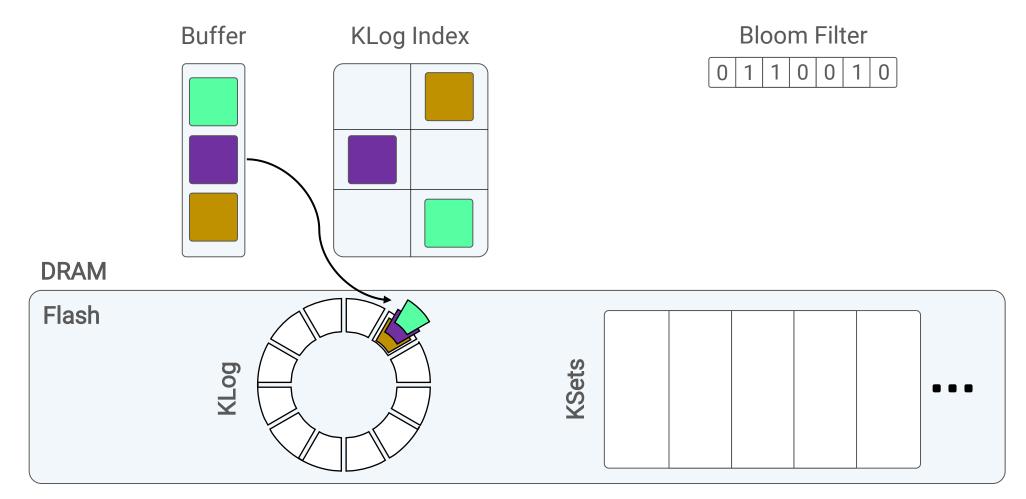


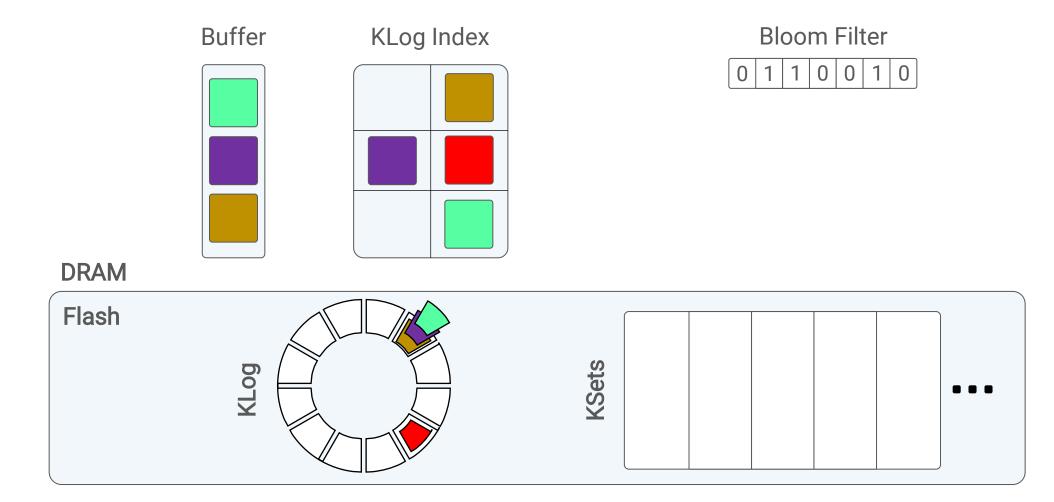
• Have advantage of both the design: 5% to KLog

