



CS618: Program Analysis

2016-17 1st Semester

Interprocedural Data Flow Analysis

Functional Approach

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Interprocedurally Valid Paths

- ▶ G^* ignores the special nature of call and return edges
- ▶ Not all paths in G^* are feasible
 - ▶ do not necessarily represent potentially valid execution paths
- ▶ $\text{IVP}(r_1, n)$: set of all interprocedurally valid paths from r_1 to n
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↳ $\text{path}_{G^*}(r_1, n) \subseteq \text{IVP}(r_1, n)$ but $\text{IVP}(r_1, n) \not\subseteq \text{path}_{G^*}(r_1, n)$



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Proper sequence

- ▶ q_1 without any return edge is proper
- ▶ let $q_1[i]$ be the first return edge in q_1 . q_1 is proper if

$i = 1, 2, \dots, n$

$\forall i \in [n], \exists j \in [n] : q_1[i] = q_1[j] \text{ and } q_1[i] \neq q_1[1]$

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- ▶ q_1 without any return edge is proper
- ▶ let $q_1[i]$ be the first return edge in q_1 . q_1 is proper if
 - ▶ $i > 1$; and
 - ▶ $q_1[i - 1]$ is call edge corresponding to $q_1[i]$; and
 - ▶ q'_1 obtained from deleting $q_1[i - 1]$ and $q_1[i]$ from q_1 is proper



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Interprocedurally Valid Complete Paths

- ▶ $\text{IVP}_0(r_p, n)$ for procedure p and node $n \in N_p$
 - ▶ set of all interprocedurally valid paths q in G^* from r_p to n s.t.
 - ▶ Each call edge has corresponding return edge in q restricted to E^*



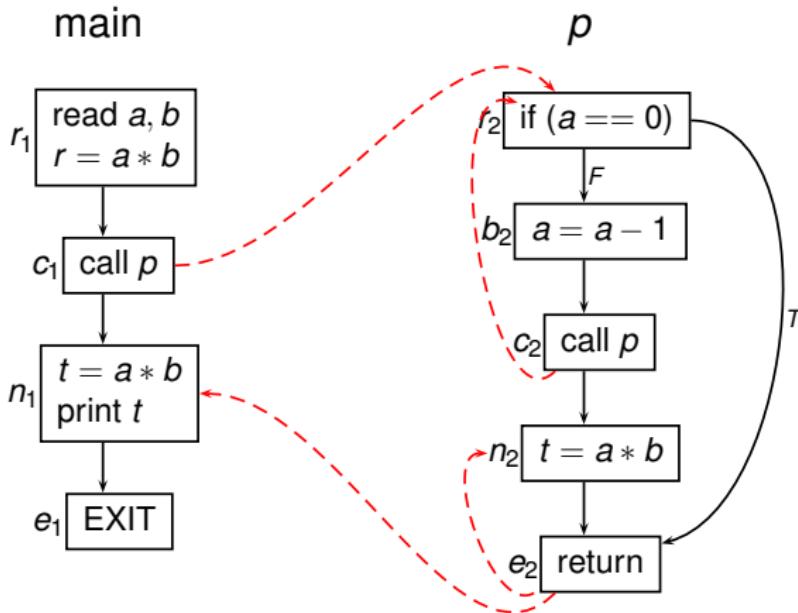
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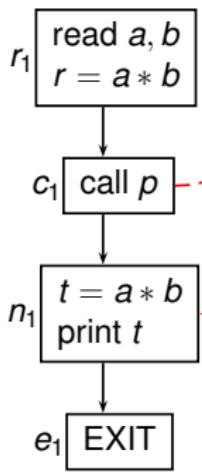
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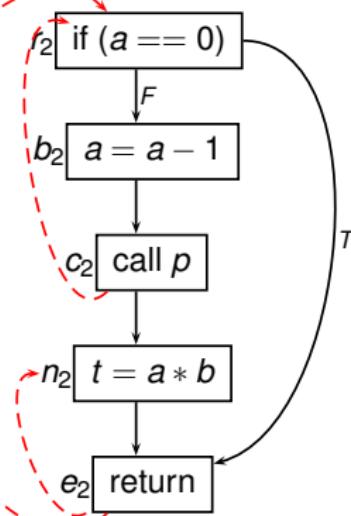




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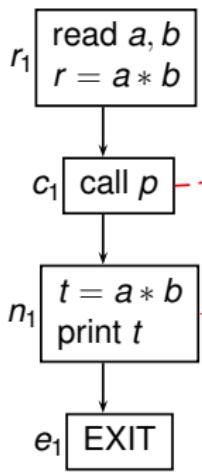
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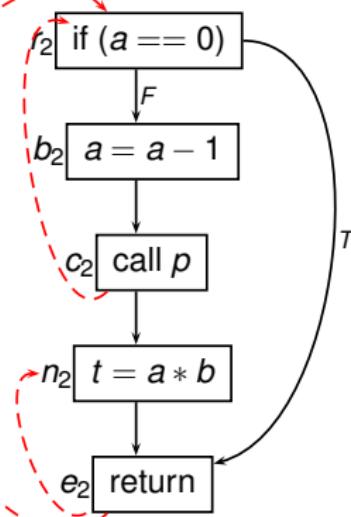
$r_1 \rightarrow c_1 \rightarrow r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow e_2 \rightarrow n_2 \rightarrow e_2 \rightarrow n_1 \rightarrow e_1$
 $\in \text{IVP}(r_1, e_1)$



