

Number Theory

- Properties of numbers. It's considered the most beautiful area in maths.
- Much of modern maths was developed to answer number theory questions.

Basics

- [Integer division] Lemma: Given numbers $a, b \in \mathbb{Z}$, we can uniquely find $a = qb + r$, s.t. $q, r \in \mathbb{Z}$ & $0 \leq r < b$ (Assume $b > 0$).
- Pf: - Existence: Follows from long division.

• Unique: $a = qb + r = q'b + r'$.
 $\Rightarrow b \mid (r - r')$ [i.e. b divides $r - r'$].
& $|r - r'| < b$.
 $\Rightarrow r - r' = 0 \Rightarrow$ unique r . \square

- Defn: • b divides a , if $r = 0$.
or $b \mid a$.

- Ex. $7 \nmid 101$, $7 \mid 105$.

• GCD of a, b , denoted $\gcd(a, b)$, is the largest $c \in \mathbb{N}$ s.t. $c \mid a$ & $c \mid b$.

• If $\gcd(a, b) = 1$ then we say a & b are coprime.

$$\begin{aligned}
 -1g. \quad (5, 11) &= (5, 5 \times 2 + 1) = (5, \underline{1}) = (5 \times 1 + 0, \underline{1}) = \underline{(0, 1)} \\
 (5, 12) &= (5, 5 \times 2 + 2) = (5, 2) = (2 \times 2 + 1, 2) \\
 &= (\underline{1}, 2) = (1, 2 \times 1 + 0) = (\underline{1}, 0). \\
 (6, 9) &= (6, 6 \times 1 + 3) = (6, 3) = (3 \times 2 + 0, 3) = (\underline{0}, 3) \\
 &= 3.
 \end{aligned}$$

$$\triangleright (a^{\geq}, b) = (qb + r, b) = (r^{\leq}, b). \quad [\text{Exercise}]$$

Euclid's gcd algorithm (~300 B.C.): gcd (a^{\geq}, b) :

- if $b = ? 0$ then OUTPUT a .
- if $b = ? 1$ " " " 1.
- Compute $a =: bq + r$. OUTPUT $\text{gcd}(b, r)$.

Exercise: Prove its correctness, by induction.

The underlying equations are:

$$\gcd(a, b) \rightarrow a =: q_1 \times b + r_1 .$$

$$\gcd(b, r_1) \rightarrow b =: q_2 \times r_1 + r_2 .$$

⋮

$$\gcd(r_{k-1}, r_k) \rightarrow r_{k-1} =: q_{k+1} \times r_k + \underline{0}$$

$$\gcd(r_k, 0) \rightarrow r_k .$$

Theorem (Bézout's identity): $\exists \alpha, \beta \in \mathbb{Z}$ s.t.
 $\gcd(a, b) = \alpha \cdot a + \beta \cdot b .$

Pf:

- $r_1 = a - q_1 \cdot b$
- $r_2 = b - q_2 \cdot r_1 = b - q_2 \cdot (a - q_1 \cdot b)$
 $= (-q_2) a + (1 + q_2 q_1) \cdot b$

(by induction
on k) ;

- $r_k = \alpha \cdot a + \beta \cdot b ; \alpha, \beta \in \mathbb{Z}.$ □

▷ [Extended Euclid gcd algorithm] α, β are (efficiently) computable.

Ex.: If a, b are n -bits, how large can k be? [$2n$?] [Construct an eg.]

$$\begin{aligned} \triangle (a, b) &= \alpha \cdot a + \beta \cdot b \\ \text{[Reduction]} &= (\alpha + \delta \cdot b) \cdot a + (\beta - \delta a) \cdot b \end{aligned}$$

Theorem (Unique Bézout's id.): Let $a > b > 0$,

$$\exists \text{ unique } \alpha, \beta : (a, b) = \alpha \cdot a + \beta \cdot b,$$

$$\text{where } 0 \leq \alpha < b' \quad \& \quad \underline{a'} := \frac{a}{(a, b)}, \quad \underline{b'} := \frac{b}{(a, b)}.$$

$$-a' < \beta \leq 0$$

Pf: • a', b' are the coprime parts.