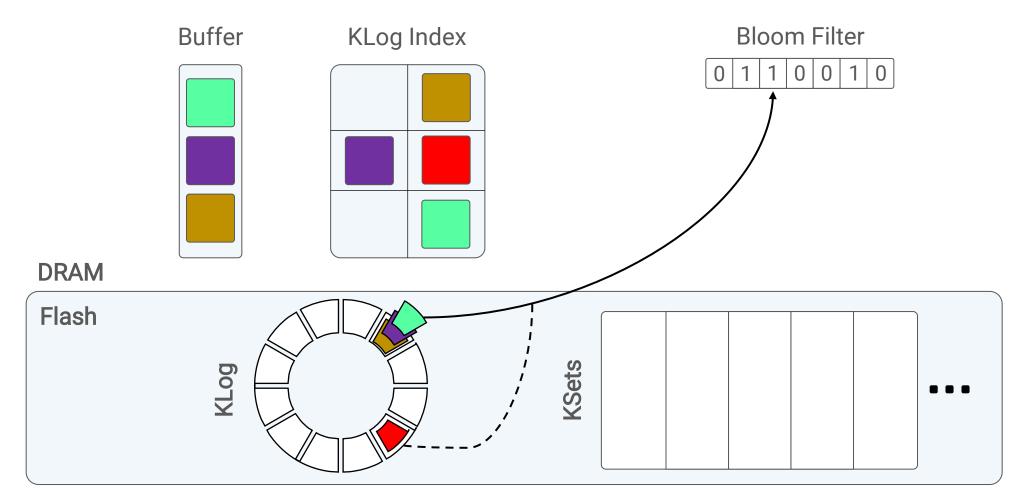


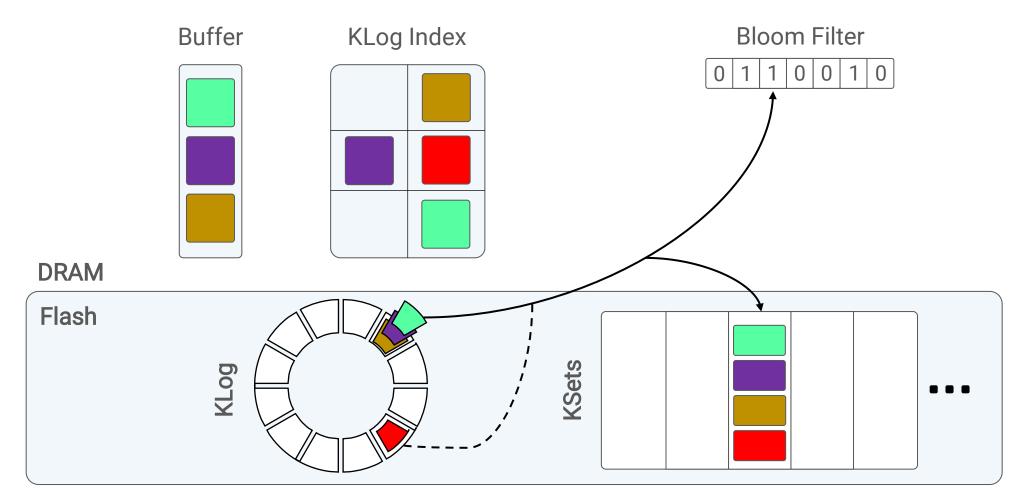


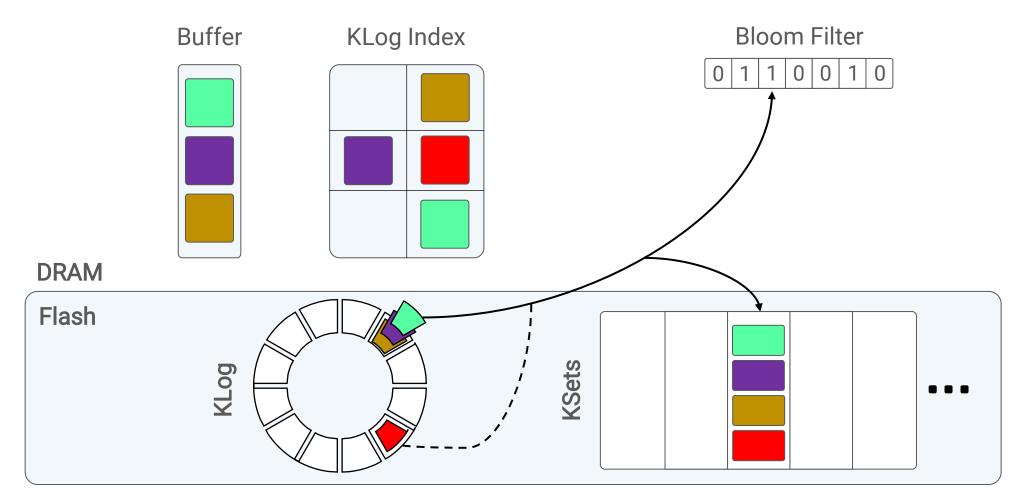












• Admission policies to reduce flash write

Kangaroo: Evaluation

Kangaroo: Evaluation (Miss Ratio)

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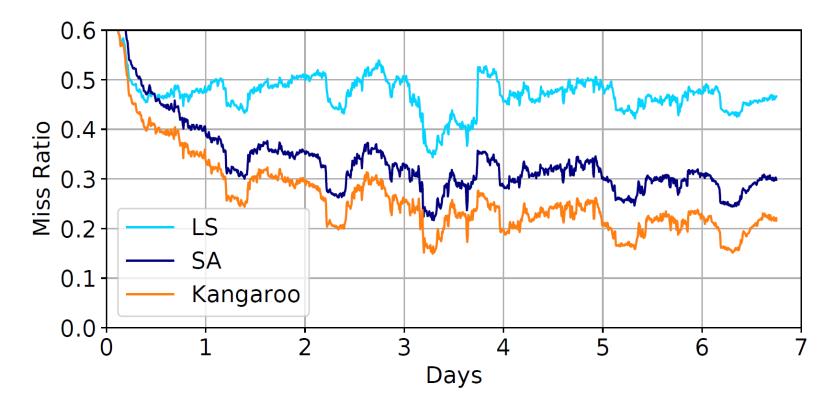