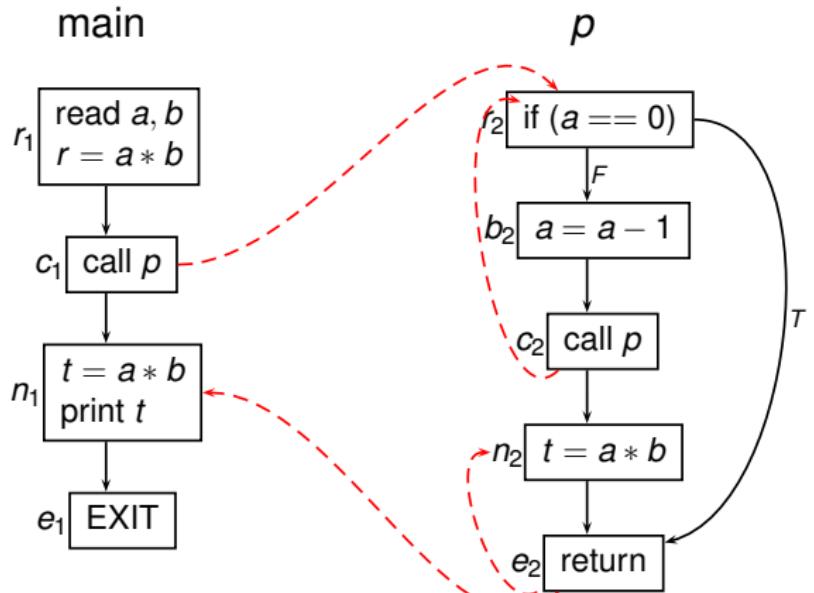
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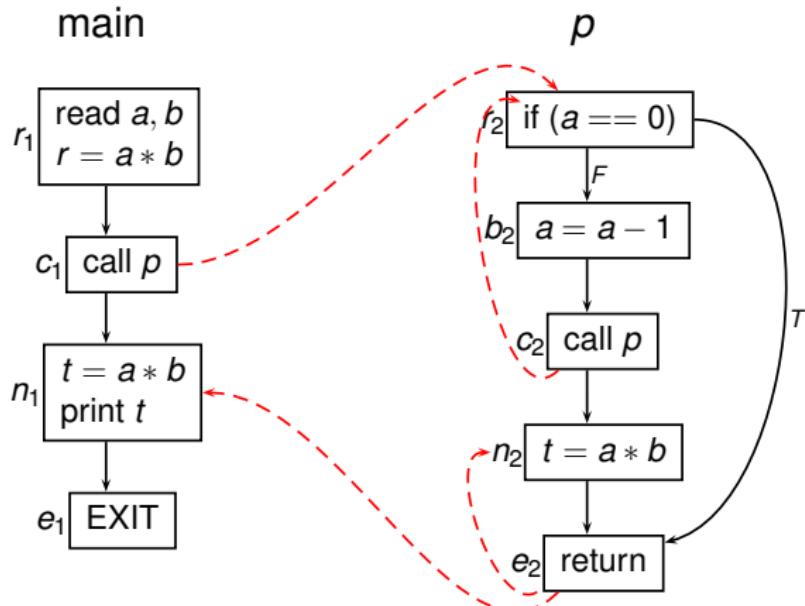
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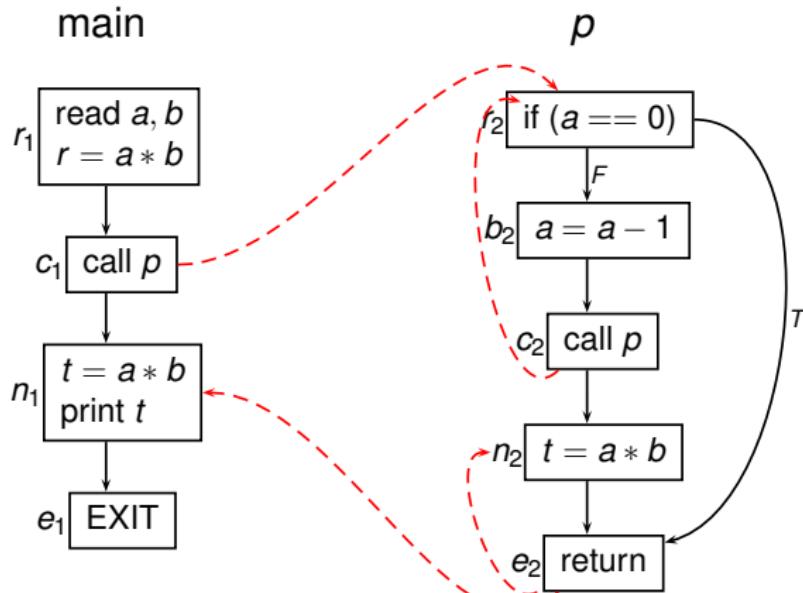


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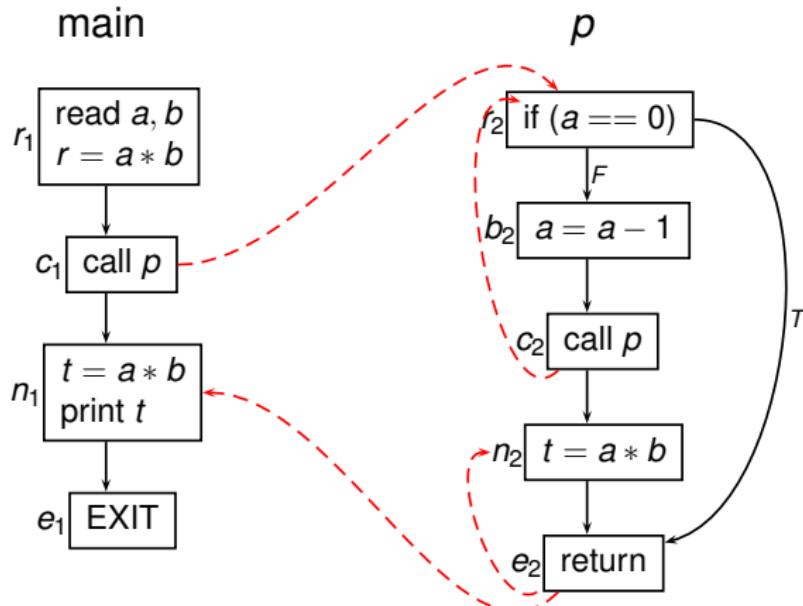


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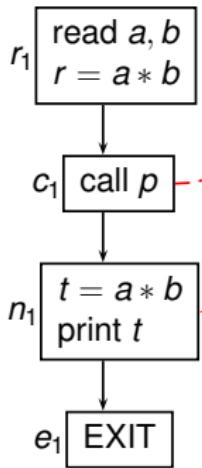
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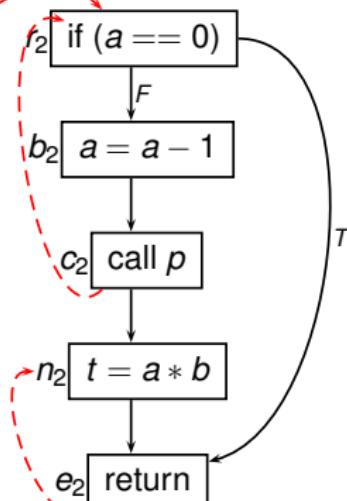
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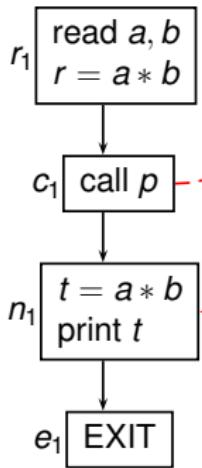
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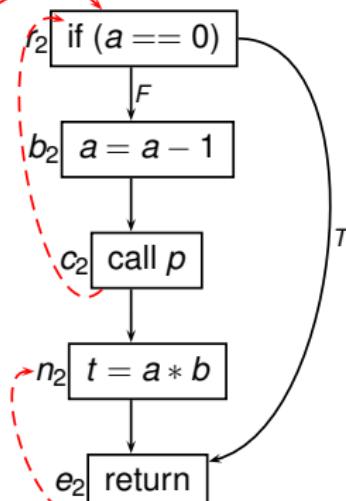
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Path Decomposition

$$q \in \text{IVP}(r_{\text{main}}, n)$$

\Leftrightarrow

$$q = q_1 \parallel (c_1, r_{p_2}) \parallel q_2 \parallel \cdots \parallel (c_{j-1}, r_{p_j}) \parallel q_j$$

where for each $i < j$, $q_i \in \text{IVP}_0(r_{p_i}, c_i)$ and $q_j \in \text{IVP}_0(r_{p_j}, n)$



