

ASSIGNMENT 3

POINTS: 50

DATE GIVEN: 07-MAR-2019

DUE: 28-MAR-2019

Rules:

- You are strongly encouraged to work *independently*. That is the best way to understand & master the subject.
- Write the solutions on your own and honorably *acknowledge* the sources if any. <http://cse.iitk.ac.in/pages/AntiCheatingPolicy.html>
- Submit your solutions, before time, to your TA. Preferably, give the TA a printed copy of your LaTeXed or Word processed solution sheet.
- Clearly express the fundamental *idea* of your proof/ algorithm before going into the other proof details. The distribution of partial marks is according to the proof steps.
- There will be a penalty if you write unnecessary or unrelated details in your solution. Also, do not repeat the proofs done in the class.
- Problems marked '0 points' are for practice.

**Question 1:** [3+4 points] For an  $m \times n$  matrix  $A$ , over a field, one can define the *rank* via columns or via rows. Prove that  $\text{col-rk}(A) = \text{row-rk}(A)$ . Thus, we could talk about *the*  $\text{rk}(A)$ .

Let  $B$  be another  $m \times n$  matrix. Prove the *subadditivity* property:  $\text{rk}(A + B) \leq \text{rk}(A) + \text{rk}(B)$ .

**Question 2:** [10 points] Let  $X$  be a nonnegative real-valued random variable with mean  $\mu$ . Let  $\delta \in (0, 1)$ . We are interested in estimating the probability of deviation of  $X$  from  $\mu$  (multiplicatively). Show that:

- (1)  $\text{Prob}[X > (1 + \delta)\mu] < \left(\frac{e^\delta}{(1+\delta)^{1+\delta}}\right)^\mu$ , and
- (2)  $\text{Prob}[X < (1 - \delta)\mu] < \left(\frac{e^{-\delta}}{(1-\delta)^{1-\delta}}\right)^\mu$ .

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**Question 3:** [6 points] Let  $n = d^2$  and consider the  $n$ -variate permanent polynomial,  $\text{per}_d$ , over  $\mathbb{F}_q$ . Prove that  $\text{Prob}_{\alpha \in \mathbb{F}_q^n}[\text{per}_d(\alpha) \neq 0] \geq \frac{1}{4}$ .

**Question 4:** [9+4+4 points] We have extensively used the binomial estimates. They originate from Stirling's approximation. Prove that for all  $n \in \mathbb{N}$ ,

$$\sqrt{2\pi} \leq \frac{n!}{n^{n+0.5}e^{-n}} \leq e.$$

Using this prove that:

- (1) For  $f + g = o(h)$ ,  $\ln \frac{(h+f)!}{(h-g)!} = (f + g) \ln h \pm O\left(\frac{(f+g)^2}{h}\right)$ ,
- (2) For constants  $\alpha \geq \beta > 0$ ,  $\ln \binom{\alpha n}{\beta n} = \alpha n \cdot H(\beta/\alpha) - O(\ln n)$ .

**Question 5:** [10 points] Prove that the homogeneous- $\Sigma\Pi\Sigma$  complexity of  $\det_d$  is  $2^{\Omega(d)}$  (over any field).

**Question 6:** [0 points] What is the formula complexity of determinant?

**Question 7:** [0 points] Show that width-two ABP is an incomplete model.

**Question 8:** [0 points] Over constant-size  $\mathbb{F}_q$ , show that the  $d^2$ -variate polynomial  $\text{sym}_{\leq d}$  requires depth-3 circuits of size  $d^{\Omega_q(d)}$ .

**Question 9:** [0 points] Complete the missing details in the exponential lower bound proof we did in the class; to express determinant as a multilinear depth-3 circuit.

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