Fig: Miss ratio for all three systems over a 7-day Facebook trace. All systems are run with 16 GB DRAM, a 1.9 TB drive, and with write rates less than 62.5 MB/s

Kangaroo: Evaluation (Miss Ratio)

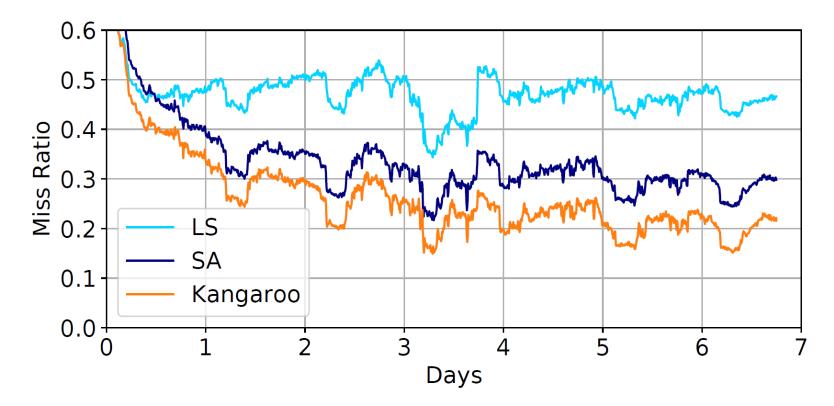
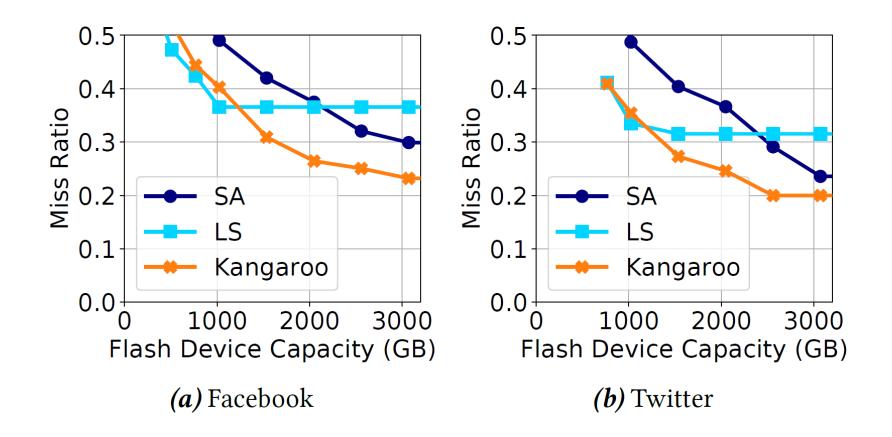
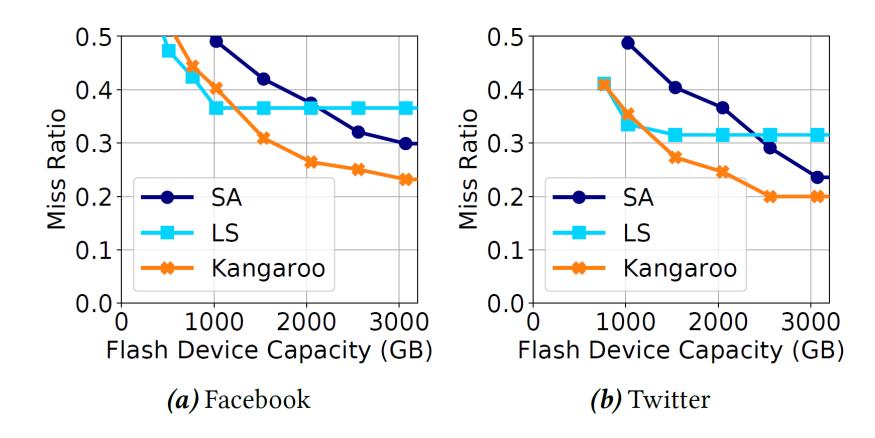