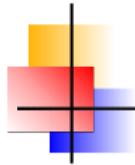
Functional Approach

- ▶ (L, F) : a *distributive* data flow framework
 - ▶ Procedure p , node $n \in N_p$
 - ▶ $\phi_{(r_p, n)} \in F$ describes flow of data flow information from start of r_p to start of n



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- ▶ $\phi_{(r_p, n)} \in F$ describes flow of data flow information from start of r_p to start of n
 - ↳ $\phi_{(r_p, n)} = \phi_{(r_p, n)}(P_0(r_p, n))$



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Functional Approach Constraints

$$\phi_{(r_p, r_p)} \leq id_L$$

$$\phi_{(r_p, n)} = \bigwedge_{(m, n) \in E_p} (h_{(m, n)} \circ \phi_{(r_p, m)}) \quad \text{for } n \in N_p$$

where

$$h_{(m, n)} = \begin{cases} f_{(m, n)} & \text{if } (m, n) \in E_p^0, \\ & f_{(m, n)} \in F \text{ associated flow function} \\ \phi_{(r_q, e_q)} & \text{if } (m, n) \in E_p^1 \text{ and } m \text{ calls procedure } q \end{cases}$$

Information x at r_p translated to information $\phi_{(r_p, n)}(x)$ at n



Solving ϕ Constraints

- ▶ Round-robin iterative approximations to initial values

$$\phi_{(r_p, r_p)}^0 \leq id_L$$

$$\phi_{(r_p, n)}^0 \leq f_\Omega \quad n \in N_p - \{r_p\}$$

- ▶ Reach maximal fixed point



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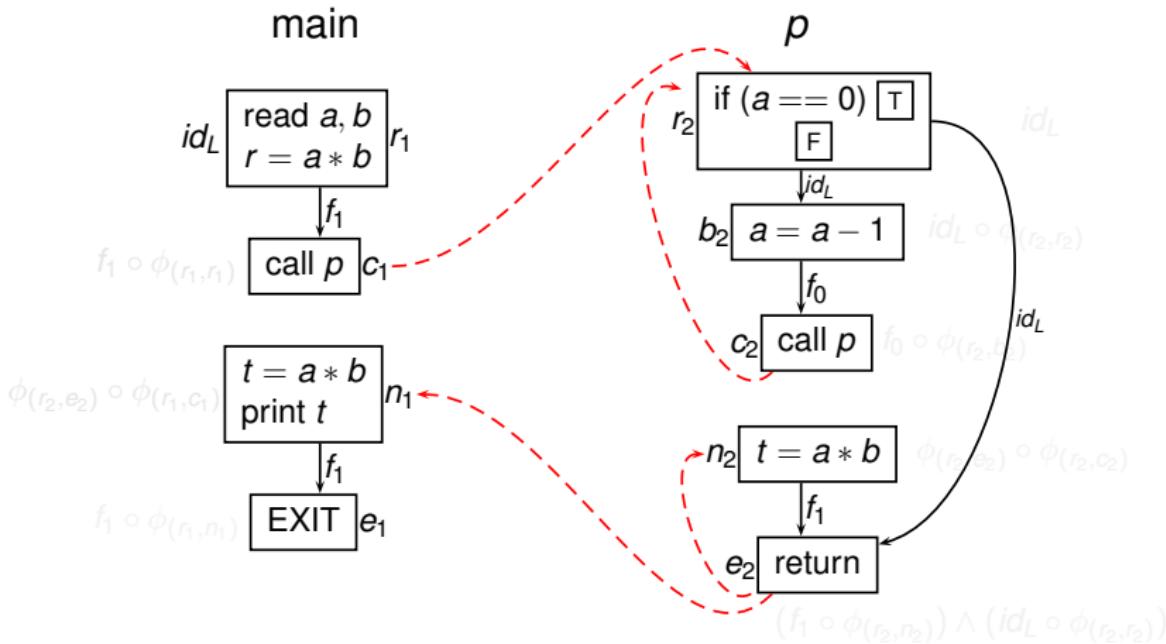
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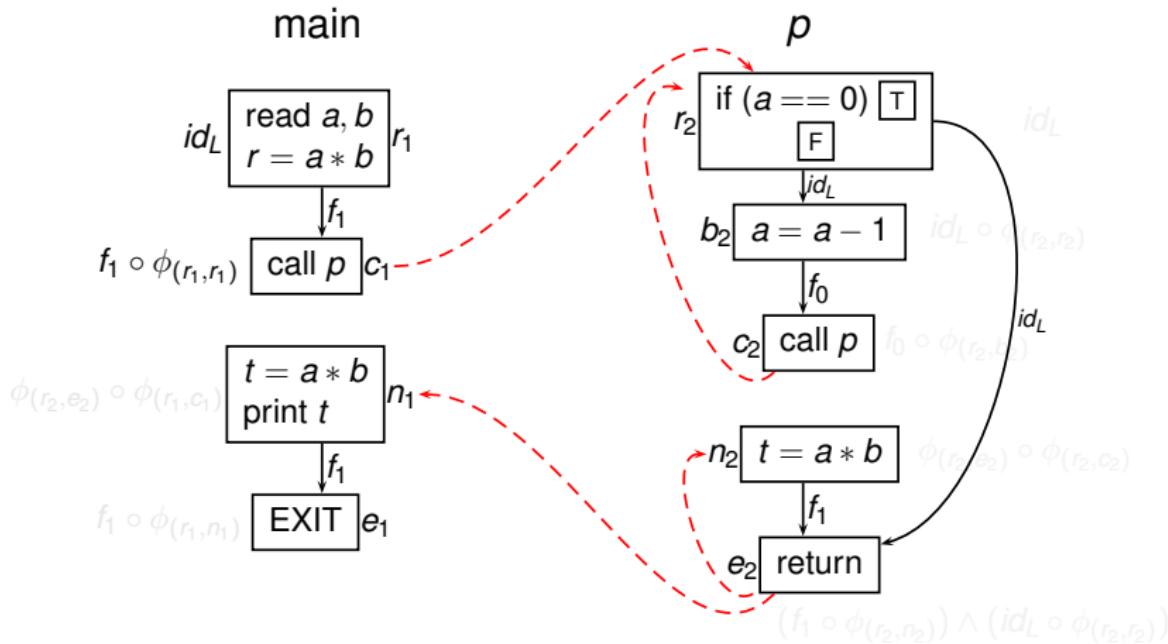
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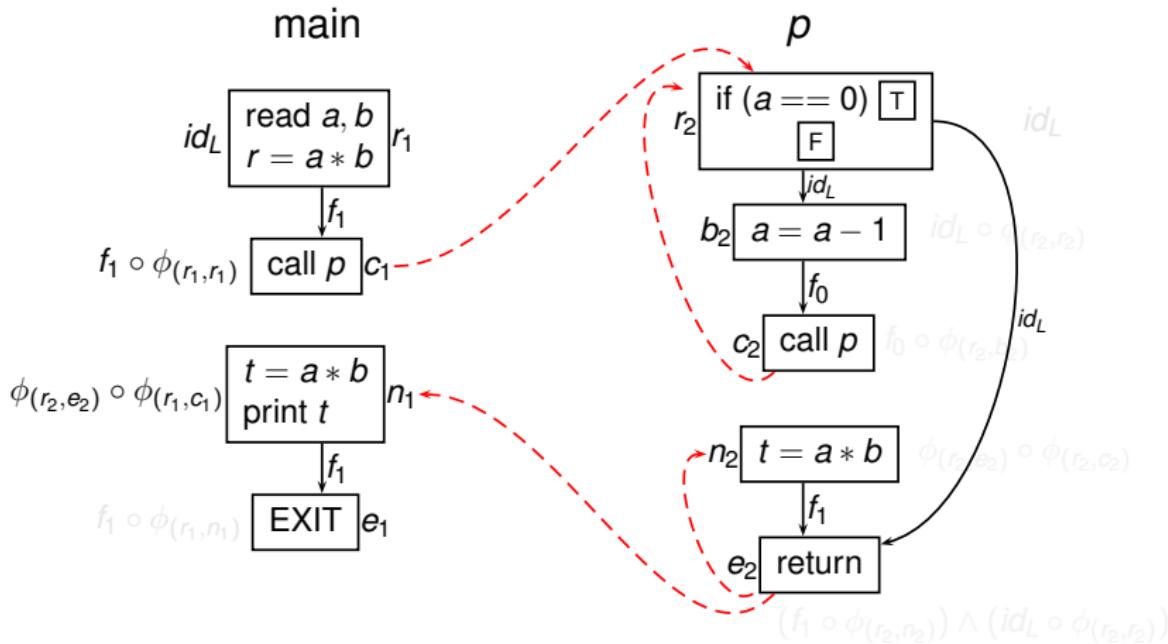
Example



