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biocomputing. image processing, graph grammars, graph transformation, and pattern recognition and computer vision, computer-aided design, pattern matching that finds practical application in areas such as Subgraph isomorphism is an important and very general form of

graph distance will be reviewed. be discussed and recent results relating subgraph isomorphism, maximum common subgraph, minimum common supergraph, and In this talk, several problems related to subgraph isomorphism will

- Read, R. C. and Corneil, D. G. (1977). The graph isomorphism disease. Journal of Graph Theory 1, 339-363.
- Gati, G. (1979). Further annotated bibliography on the isomorphism disease. Journal of Graph Theory 3, 95–109.

- Contents
- Introduction
- Non-mathematical motivation
- A hierarchy of pattern matching problems
- Solving NP-complete problems
- Complexity and approximation results
- Recent results on exact solutions
- Recent results on approximate solutions
- Open problems

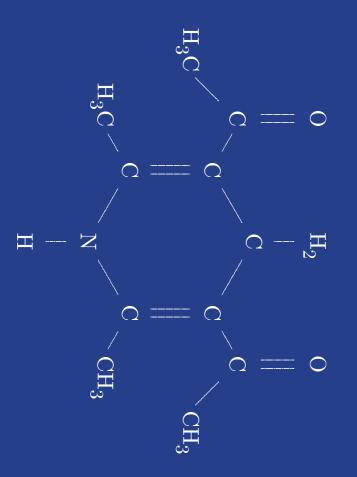
INTRODUCTION

- Mathematical motivation
- NP-complete problems are a challenge to theoretical computer science
- Non-mathematical motivation
- Pattern recognition and computer vision
- Computer-aided design
- Image processing
- Graph grammars and graph transformation
- Biocomputing

- Subgraph isomorphism is a common generalization of many important graph problems
- Clique
- Independent set
- Hamiltonian cycle
- Matching
- Girth
- Shortest path

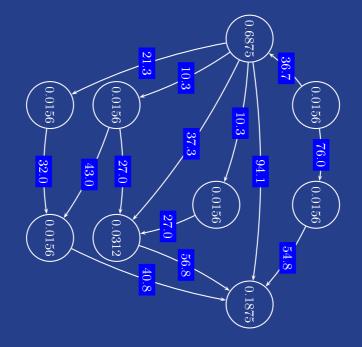
- Variations of subgraph isomorphism have been used to model important practical problems
- Information retrieval
- Scene analysis
- Computer-aided design
- Pattern recognition
- Graph grammars and graph transformation

#### Information retrieval



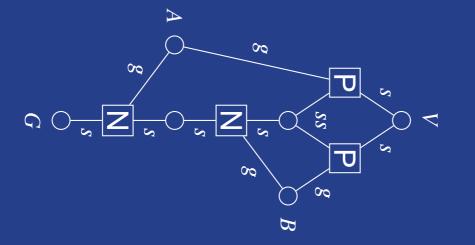
#### Scene analysis

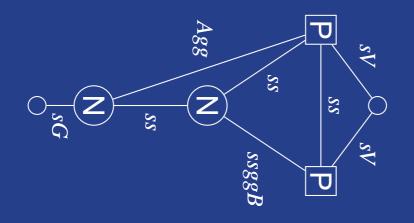
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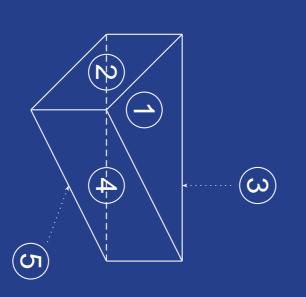
Baeza-Yates, R. and Valiente, G. (2000). An image similarity String Processing and Information Retrieval (2000), pp. 28–38. IEEE Computer Science Press. measure based on graph matching. In Proc. 7th Int. Symp. on