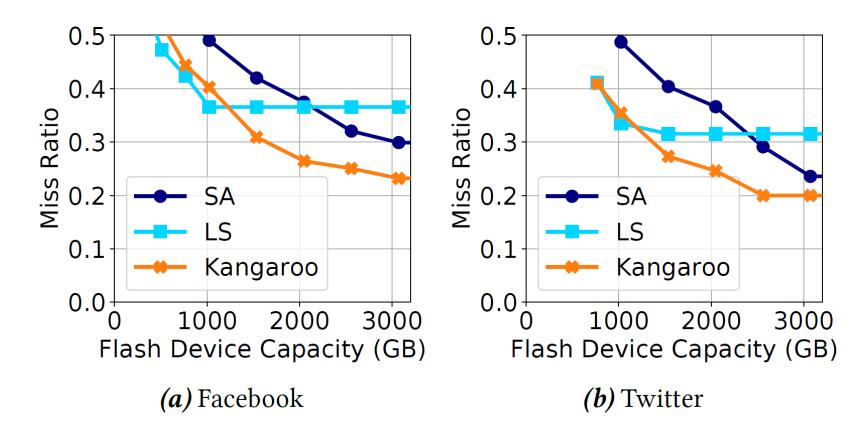
Fig: Miss ratio for all three systems over a 7-day Facebook trace. All systems are run with 16 GB DRAM, a 1.9 TB drive, and with write rates less than 62.5 MB/s