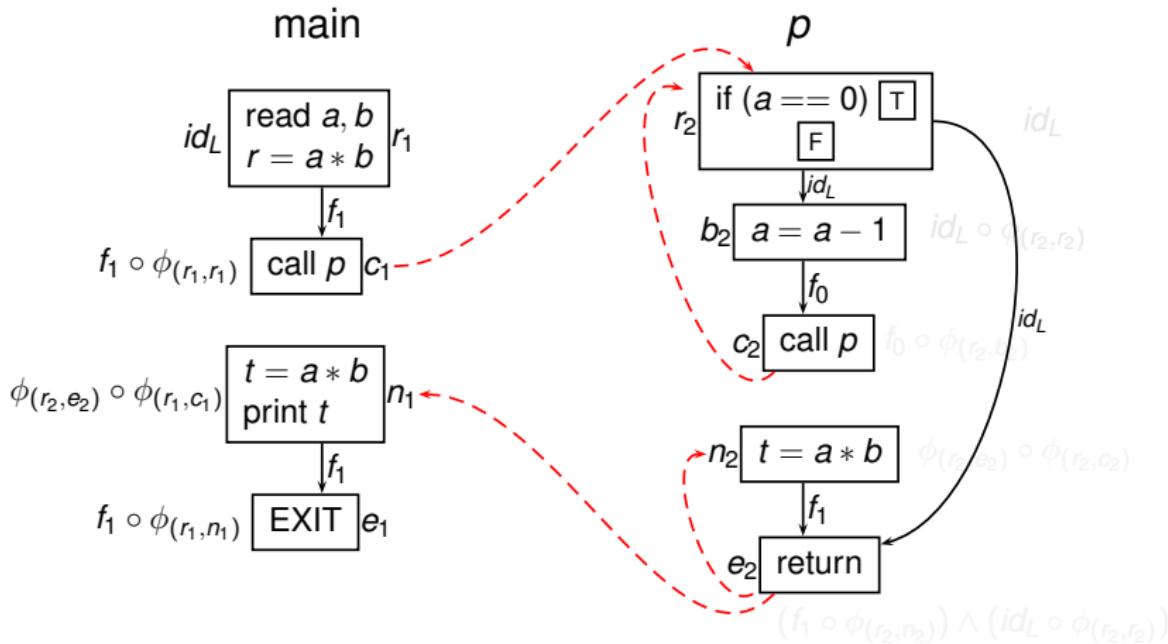
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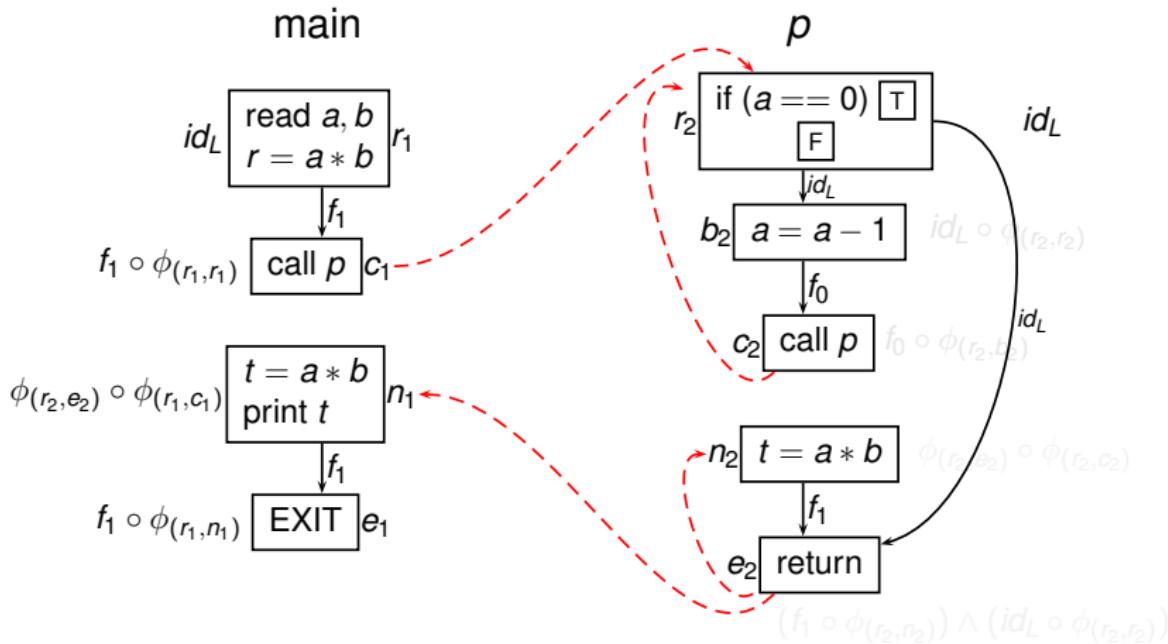


