

## 1 Recall

**Fact 1.1**  $\text{Res}(f, g) = 0$  iff  $\text{gcd}(f, g) > 1$

**Fact 1.2** There exists  $y \in F_q$  such that  $\text{gcd}(e(x) - y, f(x)) > 1$

We want  $y \in F_q$  such that  $\text{Res}(e(x) - y, f(x)) = 0$

$\text{Res}(e(x) - y, f(x))$  is a polynomial in  $y$  over  $F_q$  of degree  $\leq 2d - 1$

Let this polynomial be  $g(y)$ .

If we can find a root of  $g$  in  $F_q$  then we can factorize  $f$ .

Let  $\hat{g}(y) = \text{gcd}(g(y), y^q - y)$

All roots of  $g(y)$  in  $F_q$  are roots of  $\hat{g}(y)$  too.

Now, the remaining problem is to find roots of a given polynomial over a finite field  $F_q$ .

No polynomial time algorithm is known for this problem.

## 2 A Randomized polynomial time algorithm for root finding

Let  $f(x)$  be a square-free polynomial over  $F_q$  of degree  $d$  and such that  $f$  factors completely over  $F_q$ .

Let  $f(x) = \prod_{i=1}^d (x - \alpha_i)$

Note that  $\alpha_i \neq \alpha_j$ .

Let  $f_{ss}(x) = f(x^2 + \beta) = \prod_{i=1}^d (x^2 + \beta - \alpha_i)$

If there exist  $\alpha_i$  and  $\alpha_j$  such that  $x^2 + \beta - \alpha_i$  is reducible and  $x^2 + \beta - \alpha_j$  is irreducible, then  $f_{ss}$  can be factored.

Using factors of  $f_{ss}$ ,  $f$  can be factored.

Fix  $\{\alpha_i, \alpha_j\} = \{\alpha_1, \alpha_2\}$

$$\begin{aligned}
& \text{Prob}[x^2 + \beta - \alpha_1 \text{ and } x^2 + \beta - \alpha_2 \text{ are both reducible or irreducible}] \\
&= \text{Prob}[\text{both } \alpha_1 - \beta \text{ and } \alpha_2 - \beta \text{ are squares in } F_q \text{ or neither is}] \\
&= \text{Prob}[\beta \in F_q : (\alpha_1 - \beta)^{\frac{q-1}{2}} = (\alpha_2 - \beta)^{\frac{q-1}{2}}] \\
&= \frac{1}{|F_q|} (\text{number of roots of polynomial } (\alpha_1 - z)^{\frac{q-1}{2}} - (\alpha_2 - z)^{\frac{q-1}{2}}) \\
&\leq \frac{q-1}{2q} < \frac{1}{2}
\end{aligned}$$

Choose  $k$  values of  $\beta$ .

$$\text{Prob}[\text{no value of } \beta \text{ helps factor } f_{ss}] < \frac{1}{2^k}$$

Repeating this algorithm makes the probability of error very small.

Roots of  $f$  can be computed using repeated applications of the algorithm.

There exist randomized polynomial time algorithms for factoring multivariate polynomials in compact representation.

A polynomial over the field of rationals can be factored in polynomial time.