

Incompressibility \rightarrow ML Test

Fix a prefix-free machine M . Consider

$$\text{Comp}_b = \{ x \in \Sigma^* \mid K(x) \leq |x| - b \}$$

and $R_b = \bigcup_{x \in \text{Comp}_b} [x]$

i.e. set of all infinite sequences with some b -compressible prefix.

(1) R_b is c.e. given b : We can dovetail over all programs and accept x as a member of Comp_b when some program of length $\leq |x| - b$ halts & outputs x . (Otherwise, we run forever, but that is fine since we need only c.e.)

(2) R_1, R_2, \dots are uniformly ce. in the index - easy to see since the above conversion from b to the program accepting Comp_b can be done by a program.

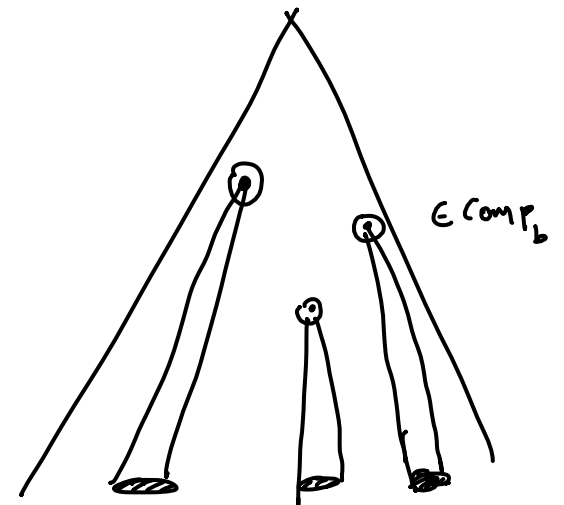
(3) Claim. $\mu(R_b) \leq \frac{1}{2}b.$

Proof

Let V_b be the largest prefix-free subset of Comp_b .

Then
$$\sum_{x \in V_b} \frac{1}{2^{|x|}} = \mu(R_b).$$

Let $x \in V_b$. Pick the shortest program p_x for x . Then $|p_x| \leq |x| - b$, since $x \in V_b \subseteq \text{Comp}_b$. Then $2^{-|p_x|} \geq 2^{+b} \cdot 2^{-|x|}$. Hence,



$$1 \geq \sum_{x \in V_b} \overset{\text{Kraft inequality}}{\rightarrow}$$

$$\frac{1}{2^{|P_x|}} \geq 2^b \left(\sum_{x \in V_b} \frac{1}{2^{|x|}} \right) = 2^b \cdot \mu(R_b)$$

$|P_x| \leq |x| - b$

Thus $\mu(R_b) \leq 2^{-b}$

Hence, by (1), (2), (3), it follows that R_1, R_2, \dots is a Martin-Löf test.