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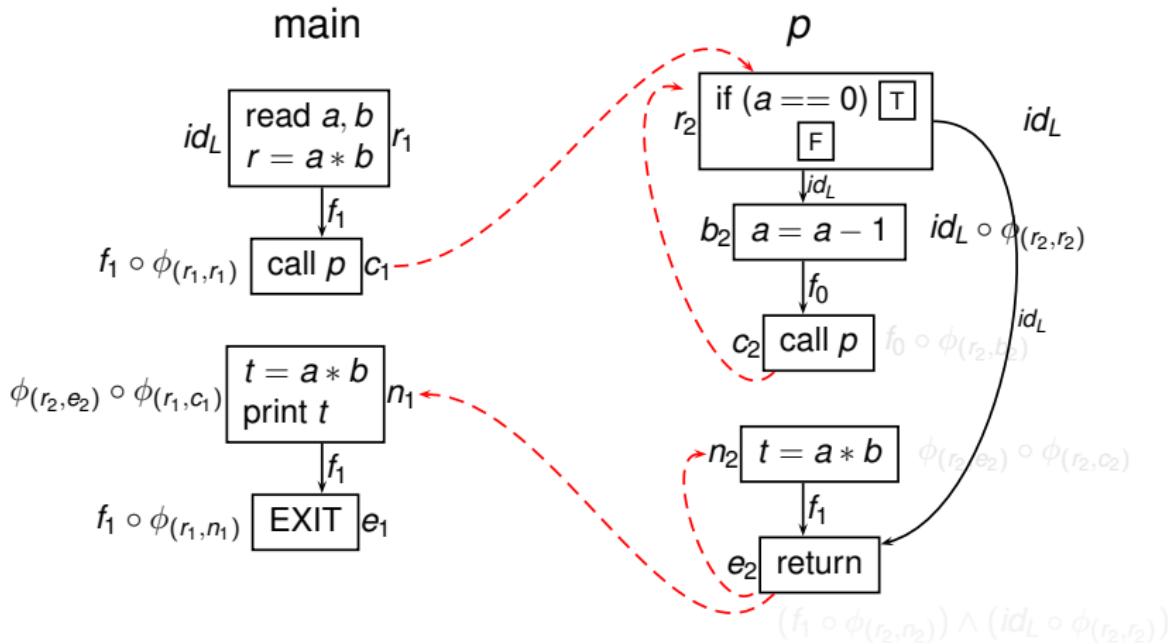


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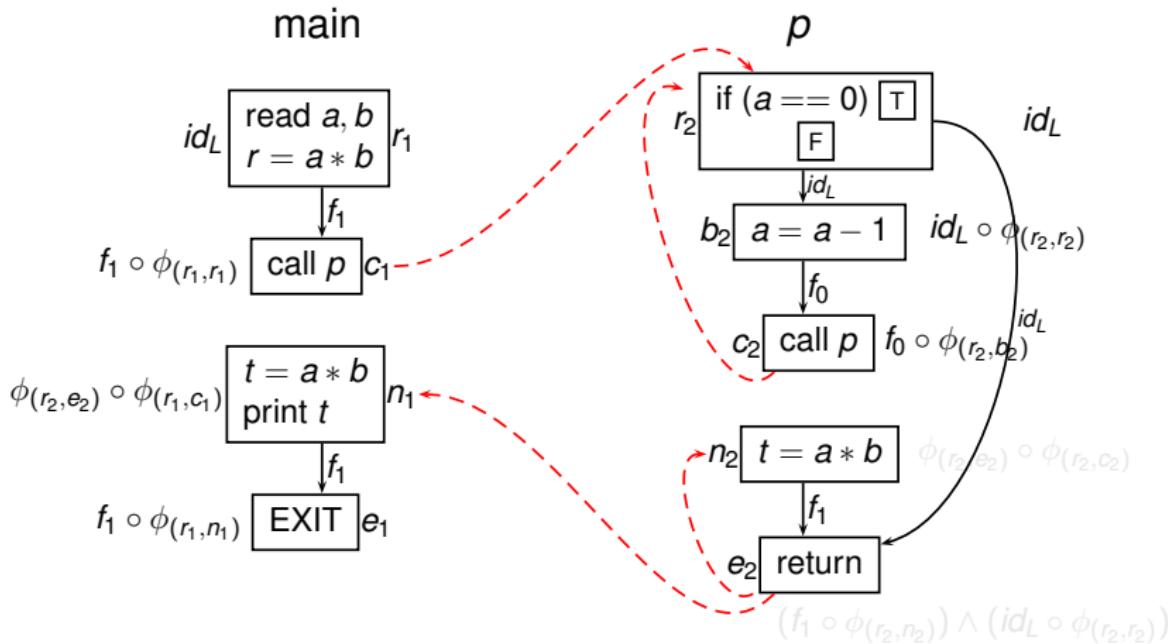


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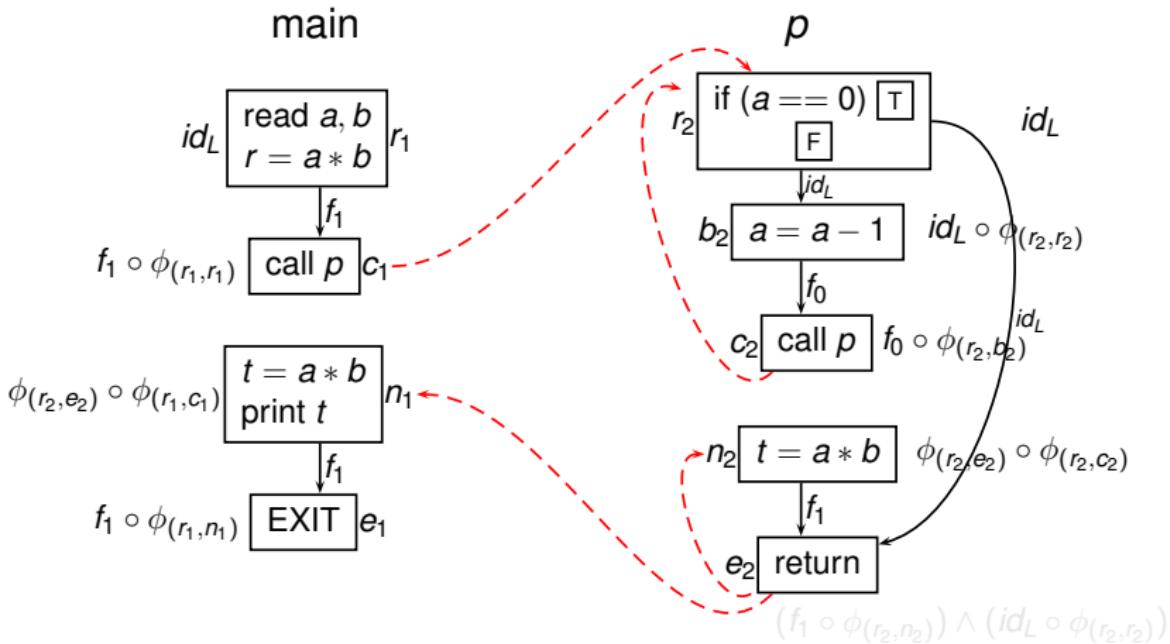




