

## Computing the Zeta function

- The tools we developed till now are not yet enough to actually compute  $Z(t, C)$ .
- The key idea, due to Weil, that we now study is : The Frobenius map  $F$  acts linearly on the Jacobian variety  $J(C)$ , or  $C_{\bar{\mathbb{F}}_p}(C)$ , and its characteristic polynomial is  $Z(t, C)$  !
- To get an algorithm out of this we need the addition algorithm on  $J(C)$ . This we sketch only for the curve  $C$ :  
 $y^2 = f(x)$ , called a hyperelliptic curve.
- The details we skip could be found in standard articles.

## Frobenius interpretation of $Z(t, c)$

- Recall that  $Z(t, c) = \frac{L(t)}{(1-t)(1-qt)}$ , where,  $\deg L = 2g$ ,  $L(0) = 1$ , and  $L$  is an integral polynomial.
- Recall that  $Z(t, c) = \exp\left(\sum_{m \geq 1} \frac{N_m(c)}{m} \cdot t^m\right)$ .
- We want to interpret  $N_m$  as fixed pts. of the Frobenius map  $F: \bar{C} \rightarrow \bar{C}$ , where  $\bar{C}$  is the curve  $C$  over  $\bar{k} = \overline{\mathbb{F}_q}$ .
  - In this case  $F$  maps an actual point  $P = (a_i)$  to  $F(P) = (a_i^{q^i})$ .
- So,  $N_m(c) = N_m(\bar{C})$  can be viewed as the  $\#\{\text{fixed pts. of } F^m \text{ on } \bar{C}\}$ .
- Now let us consider the action of  $F$  on  $C\ell(\bar{C})$ .

- We denote the  $m$ -th linear map as  $(F^m; J(\mathbb{C}))$ .
- Since  $J$  is finite, we have for any prime  $\ell$ ,  
 the  $\ell$ -torsion points  $J[\ell] \subset J$  are finite.  
 $\rightarrow$  (I.e.  $D \in J$  s.t.  $\ell \cdot D = 0$ .)
- Thus,  $J[\ell]$  is a finite dimensional vector space over  $\mathbb{F}_\ell$ .
- It can be shown that, for prime  $\ell \neq p$ ,  $F$  is a  $\mathbb{F}_\ell$ -linear automorphism of  $J[\ell]$ .
- Using these ideas, Weil gave a second proof of the RHT via  $\ell$ -adic means!
- In particular, he showed that:  
 $L(t, c) \equiv \det(I - Ft; J[\ell]) \pmod{\ell}$ .
- ▷ I.e.  $L(t)$  is the characteristic polynomial of the Frobenius map on  $J[\ell]$ .

## Sketch:

- Weil gave three  $\mathbb{Q}_\ell$ -vector spaces  $H^i(C)$ , ( $i=0, 1, 2$ )  
 s.t.  $N_m(C) = \#\{\text{fixed pts. of } F^m \text{ on } \bar{C}\}$   
 $= \sum_{i=0}^2 (-1)^i \cdot \text{Tr}(F^m; H^i(C))$

$$\Rightarrow Z(t, C) = \exp \left( \sum_{\substack{0 \leq i \leq 2 \\ m \geq 1}} (-1)^i \cdot \text{Tr}(F^m; H^i(C)) \frac{t^m}{m} \right)$$

$$= \prod_{i=0}^2 \left\{ \exp \left( \sum_{m \geq 1} \text{Tr}(F^m; H^i(C)) \frac{t^m}{m} \right) \right\}^{(-1)^i}$$

- This can be simplified!
- Claim: If  $\varphi$  is an endomorphism of a finite dim. vector space  $V$  over  $\mathbb{F}$ , then  
 $\exp \left( \sum_{m \geq 1} \text{Tr}(\varphi^m; V) \cdot \frac{t^m}{m} \right) = (\det(1 - \varphi t; V))^{-1}$ .

Pf: • Let  $\dim V = 1$ . Then,  $\exists \lambda \in \mathbb{F}$  s.t.

$$\varphi: v \mapsto \lambda v, \quad \forall v \in V.$$

$$\bullet \text{ In this case, LHS} = \exp \left( \sum_{m \geq 1} \lambda^m \frac{t^m}{m} \right)$$

$$= \exp(-\log(1 - \lambda t)) = (1 - \lambda t)^{-1} = \text{RHS.}$$

• If  $\dim V > 1$ , do induction on  $\dim V$ ; by picking a

$\varphi$ -invariant subspace  $V' \subset V$ .

• So, LHS =  $\exp\left(\sum_{m \geq 0} \text{Tr}(\varphi^m; V') \frac{t^m}{m}\right)$ .

$$\exp\left(\sum_{m \geq 0} \text{Tr}(\varphi^m; V/V') \frac{t^m}{m}\right)$$

$$\begin{aligned} &= \det(1 - \varphi t; V')^{-1} \cdot \det(1 - \varphi t; V/V')^{-1} \\ &= \det(1 - \varphi t; V)^{-1} = \text{RHS}. \quad \square \end{aligned}$$

- Thus,  $Z(t, c) = \frac{\det(1 - F_t; H^1(c))}{\det(1 - F_t; H^0(c)) \cdot \det(1 - F_t; H^2(c))}$

- It turns out that:

$$L(t, c) = \det(1 - F_t; H^1(c)),$$

the other two give  $(1-t) \cdot (1 - q^t)$ , &  
 $H^1(c) \pmod{\ell}$  is exactly  $\mathbb{J}[c]$ .

- This is called the cohomological interpretation of the zeta fn.

- Weil conjectured this pattern also for the general varieties!

- Since the coefficients of  $L(t)$ , by RH, are of magnitude  $\leq \binom{2g}{g} (\sqrt{q})^{2g}$ ; it will suffice in Chinese remainderring to compute  $L(t) \pmod{\ell}$  for the first  $\log(2\binom{2g}{g} q^g) = (\log q)^{O(g)}$  many primes.

- Adleman-Huang (2001) gave an algorithm to do this in time  $(\log q)^{O(g^2 \log g)}$ , assuming that:  
 $C$  is a hyperelliptic curve, i.e.  
the fn. field  $K = \mathbb{F}_q(x)[y] / \langle y^2 - f(x) \rangle$ .
- For general curves  $C$  we may not even be able to compute within  $J(C)[e]$ .