• Kangaroo reduces cache misses by 29% vs. SA and by 56% vs. LS



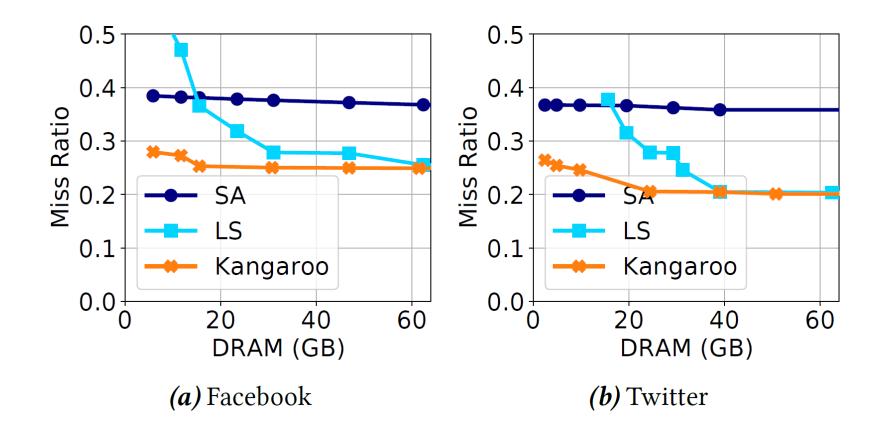


LS saturates faster: Fixed DRAM size

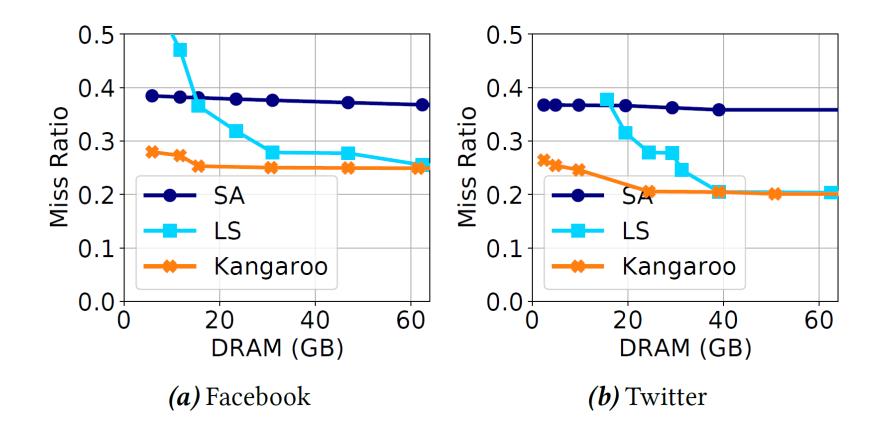


- LS saturates faster: Fixed DRAM size
- SA performs worst than Kangaroo: FIFO eviction and higher write amplification

Kangaroo: Evaluation (DRAM Size)

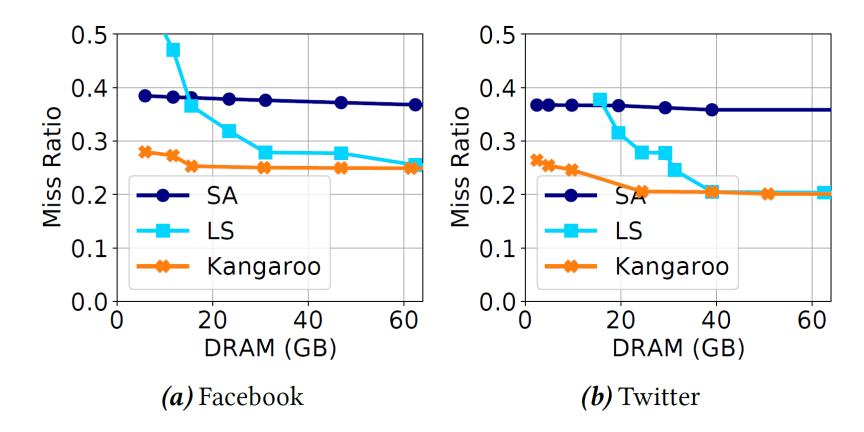


Kangaroo: Evaluation (DRAM Size)



• LS improves miss ratio: uses available DRAM

Kangaroo: Evaluation (DRAM Size)



- LS improves miss ratio: uses available DRAM
- o SA has no effect, however Kangaroo perform better than all

Conclusion

 $\circ~$ Have advantages of both the designs, i.e. LS and SA

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- Better miss rate than both the designs

Conclusion

- $\circ~$ Have advantages of both the designs, i.e. LS and SA
- Better miss rate than both the designs
- Throughput and latency are not better than both designs but very well suited in production

Questions ?

Thank You