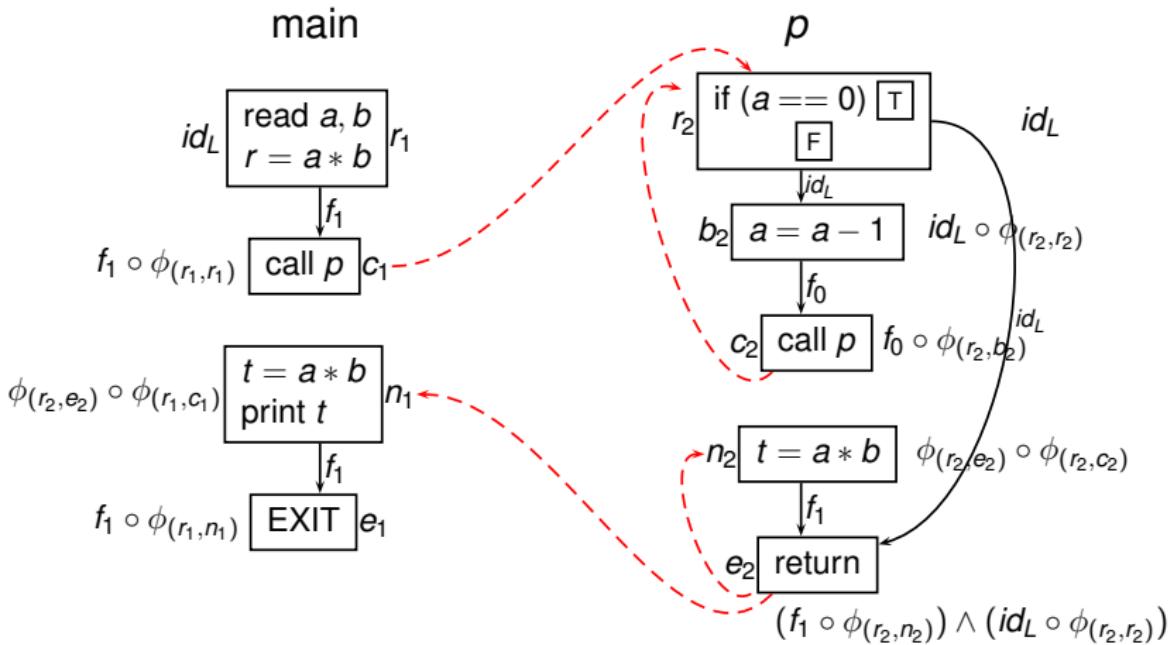
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Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
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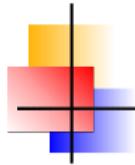
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$\phi(r_2, b_2)$	$id_L \circ \phi(r_2, r_2)$	f_Ω	id	id	id
$\phi(r_2, c_2)$	$f_0 \circ \phi(r_2, b_2)$	f_Ω	f_0	f_0	f_0
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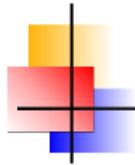
Iterative Solution

Function	Constraint	Iteration #			
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$\phi(r_1, r_1)$	id_L	id	id	id	id
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Solving Data Flow Problem

- ▶ The above process gives solution to ϕ functions
- ▶ Use it to compute data flow information x_n associated with start of block n

$$x_{r_{\text{main}}} = \textit{BoundaryInfo}$$

for each procedure p

$$x_{r_p} = \bigwedge \left\{ \phi_{(r_q, c)}(x_{r_q}) : \begin{array}{l} q \text{ is a procedure and} \\ c \text{ is a call to } p \text{ in } q \end{array} \right\}$$

$$x_n = \phi_{(r_p, n)} \quad n \in N_p - \{r_p\}$$

- ▶ Iterative algorithm for solution, maximal fixed point solution



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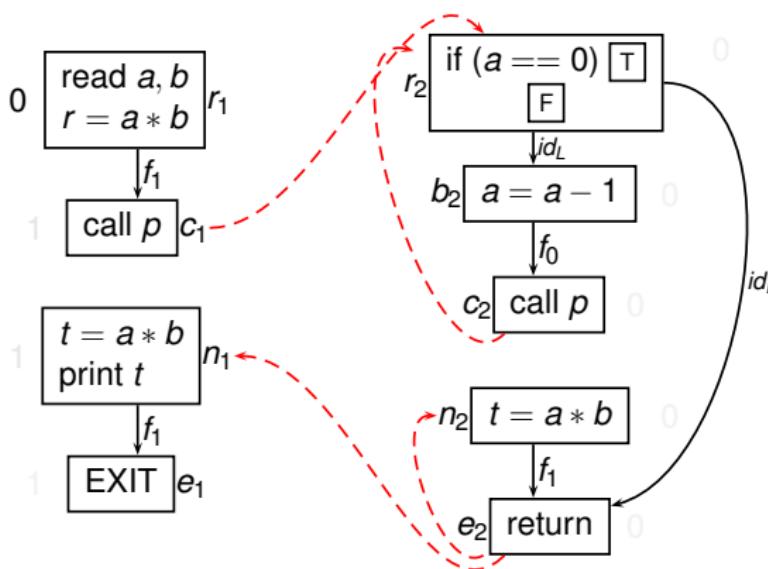
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Example