$$\bullet \text{ By Euclid, } \exists \alpha, \beta : \alpha a' + \beta b' = 1$$

— w.l.o.g. assume $0 \leq \alpha < b'$

[By Reduction]

$$\Rightarrow |\beta b'| = |1 - \alpha a'| < |b' a'| \Rightarrow |\beta| < a'.$$

• [Uniqueness] Suppose $(a, b) = \alpha_1 a + \beta_1 b = \alpha_2 a + \beta_2 b$
 $\Rightarrow (\alpha_1 - \alpha_2) a' = (\beta_2 - \beta_1) b' \Rightarrow b' | (\alpha_1 - \alpha_2) a'$.

• We've $|\alpha_1 - \alpha_2| < b'$ $\stackrel{?}{\Rightarrow} \alpha_1 - \alpha_2 = 0 = \beta_1 - \beta_2$

(cancel)

Claim 1: $(m, n) = 1$ & $m | n\alpha \Rightarrow m | \alpha$.

Pf: • Bézout's $\Rightarrow u \cdot m + v \cdot n = 1$

$$\Rightarrow u \cdot m \cdot \alpha + v \cdot n \cdot \alpha = \alpha$$

$$\Rightarrow m | \alpha. \quad \square$$

$\Rightarrow (\alpha_1, \beta_1) = (\alpha_2, \beta_2) \Rightarrow (\alpha, \beta)$ is unique.

\square

- Fundamental theorem of arithmetic:

Thm (Unique factorization) $\forall n \in \mathbb{N}$,
 \exists $2 \leq p_1 < \dots < p_k$ primes: $n = \prod_{i=1}^k p_i^{e_i}$ in a
unique way.

Pf: • [Existence]: Either n is prime or n has
largest prime $\leq n/2$. \Rightarrow Gives $n = \prod_i p_i^{e_i}$.

• [Uniqueness] $n = \prod_{i=1}^k p_i^{e_i} = \prod_{j=1}^l q_j^{f_j}$.

\Rightarrow (Cancellation) $p_1 = q_1$ & $e_1 = f_1$,
....., $p_l = q_l$ & $e_l = f_l$.

\Rightarrow uniqueness of factors. \square

- We can study numbers via primes (& prime-powers).

Modular Arithmetic

- Ex. Sunday is the first day. What is the 184th day? **Mod 7?**
- Last digit of $2^{2^{100}}$? **Mod 10?**
 - ↳ Impossible to store/compute
 - ↳ But we can work with $\{0, 1, \dots, 9\}$.
- Defn: $r_d(n) :=$ remainder of n/d .
 - Ex. $r_{10}(n)$, $n := 2^{2^{100}}$.

$$\triangleright r_d(ab) = r_d(r_d(a) \cdot r_d(b))$$

Pf:

$$\begin{aligned} \cdot a &= d \cdot q_1 + r_d(a) \\ \cdot b &= d \cdot q_2 + r_d(b) \\ \cdot ab &= d \cdot q_3 + r_d(ab) \end{aligned} \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} \Rightarrow ab = d \cdot (-) + \underbrace{r_d(a) \cdot r_d(b)}$$

& $r_d(\cdot) \in [0, \dots, d-1]$. \Rightarrow done! \square

Given $n: \forall a, b \in \mathbb{Z}$:

- Defn: $a \equiv b \pmod{n}$ if $n \mid (a-b)$.

\triangleright "mod n " is an equivalence relation of \mathbb{Z} .

