

- Obviously, the #unknowns  $m$  is still  $< 3d^3$ .
- If it has no solution, then the corresponding  $m \times m$  matrix  $M$  (with entries as coefficients of  $g_k$ ) has a nonzero determinant  $D(\bar{y})$ .
 
$$\Rightarrow \deg D(\bar{y}) < m \cdot 2^k \leq 3d^3 \cdot 2d^2 = 6d^5.$$

$$\Rightarrow \Pr_{\bar{a} \in S^n} [D(\bar{a}) = 0] \leq 6d^5 / |S|.$$

- On the other hand, by the hypothesis, the system has a solution for "many"  $\bar{y} = \bar{a} \in S^n$ , in which cases  $D(\bar{a}) = 0$ .

This contradiction implies  $D(\bar{y}) = 0$ .

$\Rightarrow g(x, t, \bar{y}) \& l_k$  do exist!

- The  $\sum_i \deg_{y_i} g \leq \deg |M|$  follows from the Cramer's rule of solving linear system of equations.  $\square$

- Finally, we want to use  $g(x, t, \bar{y})$  to factor  $f(x, \bar{y}t)$ .

- actually,  
 $t=1$  suffices  
here  $\rightarrow$
- Consider  $r(t, \bar{y}) := \gcd_x(f(x, \bar{y}t), g(x, t, \bar{y}))$ .  
 $\Rightarrow \deg r \leq d \cdot (d+d+6d^5) < 7d^6 [\because d \geq 2]$   
[However,  $\deg_t r \leq d \cdot d = d^2 < 2^k$ .]
  - On the other hand, we know from "bivariate factoring" proof & the construction of  $g$  that  $r(t, \bar{a}) = 0$ , for a "good" fraction of  $\bar{a} \in S^n$ .  
 $\Rightarrow r(t, \bar{y}) = 0$ .  
 $\Rightarrow \gcd_x(f(x, \bar{y}t), g(x, t, \bar{y})) \neq 1$ .  
 $\Rightarrow f$  is reducible.
  - This proves HIT !  $\square$

### Blackbox factoring algorithm

Oracle to

Input:  $f(x, \bar{y}) \in \mathbb{F}[x, \bar{y}]$  of  $\deg d$  &  $S \subseteq \mathbb{F}$  s.t.  $|S| > 7d^7$ .  
 $f$  is almost-monic in  $x$  &  $\partial_x f \neq 0$ .

Output: Blackboxes to the factors of  $f$ .

Algo: 1) We compute the number of factors by :  
1.1) Pick  $\bar{a}, \bar{b} \in S^n$  randomly.

1.2) Factor  $f_{\bar{a}, \bar{b}}(x, t) := f(x, \bar{a}t + \bar{b})$ .

Let  $\{\tilde{f}_i(x, t) \mid i \in [\ell]\}$  be the irreducible factors.

[ $\triangleright$  Why  $\ell$  is the number of factors of  $f(x, \bar{y})$ .

Pf: Let  $f_i(x, \bar{y})$ ,  $i \in [\ell']$ , be the actual factors.

These are all almost-monic irreducibles.

By H.I.T.:  $f_i(x, \bar{a}t + \bar{b})$  is reducible with prob.  
 $< 7d^6/|S|$ .

$\Rightarrow \Pr[\exists i, f_i(x, \bar{a}t + \bar{b}) \text{ reduces}] < 7d^7/|S|$ .  $\square$

2) Assuming that  $\tilde{f}_i(x, t)$  is the projection of an actual factor, i.e.  $\tilde{f}_i = f_i(x, \bar{a}t + \bar{b})$ , we want to compute the value  $f_i(\alpha, \bar{\beta})$  for any given  $(\alpha, \bar{\beta}) \in \mathbb{F}^{n+1}$ .

For this we define a trivariate that "contains" both the projections of  $f$  to the line  $\bar{a}t + \bar{b}$  & the point  $(\alpha, \bar{\beta})$ :

$$g(x, t_1, t_2) := f(x, \bar{a}t_1 + \bar{b} + (\bar{\beta} - \bar{b})t_2).$$

$\triangleright g(x, t, 0) = f(x, \bar{a}t + \bar{b})$  &  $g(\alpha, 0, 1) = f(\alpha, \bar{\beta})$ .

3) Now we factor  $g$  to compute  $f_i(\alpha, \bar{\beta})$ :

3.1) Using 3-variate factoring, find the irreducible factors  $\{g_j(x, t_1, t_2) \mid j \in [e]\}$  whp.

3.2) Find the index  $j$  s.t.  $\tilde{f}_i(x, t) = g_j(x, t, 0)$ .

3.3) Output  $g_j(\alpha, 0, 1)$ .

[ Whp we will get the factors  $g_j$  that exactly are projections like  $f_i(x, \bar{\alpha}t_1 + \bar{t}_1 + (\bar{\beta} - \bar{t}_1)t_2)$ . ]

Theorem (Kaltofen & Trager, 1990): Given  $f(x, \bar{y})$ , as a blackbox, one can factorize  $f$  (as blackboxes) in randomized  $\text{poly}(n, d)$  time (assuming that univariate factoring can be done).