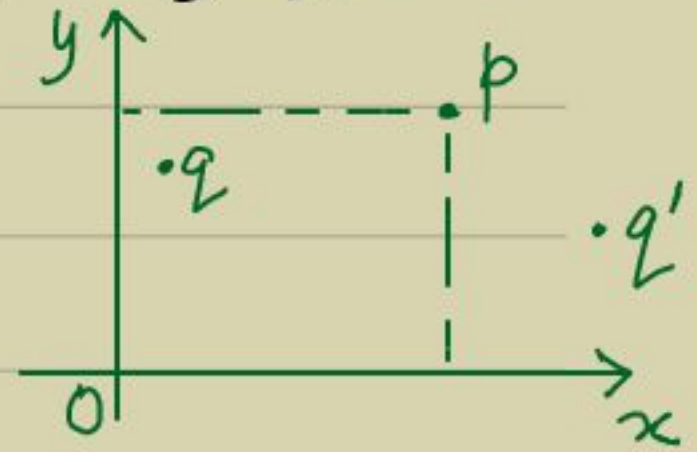


## Non-dominated points problem

- Defn: A point  $p$  dominates  $q$  if  $x(p) > x(q)$  &  $y(p) > y(q)$ .



Input: Given  $n$  points  $S \subset \mathbb{R}^2$ .

Output: Points  $p$  in  $S$  that are not dominated by any point in  $S$ .

- Brute-force algorithm:

Go over every  $p \in S$  & compare with each  $q \in S$ .

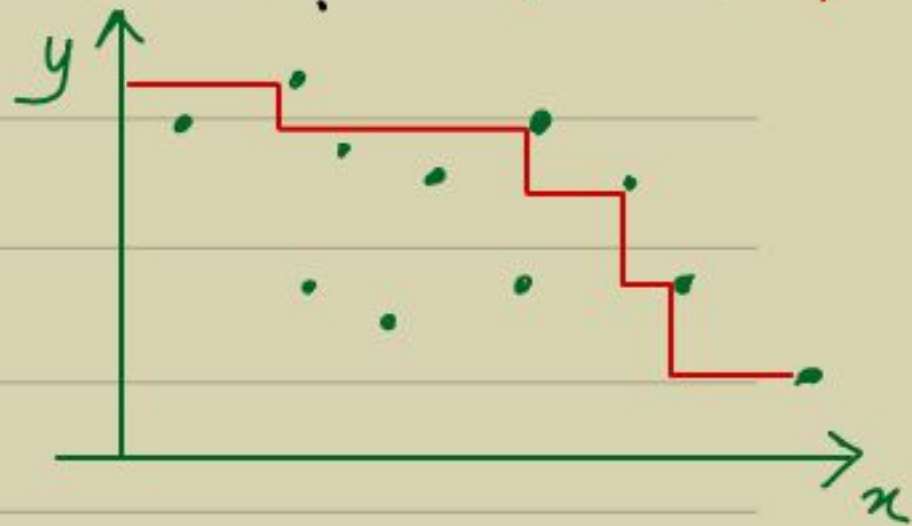
Takes  $O(|S|^2)$  time.

- Can geometry help?

- What is the structure of non-dominated points?

- They form a staircase! (Exercise)

- Thus, these are extremal points in a sense.



▷ A point with max.  $x$ -coordinate is a non-dominated point.

- Idea 1: Among the points in  $S$  with the max.  $x$ -coordinate, pick the point  $p$  with max.  $y$ -coordinate.

- Declare  $p$  non-dominated.
- Delete all the points  $q \in S$  with  $y(q) < y(p)$ .
- Repeat till  $S \neq \emptyset$ .

▷ If  $h = \#$  non-dominated points in  $S$ , then the time complexity is  $O(nh)$ .

output-sensitive algorithm

- Idea 2: Divide  $S$  into two halves. Solve each & Merge the two staircases.

- Let the two staircases be  $T_L$  &  $T_R$ .

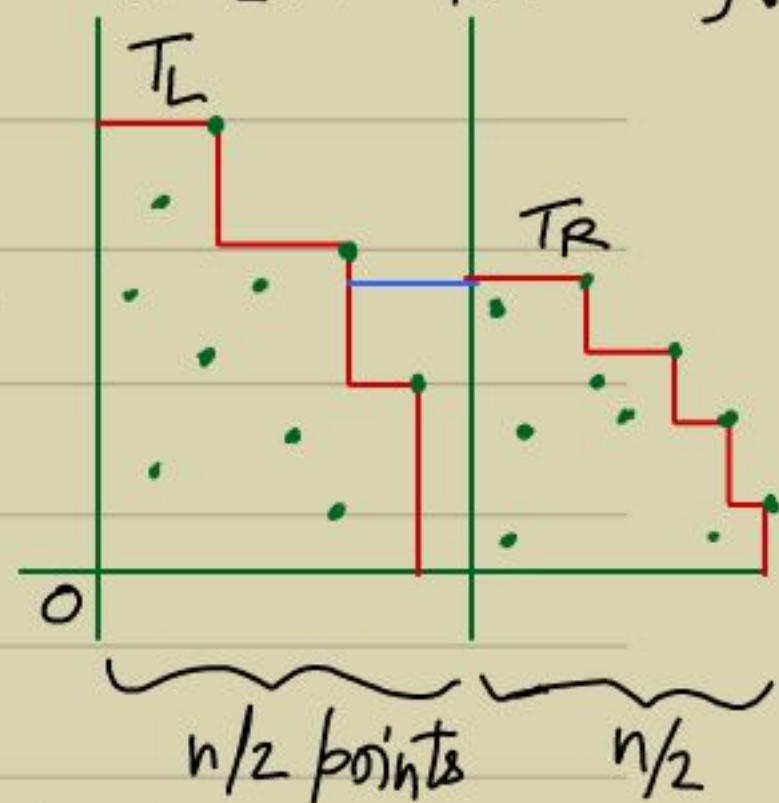
- During the merge step scan  $T_R$  to find the max- $y$ -coordinate point  $p \in T_R$ .

Scan  $T_L$  & delete those points  $q$  s.t.  $y(q) < y(p)$ .

Take the union of  $T_L$  &  $T_R$  that remain.

- Merge steps are  $O(|T_L| + |T_R|)$  many.

Theorem: Non-dominated points in  $S$  can be computed in  $O(|S| \lg |S|)$  time.



- by Kung, Luccio, Preparata (1975).