- Ex. $a \equiv b \pmod{n} \stackrel{?}{\iff} a = b$. (If $|a-b| < n$.)

$$\begin{aligned} - \text{eg, } \pi_{10}(2^{64}) : & \quad 2^8 \equiv 6 \pmod{10} \\ & \Rightarrow 2^{8 \times 2} \equiv 36 \equiv 6 \pmod{10} \\ & \dots \Rightarrow 2^{64} \equiv 6 \pmod{10}. \\ & \dots \Rightarrow 2^{2^n} \equiv 6 \pmod{10}, \quad \forall n \geq 2. \end{aligned}$$

[The notation invented by Gauss, 1800s]

$$\triangleright a \equiv b \pmod{n} \Leftrightarrow \exists k : a = b + kn. \leftarrow \text{AP!}$$

- Defn: $\{b + kn \mid k \in \mathbb{Z}\}$ is the residue class of
 $b \pmod{n}$.

▷ $\forall n, \mathbb{Z}$ partitions into n residue classes $\pmod n$.

- Defn: • The residue classes are $0, 1, \dots, n-1$.
• We call it $\mathbb{Z}/n := \{0, 1, \dots, n-1\}$.

▷ $\forall a, b \in \mathbb{Z}/n$: (i) $a \pmod n + b \pmod n \equiv (a+b) \pmod n$
- eg. $1 + (n-1) \equiv n \equiv 0 \pmod n$.

(ii) $(a \pmod n) \cdot (b \pmod n) \equiv (ab \pmod n)$.

- eg. $2 \cdot (n-1) \equiv -2 \equiv n-2 \pmod n$.

- eg. $2^{39} \pmod{10} \equiv 2^{32+4+2+1} \equiv 6 \cdot 2^2 \cdot 2^1 \equiv 6 \cdot 8 \equiv 8$.

Qn: What's division $\pmod n$?

- Undefined, eg, $\frac{1}{2} \bmod 10 =: x \in \mathbb{Z}/10$.

$$\Rightarrow 1 \equiv 2x \pmod{10}$$

$$\Rightarrow 5 \equiv 0 \pmod{10} \Rightarrow \text{no solution}$$

$$\Rightarrow \nexists 2^{-1} \in \mathbb{Z}/10.$$

$\triangleright m^{-1} \bmod n$ doesn't exist, if $m|n$.

\rightarrow if $(m,n) \neq 1$.

- eg, $3^{-1} \bmod 10 \equiv 7$. [$\because 3 \times 7 \equiv 21 \equiv 1$]

$\triangleright \mathbb{Z}/n$ may not have "cancellation rule", i.e.

$$a \cdot b \equiv 0 \not\Rightarrow a \equiv 0 \vee b \equiv 0.$$

$$a \cdot b \equiv a \cdot c \not\Rightarrow a \equiv 0 \vee b \equiv c.$$

- Yes, $5^n \bmod 10 \equiv 5$ [$\because 5^2 \equiv 5$]

- $n = x_n x_{n-1} \dots x_0$ in base-10.

$$n \bmod 9 \equiv x_n + x_{n-1} + \dots + x_0$$

$$[\because n = x_n 10^n + x_{n-1} 10^{n-1} + \dots + x_0 \equiv \sum_i x_i.]$$

$$n \bmod 2 \equiv x_0.$$

$$n \bmod 11 \equiv \sum_i x_i \cdot (-1)^i.$$

- Ex: • Find $(x, y) \in \mathbb{Z}^2$: $25x + 31y = 6$.

• Find $(x, y) \in \mathbb{Z}^2$: $x^2 - 2y^2 = 1$.

Solving linear equations in \mathbb{Z}/n

- How do you solve $ax \equiv b \pmod{n}$?
- If $a^{-1} \pmod{n}$ exists, then $x \equiv ba^{-1} \pmod{n}$.

Lemma: $a^{-1} \pmod{n}$ is easy to compute (or say that it doesn't exist) & is unique.

Pf: • Use Euclid gcd algo. to write:

$$(a, n) =: l \cdot a + k \cdot n \Rightarrow \underline{l \cdot a} \equiv (a, n) \pmod{\underline{n}}$$

• If $(a, n) = 1 \Rightarrow l \equiv a^{-1} \pmod{n}$.

