

- Now we need to study the effect of a random partitioning on a t-product.

Lemma: Let  $f(x)$  be n-variate & computable by a size- $s$  multilinear depth- $\Delta$  formula.

If  $X = Y \sqcup Z$ ,  $|Y| = |Z| = n/2$ , is random then with probability  $1 - s \cdot \exp(-n^{\Omega(1/\Delta)})$ :

$$\Gamma_{Y,Z}^t(f) = s \cdot 2^{n/2} \cdot \exp(-n^{\Omega(1/\Delta)}).$$

Proof:

- By the previous lemma, write  $f = g_0 + \sum_{i=1}^s g_i$  where  $\deg g_0 \leq n/100$  &  $g_1, \dots, g_s$  are multilinear t-products.

- Note that  $g_0$ 's sparsity can be at most

$$\sum_{i \leq n/100} \binom{n}{i} = 2^{H_2(1/100) \cdot n - O(\log n)} < 2^{n/10}.$$

$$\Rightarrow \Gamma_{Y,Z}^t(g_0) < 2^{n/10} \text{ (sub-additivity).}$$

- All that remains is to bound  $\Gamma_{Y,Z}^t(g_1)$  for a random partition  $X = Y \sqcup Z$ .

- Let  $g = h_1 \cdots h_t$ ,  $h_i \in \mathbb{F}[x_i]$ , be a  $T$ -product for  $X = \bigsqcup X_i$ .

Let  $y_i := x_i \cap Y$  &  $z_i := x_i \cap Z$ .

- Let  $d_i := |\#y_i - \#z_i|/2$  be the imbalance between  $y_i, z_i$  in  $x_i$ .

$x_i$  is called  $k$ -imbalanced if  $d_i \geq k$ .

Let  $b_i := (\#y_i + \#z_i)/2 = \#x_i/2$ .

$$\begin{aligned} \text{We have } T_{Y,Z}(g) &= \prod_i T_{Y,Z}(h_i) \leq \prod_i 2^{\min(|Y_i|, |Z_i|)} \\ &= \prod_i 2^{b_i - d_i} = 2^{|X|/2} / \prod_i 2^{d_i}. \end{aligned}$$

$\Rightarrow$  it suffices to show that one of the  $x_i$ 's is imbalanced (i.e.  $d_i$  is large).

- We need to estimate  $|Y_i|$  on choosing a random  $Y \in \binom{[n]}{n/2}$ .

- The relevant probability is that of the hypergeometric distribution.

Claim: For a fixed set  $A \in \binom{[n]}{a}$ ,  $k \leq a \leq 2n/3$ ,

$$\Pr_{R \in \binom{[n]}{n/2}} [|R \cap A| = k] = O(1/\sqrt{a}).$$

Proof:

- $\Pr_R [|R \cap A| = k] = \binom{a}{k} \cdot \binom{n-a}{n/2-k} / \binom{n}{n/2}$
- Call it  $P(k)$ .
- $P(k+1) > P(k)$  iff  
 $(a-k)(\frac{n}{2}-k) > (k+1)(\frac{n}{2}-a+k+1)$  iff  
 $\frac{an}{2} - k(a+\frac{n}{2}) > (\frac{n}{2}-a+1) + k(\frac{n}{2}-a+2)$
- iff  $k < \frac{a-1}{2}$ .
- Thus,  $P(k) \leq \binom{a}{\frac{a-1}{2}} \cdot \binom{n-a}{\frac{n}{2}-\frac{a-1}{2}} / \binom{n}{n/2}$   
 $(\text{Stirling's approx.})$ 
 $= O\left(\sqrt{\frac{n}{a(n-a)}}\right) = O\left(\frac{1}{\sqrt{a}}\right).$  □

( $R$ -balanced)

- Let  $\Sigma_i$  denote the event that  $d_i < k$ .
- We have  $\Pr[\bigwedge_{i=1}^t \Sigma_i]$  equal to  
 $\Pr[\Sigma_1] \cdot \Pr[\Sigma_2 | \Sigma_1] \cdot \Pr[\Sigma_3 | \Sigma_1 \wedge \Sigma_2] \dots$ .
- $\Pr[\Sigma_1] = \Pr[Y \cap X_1 \in [b_1 - k, b_1 + k]]$

which the above claim estimates as:

$$k \cdot O(1/\sqrt{b_i}) \quad (\text{assuming } k \leq b_i/2).$$

- Consider the event  $\Sigma_i$  given  $\Sigma_1, \dots, \Sigma_{i-1}$ . Since  $x_1, \dots, x_{i-1}$  have been partitioned in a fairly balanced way ( $\forall j \in [i-1], d_j < k$ ), we deduce that  $|Y \cap (x_1 \cup \dots \cup x_{i-1})^c|$
- $$\begin{aligned} &= |Y \cap x| - |Y \cap (x_1 \cup \dots \cup x_{i-1})| \\ &< n/2 - (b_1 - k + \dots + b_{i-1} - k) \\ &= (n/2 - b_1 - \dots - b_{i-1}) + (i-1)k \end{aligned}$$
- $\Rightarrow$  The partition of  $x' := (x_1 \cup \dots \cup x_{i-1})^c$  by  $Y \cup Z$  is  $(ik)$ -balanced.