main

p



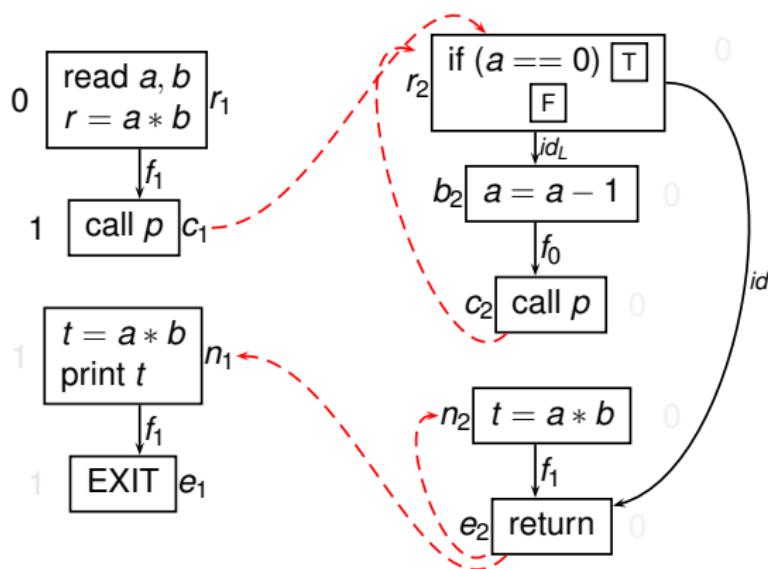
Functions

$\phi(r_1, r_1)$	$= id_L$
$\phi(r_1, c_1)$	$= f_1$
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$\phi(r_1, e_1)$	$= f_1$
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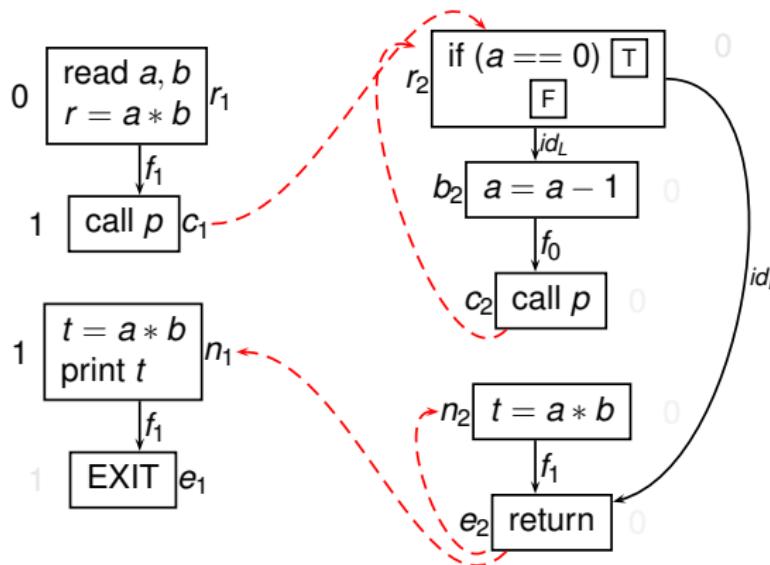
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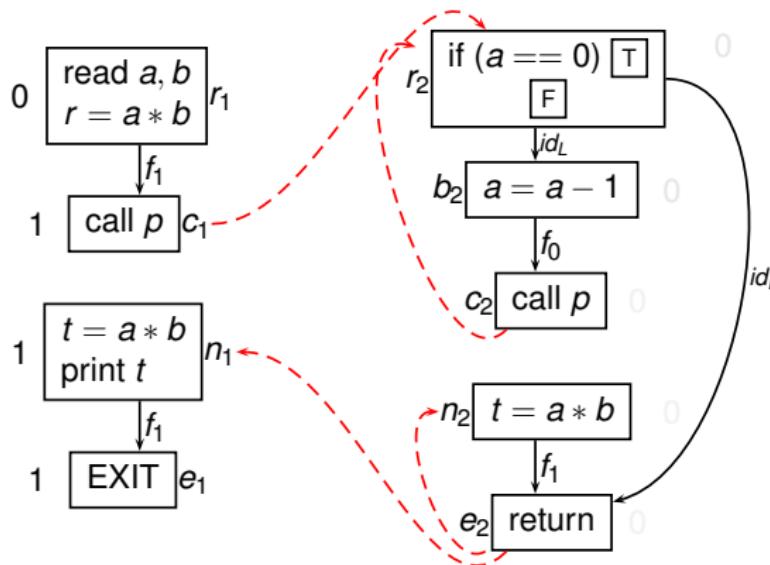
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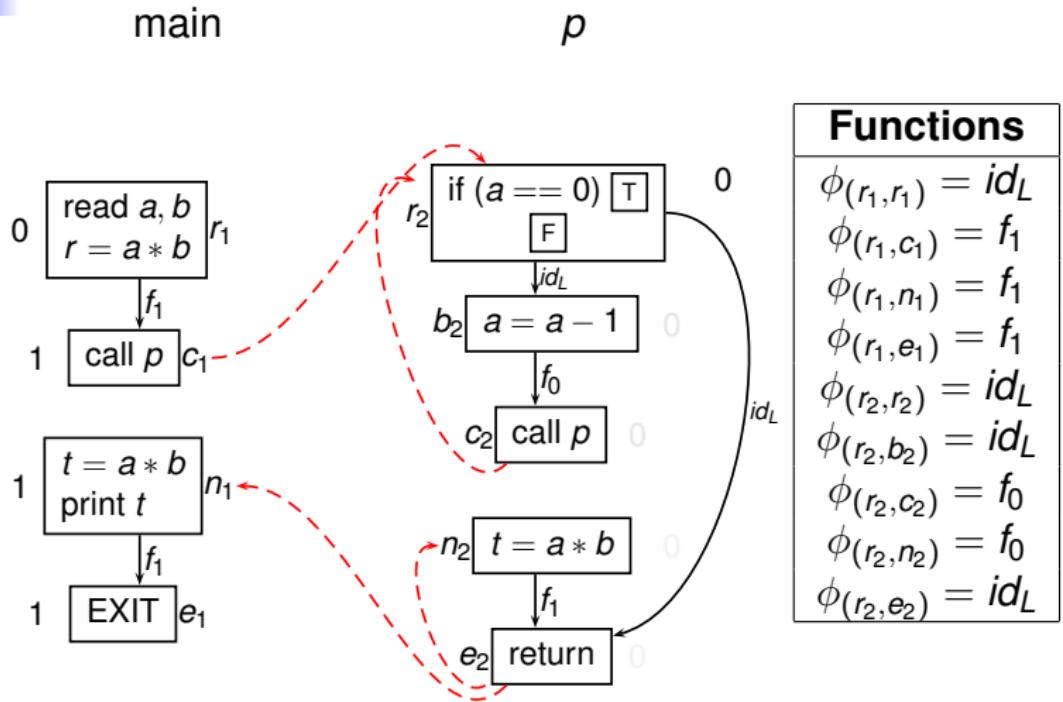
p