• Say, $(a, n) =: g > 1 \Rightarrow a^{-1} \pmod{n}$ doesn't exist!

• Suppose $\underline{ax} \equiv 1 \pmod{n}$
 $\Rightarrow \left(\frac{n}{g}\right) \cdot ax \equiv \left(\frac{n}{g}\right) \cdot 1 \pmod{n}$

$ax_1 \equiv 1 \equiv ax_2$
 $\Rightarrow a(x_1 - x_2) \equiv 0$
 $\Rightarrow x_1 - x_2 \equiv 0 \pmod{n}$
 \Rightarrow uniqueness.

$\Rightarrow 0 \equiv \left(\frac{n}{g}\right) \pmod{n} \Rightarrow \swarrow$
 $\Rightarrow a^{-1} \pmod{n}$ doesn't exist. □

Defn: - Given coprime (a, n) , we use $\underline{(a^{-1} \pmod{n})}$ to denote the unique inverse.

$\triangleright n$ is prime & $n \nmid a \Rightarrow (a, n) = 1 \Rightarrow a^{-1} \pmod{n}$ exists.
 $(a \not\equiv 0 \pmod{n})$

- Exs.

$$2^{-1} \pmod{11} \equiv 6$$

$$16^{-1} \pmod{13} \equiv 3^{-1} \equiv 9$$

$$92^{-1} \pmod{23} \equiv 0^{-1} \Rightarrow \text{?}$$

- Ex: $25^{-1} \pmod{23}$ using Euclid gcd algo.

Theorem (Fermat's little theorem, 1640): For prime p & $p \nmid a$, $a^{p-1} \equiv 1 \pmod{p}$.

Pf: $S := \{a, 2a, 3a, \dots, (p-1)a\} \pmod{p}$.

• Consider $\prod_{i \in [p-1]} (ia) = ?$

$\triangleright ia \equiv ja \pmod{p} \Rightarrow i \equiv j \pmod{p} [\Rightarrow |S| = p-1]$

[Pf: Say, $ia \equiv ja \Rightarrow (i-j)a \equiv 0 \pmod{p}$
 $\Rightarrow i-j \equiv 0.$]

$$\triangleright S = [p-1] = (\mathbb{Z}/p) \setminus \{0\}.$$

\Rightarrow

$$\triangleright a \cdot 2a \cdot 3a \cdots (p-1)a \equiv (p-1)! \pmod{p}.$$
$$\Rightarrow a^{p-1} \cdot (p-1)! \equiv \text{"} \pmod{p}$$

$$\Rightarrow a^{p-1} \equiv 1 \pmod{p}. \quad \square$$

- Qn: $a^n \equiv a \pmod{n}$? for composite n ?

- Ex. $n=4, a=3$: $3 \not\equiv 3^4 \equiv 1 \pmod{4}$
 $2 \not\equiv 2^4 \equiv 0 \pmod{4}$

- Qn: Computing $a^n \bmod n$?

▷ Easy to compute by repeated \rightarrow squaring
& $\bmod n$ arithmetic.

- Back to $a \underline{x} \equiv b \pmod n$:

1) Compute $(a, n) =: g$ [Euclid gcd]

2) If $g=1$ OUTPUT $x \equiv (b a^{-1} \pmod n)$

3) [$g \geq 2$] $(a/g) \cdot x \equiv (b/g) \pmod{(n/g)}$

3.1) If $g \nmid b$ then OUTPUT Fail.

3.2) $a' \cdot x' \equiv b' \pmod{n'}$ [$(a', n') = 1$]

3.3) Compute $x' \equiv b'/a' \pmod{n'}$ [$\Rightarrow ax' \equiv$

3.4) OUTPUT $\{x' + k \cdot n', 0 \leq k < g\}$. $b \pmod n$]

Euler's Totient Function $\varphi(n)$

- Numbers coprime to n , in (\mathbb{Z}/n) .