- So, assuming  $ik \ll n$ , we can redo the calculation in the above claim & still get  $\Pr[\Sigma_i | \Sigma_1 \wedge \dots \wedge \Sigma_{i-1}] = k \cdot O(1/\sqrt{b_i})$ .
- $$\Rightarrow \Pr\left[\bigwedge_{i \in [t]} \Sigma_i\right] = O(k^t / \sqrt{b_1 \dots b_t})$$

$$\Rightarrow \Pr_y[T_{Y,Z}(g) > 2^{|x|/2} \cdot 2^{-k}] = O(k^t / \sqrt{b_1 \dots b_t})$$

- In particular, on fixing  $k \leq t^{1/3}$ , we get:  $\Pr_y [T_{y,z}(g) > 2^{n/2} \cdot 2^{-k}] = O\left(\prod_{i=1}^t \frac{1}{i} \cdot t^{-1/6}\right)$   
 $= O(t^{t/6}) = \exp(-n^{o(1/\Delta)})$ .

$\Omega$ : independent  
of  $\Sigma \rightarrow$   
if  $\Sigma \leq 1/3$

$$\Rightarrow \Pr_y [T_{y,z}(f) > \delta \cdot 2^{n/2-k} =: \delta 2^{n/2} \cdot 2^{-t^\varepsilon}]$$
 $= \delta \cdot \exp(-n^{o(1/\Delta)}).$

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- Thus, there is an  $\varepsilon \in (0, \frac{1}{3}]$  such that if  $\delta \leq \exp(-t^\varepsilon)$  then  $\Pr_y [T_{y,z}(f) = 2^{n/2}] < 1/10$ .

$\Rightarrow f(x)$  could compute  $\det(x)$  only if  $\delta > 2^{t^\varepsilon} = \exp(n^{o(1/\Delta)})$ .

- This finishes (Raz, Yehudayoff '09) proof for  $\text{det}_n$  or  $\text{per}_n$  against constant-depth multilinear model.

- We can also say something for multilinear formulas using the probability calculation seen above.

- The multilinear products of interest there are:

Defn: Multilinear  $f = \prod_{i=1}^t g_i$ , with partition  $X = \bigsqcup_{i \in [t]} X_i$ , is called a log-product if for all  $i$ ,  $|X|/3^i \leq |X_i| \leq 2 \cdot |X|/3^i$  and  $|X_t| = 1$ .

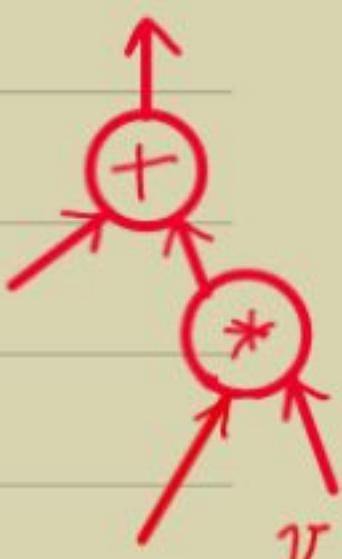
Lemma: Any size- $\delta$  multilinear formula  $\phi$  can be written as a sum of  $(\delta+1)$  log-products.

Proof: • Let  $|X| > 2$  &  $\phi$  compute  $f$ .  
• Let  $v$  be a node in  $\phi$  that

depends on variables  $X_v$  such that

$\xrightarrow{\text{fan in } 2}$   $|X|/3 \leq |X_v| \leq 2 \cdot |X|/3$

• By the formula properties, we have



$$f = \phi_v \cdot g + \phi_{v=0}$$

for some  $g \in F[X \setminus X_v]$ .

- Note that  $|X|/3 \leq |X \setminus X_v| \leq 2 \cdot |X|/3$ .
- Moreover, since  $g$  has size  $\leq d$ , we can use induction & write it as a sum of  $\leq \text{size}(g)+1$  log-products.

Similarly, for  $\phi_{v=0}$ .

$\Rightarrow f$  is a sum of  $(s+1)$  log-products.

□

- Now, we can estimate  $T_{Y,Z}(h_1 \cdots h_t)$  for a log-product  $h_1 \cdots h_t$ ,  $t = O(\lg n)$ .

Note that around  $\frac{1}{2} \lg n$  many of these  $h_i$ 's do depend on at least  $\sqrt{n}$  many variables each.

$\Rightarrow$  On doing the probability calculation we will get that  $T_{Y,Z}(h_1 \cdots h_t)$  is high with prob. smaller than  $n^{-\Omega(\lg n)}$ .

(Raz '09)  $\Rightarrow \det_n$  or  $\text{per}_n$  requires  $n^{\Omega(\lg n)}$  size!