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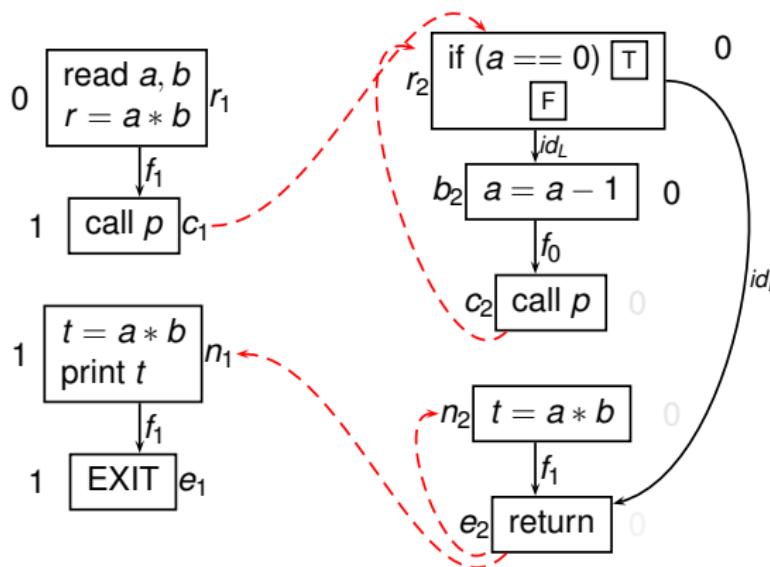
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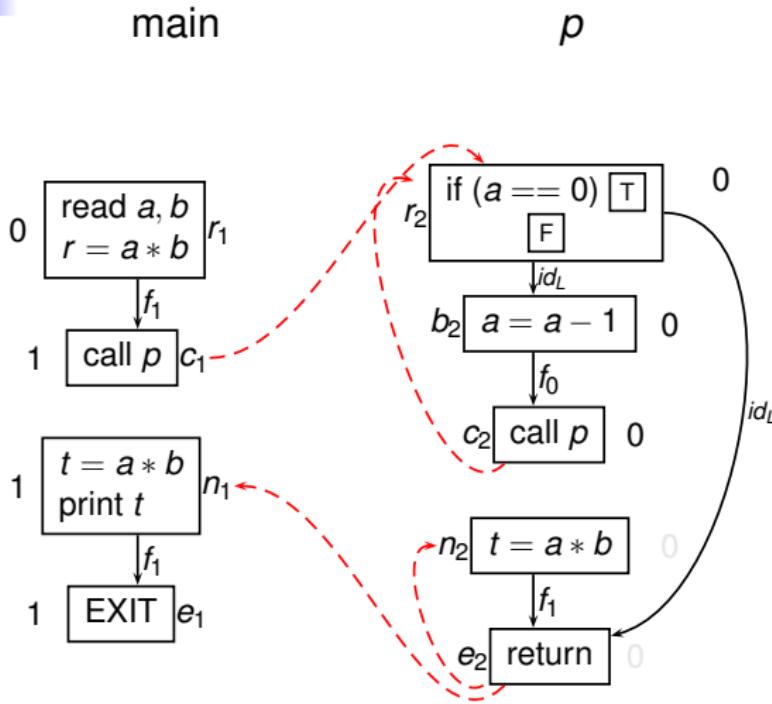
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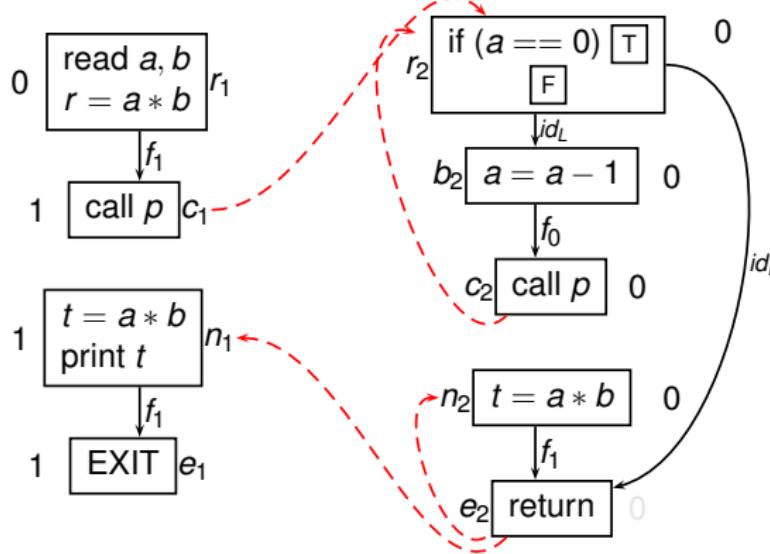
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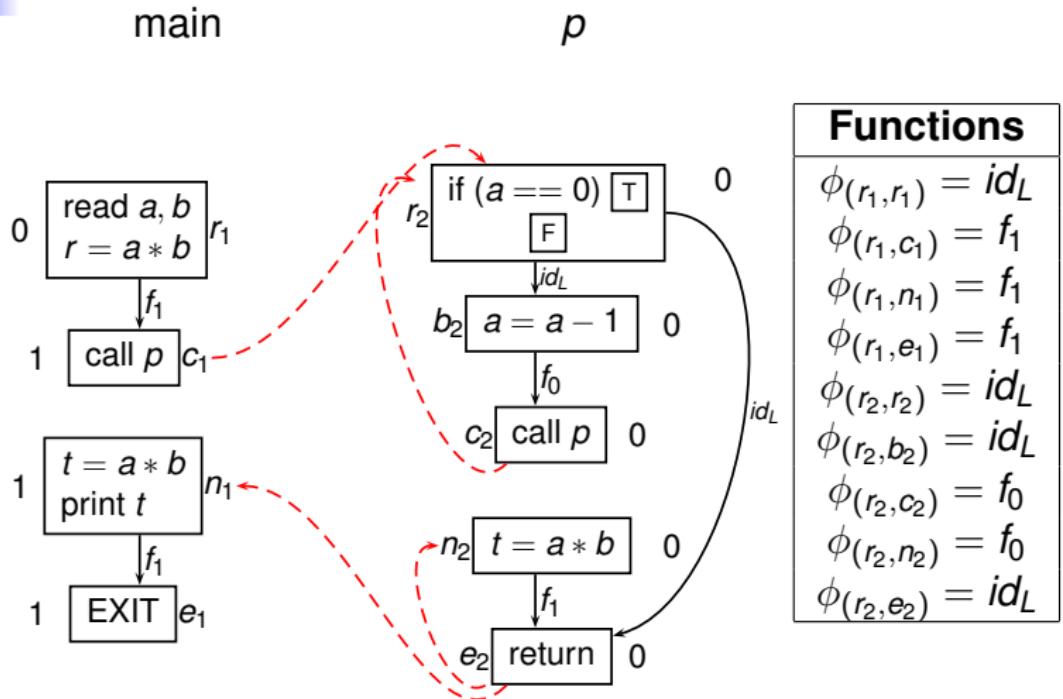
p



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Example





Interprocedural MOP

$$\begin{aligned}\Psi_n &= \bigwedge \{f_q : q \in \text{IVP}(r_{\text{main}}, n)\} \in F \quad \forall n \in N^* \\ y_n &= \Psi_n(\text{BoundaryInfo}) \quad \forall n \in N^*\end{aligned}$$

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Proof: By induction (**Exercise/Reading Assignment**)



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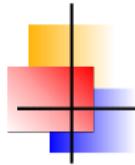
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Then

$$\Psi_n = \mathcal{X}_n$$

Proof: IVP₀ Lemma and Path decomposition

$$y_n = \Psi_n(\text{BoundaryInfo}) = \mathcal{X}_n(\text{BoundaryInfo})$$



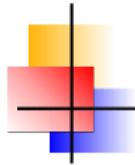
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MOP vs MFP

- ▶ F is distributive $\Rightarrow MFP = MOP$
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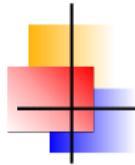
Practical Issues

- ▶ How to compute ϕ s effectively?
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 - ▶ Encoding
 - ▶ Compositions
 - ▶ Does the solution process converge?
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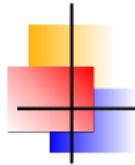
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