- Defn: $(\mathbb{Z}/n)^* := \{ a \in \mathbb{Z}/n \mid (a, n) = 1 \}$.
• $\varphi(n)$:= $|\mathbb{Z}/n^*|$.

- Ex. $\varphi(1) := 1 = \varphi(2)$; $\varphi(3) = 2 = \varphi(4) = \varphi(6)$
 $\varphi(5) = 4$.

$\triangleright \varphi(n) = n-1$, if n is prime (>1).

\triangleright For $n > 2$, $\varphi(n) \geq 2$.

- $\varphi(n) = \varphi(n') \not\Rightarrow n = n'$.

$$\triangleright \varphi(p^2) = p^2 - \#\{a \in \mathbb{Z}/p^2 \mid (a, p^2) > 1\}$$

$$\triangleright \varphi(p^k) = \overset{=}{p^k} - p^{k-1} = p^k \cdot \left(1 - \frac{1}{p}\right); \quad k \geq 1 \text{ \& prime } p.$$

Qn! For primes $p \neq q$, $\varphi(pq)$, $\varphi(p)$, $\varphi(q)$?

Theorem (Multiplicative): If $(m, n) = 1$, then

$$\varphi(mn) = \varphi(m) \cdot \varphi(n).$$

Pf: • $(\mathbb{Z}/m)^* \times (\mathbb{Z}/n)^* \xrightarrow{\psi} (\mathbb{Z}/mn)^*$ [bijection?]
 $(a, b) \mapsto \psi(a, b) =: c$

Idea: $\psi(a, b) := na + mb \pmod{mn}.$

(i) ψ is a valid map: $\cdot na + mb \pmod{mn}$
doesn't change as long as we pick
classes a & b .

• $na + mb \in (\mathbb{Z}/mn)^*$:
[Suppose not. $(na + mb, m) =: g$.
 $\Rightarrow (na, m) = g \Rightarrow (a, m) = g$
 $\Rightarrow g = 1$.]

(ii) ψ is injective : Suppose $na + mb \equiv na' + mb' \pmod{mn}$
 $\Rightarrow n(a - a') \equiv m(b' - b) \pmod{mn}$
 $\Rightarrow m \mid n(a - a') \Rightarrow m \mid (a - a') \Rightarrow \dots$
 $(a, b) \equiv (a', b') \in (\mathbb{Z}/m)^* \times (\mathbb{Z}/n)^*$.

(iii) ψ is surjective: Pick $c \in (\mathbb{Z}/mn)^*$.

Qn: Find (a, b) : $\psi(a, b) = c$?

$$na + mb \equiv c \pmod{mn}.$$

$$[\text{mod } m]: na \equiv c \pmod{m} \Leftrightarrow a \equiv c/n \pmod{m}$$

$$\Rightarrow mb \equiv c - na \pmod{mn} \quad \Rightarrow a \text{ found.}$$

$$\Rightarrow b \equiv \left(\frac{c - na}{m} \right) \pmod{n} \Rightarrow b \text{ found.}$$

[Check: $a \in (\mathbb{Z}/m)^*$ & $b \in (\mathbb{Z}/n)^*$.]

$\Rightarrow \psi$ is a bijection.

$$\Rightarrow \varphi(mn) = \varphi(m) \cdot \varphi(n).$$

□

Corollary: If $n = \prod_i p_i^{e_i}$ is the prime factorization

then

$$\varphi(n) = \prod_i \varphi(p_i^{e_i})$$

$$= \prod_i p_i^{e_i} \left(1 - \frac{1}{p_i}\right) = n \cdot \prod_i \left(1 - \frac{1}{p_i}\right)$$

$$\Rightarrow \varphi(n)/n = \prod_{\text{prime } p|n} \left(1 - \frac{1}{p}\right), \quad [n =: p_1^{e_1} \dots p_\ell^{e_\ell}]$$