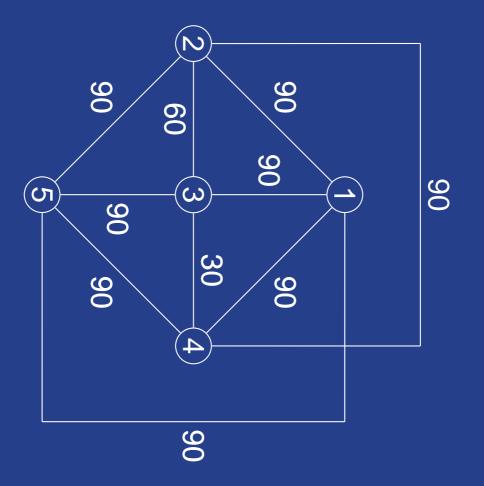
#### Computer-aided design



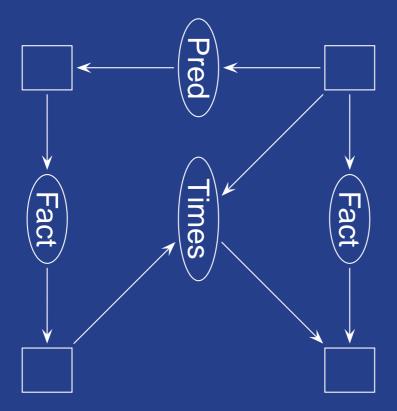


#### Pattern recognition



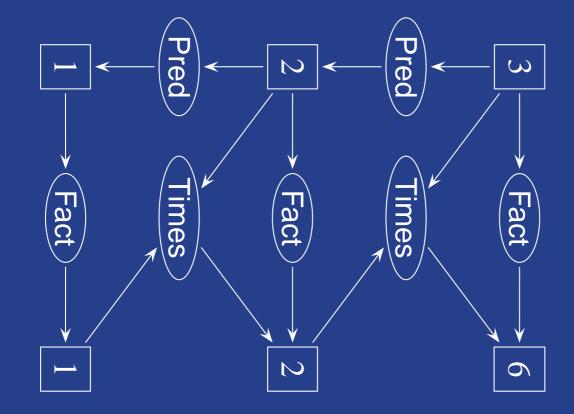


Graph grammars and graph transformation





 Graph grammars and graph transformation



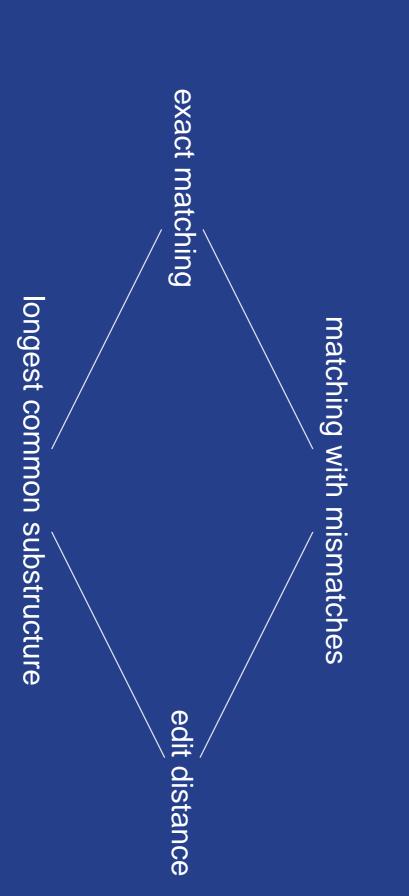
Subgraph isomorphism is an important and very general form of exact

pattern matching

- String searching
- Sequence alignment
- Tree comparison
- Pattern matching on graphs

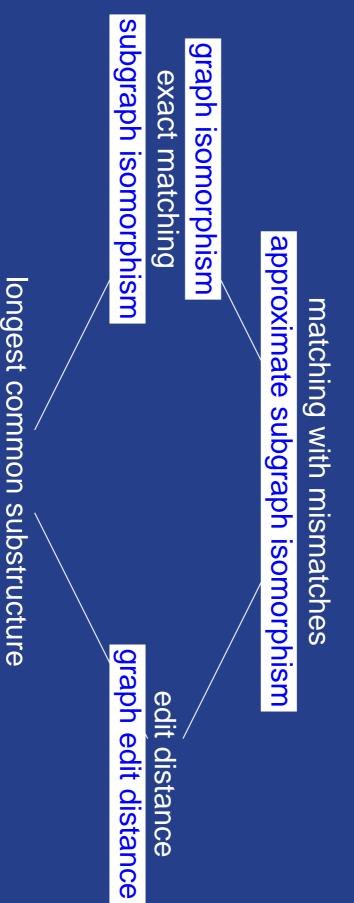
- A hierarchy of pattern matching problems
- Graph isomorphism
- Subgraph isomorphism
- Maximum common subgraph
- Approximate subgraph isomorphism
- Graph edit distance

A hierarchy of pattern matching problems





A hierarchy of pattern matching problems



maximum common subgraph

- Given a pattern G and a text H
- Decision problem

Answer whether H contains a subgraph isomorphic to G

Search problem

Return an occurrence of G as a subgraph of H

Counting problem

isomorphic to G Return a count of the number of subgraphs of H that are

Enumeration problem

Return all occurrences of G as a subgraph of H

- Given a pattern G and a text H
- General problem

Both G and H are input graphs

Restricted problem

such as trees or planar graphs Both G and H are input graphs belonging to a particular class,

Fixed problem

G is an input graph but H is a fixed graph, or vice versa

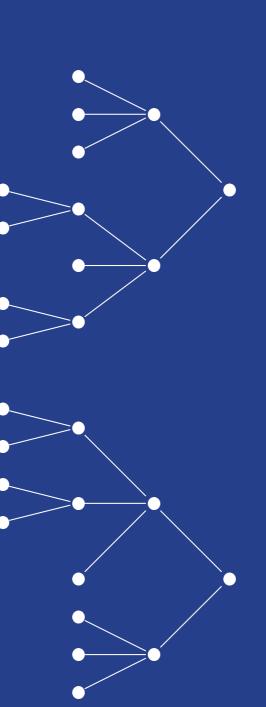
- Solving NP-complete problems
- Combinatorial enumeration
- Special cases
- Restricted
- Fixed
- Local search
- Probabilistic analysis
- Approximation algorithms
- Heuristics

pair of vertices  $v_i, v_j \in V_1$ ,  $(v_i, v_j) \in E_1$  if and only if  $(\varphi(v_i), \varphi(v_j)) \in E_2$ denoted by  $G_1 \cong G_2$ , if there is a bijection  $\varphi: V_1 \to V_2$  such that, for every **Definition.** Two graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  are isomorphic,

For input graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  with  $V_1 = \{u_1, ..., u_n\}$ and  $V_2 = \{v_1, \ldots, v_n\}$ , a necessary condition for  $G_1 \cong G_2$  is that the multisets  $\{\Gamma(u_i) \mid 1 \leqslant i \leqslant n\}$  and  $\{\Gamma(v_i) \mid 1 \leqslant i \leqslant n\}$  be equal

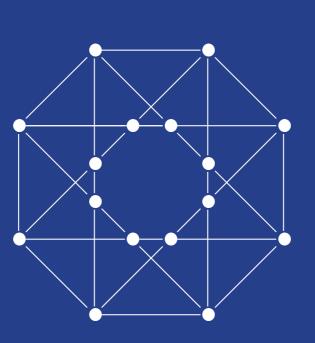
#### The Subgraph Isomorphism Problem

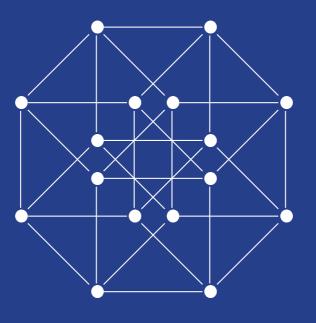




#### The Subgraph Isomorphism Problem

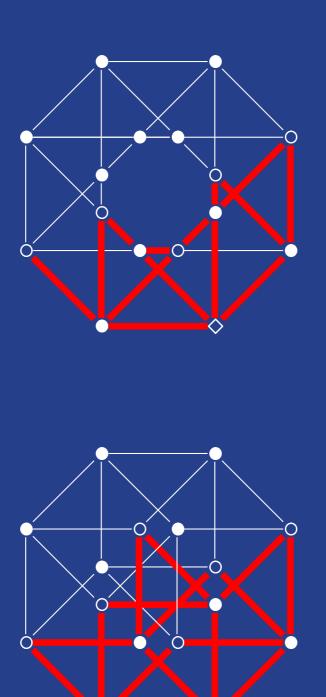
#### Are these graphs isomorphic?