— Alternate proof is by inclusion-exclusion:

$$\varphi(n) = \sum_{I \subseteq [l]} (-1)^{|I|} \cdot \left(\frac{n}{\prod_{i \in I} p_i} \right)$$

$$=: \sum_{d|n} \mu(d) \cdot \frac{n}{d}$$

where Möbius Function μ is defined as:

$$\mu(k) := \begin{cases} (-1)^r & , \text{ if } k = \text{product of } r \text{ primes} \\ 0 & , \text{ else} \\ 1 & , k=1 \end{cases}$$

Theorem (Euler, 1736): $\forall a \in (\mathbb{Z}/n)^*$,
 $a^{\varphi(n)} \equiv 1 \pmod{n}$.

Pf: • Redefine $S := \{a \cdot b \mid b \in (\mathbb{Z}/n)^*\}$.

• What's $\prod \{s \in S\} \equiv ? \pmod{n}$?

$$\equiv a^{\varphi(n)} \cdot \prod \{b \in (\mathbb{Z}/n)^*\}$$

$$\left[\begin{array}{l} \triangleright \forall s \in S, s \in (\mathbb{Z}/n)^* \\ \triangleright \forall b \neq b', ab \not\equiv ab' \pmod{n} \end{array} \right]$$

$$\equiv \prod \{b \in (\mathbb{Z}/n)^*\}$$

$$\Rightarrow a^{\varphi(n)} \equiv 1 \pmod{n} \quad \square$$

More linear equations (over \mathbb{Z}/n)

- How to solve: $(a_1 x \equiv b_1 \pmod{m}) \wedge (a_2 x \equiv b_2 \pmod{n})$

$$- \text{eg. } x \equiv 1 \pmod{4} \wedge x \equiv 2 \pmod{6} \Rightarrow \text{⚡}$$

- So, $(m, n) > 1$ creates a problem. What if $(m, n) = 1$?

- Further, assume $a_i = 1, \forall i$; i.e.
 $x \equiv b_1 \pmod{m} \wedge x \equiv b_2 \pmod{n}$.

- Assume $(m, n) = 1$:

Theorem 1 (Chinese Remainder thm, ~300 AD): An $x \in \mathbb{Z}$ exists. Moreover, $x \pmod{mn}$ is

Unique!

Pf: • Consider $\psi': \mathbb{Z}/m \times \mathbb{Z}/n \rightarrow \mathbb{Z}/mn$
 $(b_1, b_2) \mapsto \psi'(b_1, b_2) \pmod{mn}$

- Previously, $\psi(b_1, b_2) = b_1 n + b_2 m$.

- To correct ψ we redefine ψ' as:

$$\psi'(b_1, b_2) := b_1 \cdot n \cdot (n^{-1} \bmod m) + b_2 \cdot m \cdot (m^{-1} \bmod n)$$

$$\triangleq \begin{matrix} \equiv & \uparrow & b_1 & \bmod m \\ \equiv & \uparrow & b_2 & \bmod n \end{matrix}$$

$\Rightarrow x$ has solution $\psi'(b_1, b_2)$.

\triangle Moreover, ψ' is a bijection!

$\Rightarrow x \bmod mn$ is unique ($\because \psi'$ injective).

Exercise: Set of modular linear equations
can be solved efficiently. \square

Theorem 2 (CRT): Let n_1, n_2, \dots, n_k be mutually coprime integers. Then, the system $x \equiv c_i \pmod{n_i}, \forall i \in [k]$ has a unique solution $x \pmod{\prod_i n_i}$ (& can be easily computed).

Pf:

Idea:
$$\psi'(c_1, \dots, c_k) := \left(\sum_{j=1}^k c_j \cdot N_j \right) / \left(\sum_{j=1}^k N_j \right) \pmod{N}.$$

Where $N := \prod_{i=1}^k n_i$; $N_j := N/n_j$.

- Complete the pf.

□