#### The Subgraph Isomorphism Problem

#### **Example.** Graph isomorphism



 $G_2 = (V_2, E_2)$ , denoted by  $G_1 \cong S_2 \subseteq G_2$ , if there is an injection  $\varphi: V_1 \rightarrow C_2$  $(\varphi(v_i), \varphi(v_j)) \in E_2$  $V_2$  such that, for every pair of vertices  $v_i, v_j \in V_1$ , if  $(v_i, v_j) \in E_1$  then **Definition.** A graph  $G_1 = (V_1, E_1)$  is isomorphic to a subgraph of a graph

For input graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ , vertex  $u_i \in V_1$  $\deg(u_i) \leq \deg(v_j)$ , for all  $1 \leq i \leq n_1$  and  $1 \leq j \leq n_2$ cannot be mapped by a subgraph isomorphism to vertex  $v_j \in V_2$  unless

**SUBGRAPH ISOMORPHISM** 

INSTANCE Graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ 

QUESTION Does  $G_1$  contain a subgraph isomorphic to  $G_2$ ?

**Reference** Transformation from CLIQUE

Comment Contains CLIQUE, COMPLETE BIPARTITE SUBGRAPH, HAMILTONIAN CIRCUIT as special cases

a subgraph  $H_1$  of  $G_1$  and a subgraph  $H_2$  of  $G_2$  such that  $H_1 \cong H_2$ . The is not a proper subgraph of another common subgraph maximum common subgraph of two graphs is a common subgraph that **Definition.** A common subgraph of two graphs  $G_1$  and  $G_2$  consists of

 The maximum common subgraph is the largest possible common subgraph, while a common subgraph is maximal if it cannot be edges extended to another common subgraph by the addition of vertices or

MAXIMUM COMMON SUBGRAPH

**INSTANCE** Graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ , positive integer K

QUESTION Do there exist subsets  $E'_1 \subseteq E_1$  and  $E'_2 \subseteq E_2$  with  $|E'_1| = |E'_2| \ge K$  such that the two subgraphs  $G'_1 = (V_1, E'_1)$  and  $G'_2 = (V_2, E'_2)$ are isomorphic?

**Reference** Transformation from CLIQUE

MAXIMUM COMMON SUBGRAPH

INSTANCE Graphs  $G_1 = (V_1, E_1)$  and  $\overline{G_2} = (V_2, E_2)$ 

**SOLUTION** A common subgraph: graphs  $G'_1 \subseteq G_1$  and  $G'_2 \subseteq G_2$  such that  $G'_1$  and  $G'_2$  are isomorphic

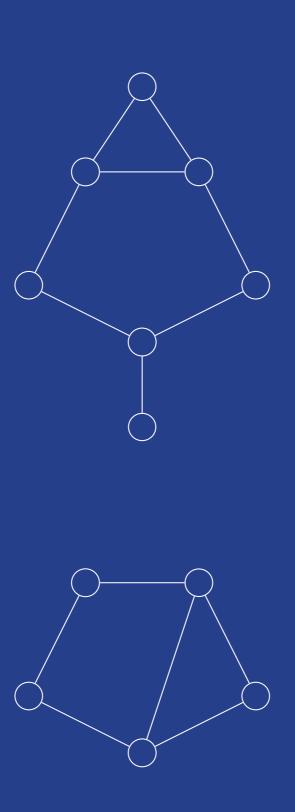
MEASURE Size of the common subgraph

one graph into the other least cost sequence of elementary graph edit operations that transform Definition. The edit distance between two graphs is the shortest or the

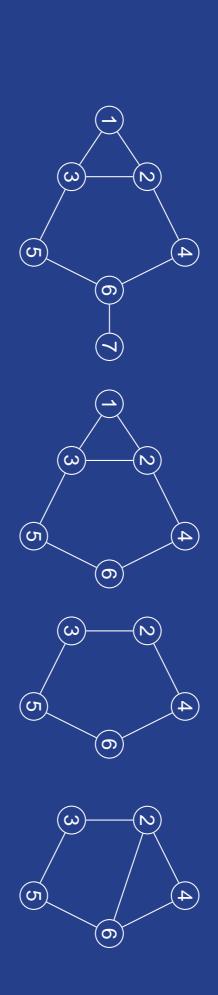
- Elementary edit operations include
- rotation
- substitution
- deletion
- insertion

of vertices and edges

What is the edit distance between the following graphs?

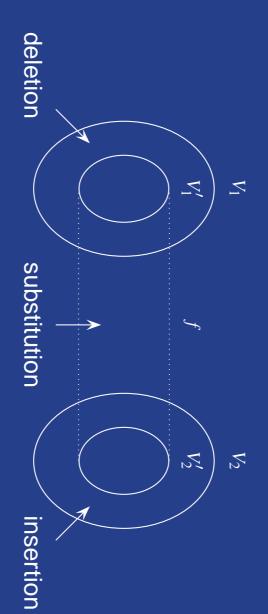


**Example.** Computing edit distance by deletion and insertion



- delete vertex 7
- delete vertex 1
- insert edge 2—6

 $G_2$  is a bijective function  $f: V_1' \rightarrow V_2'$ , where  $V_1' \subseteq V_1$  and  $V_2' \subseteq V_2$ **Definition.** An approximate graph matching from a graph  $G_1$  to a graph



real numbers **Definition.** A cost function is a tuple  $C = (c_{vd}, c_{vi}, c_{vs}, c_{es})$  of nonnegative

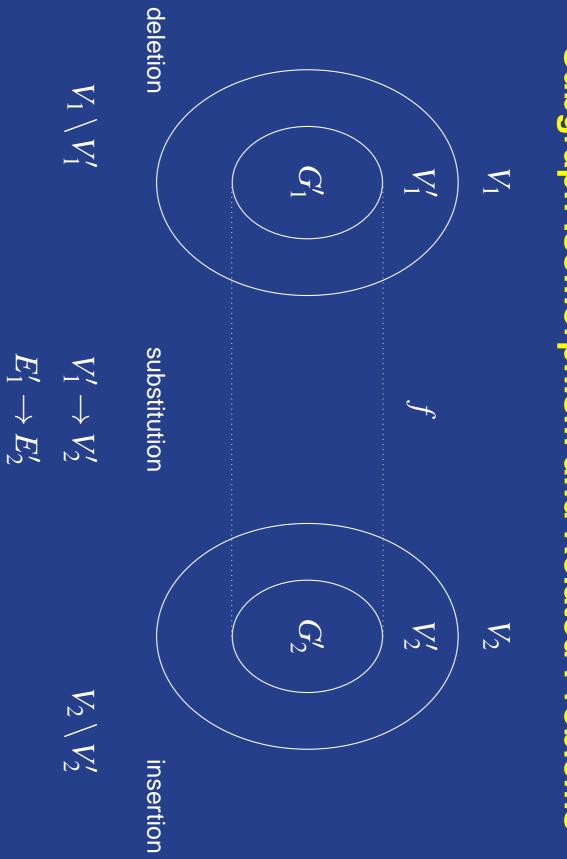
- $c_{vd}, c_{vi}, c_{vs}$  model the cost of vertex deletion, insertion, substitution
- c<sub>es</sub> models the cost of edge substitution

included in the costs of the corresponding vertex deletions and Edge deletion cost  $c_{ed}$  and edge insertion cost  $c_{ei}$  are assumed to be insertions

a graph  $G_1$  to a graph  $G_2$  is given by **Definition.** The cost of an approximate graph matching  $f: V'_1 \to V'_2$  from

$$\gamma_C(f) = \sum_{\nu \in V_1 \setminus V_1'} c_{\nu d}(
u) + \sum_{\nu \in V_2 \setminus V_2'} c_{\nu i}(
u) + \sum_{\nu \in V_1'} c_{\nu s}(
u) + \sum_{e \in E_1'} c_{es}(e)$$

where  $C = (c_{vd}, c_{vi}, c_{vs}, c_{es})$  is a cost function



The costs  $c_{\nu d}(\nu), c_{\nu i}(\nu), c_{\nu s}(\nu), c_{es}(\nu)$  correspond to

- deleting a vertex  $v \in V_1 \setminus V_1'$  from  $G_1$
- inserting a vertex  $u \in V_2 \setminus V_2'$  into  $G_2$
- substituting a vertex  $v \in V_1'$  by  $f(v) \in V_2'$
- substituting an arc  $e_1 = (u, v) \in E_1'$  by  $e_2 = (f(u), f(v)) \in E_2'$

where  $G_1' \sqsubseteq G_1$  and  $G_2' \sqsubseteq G_2$ 

of the) least cost approximate graph matching from  $G_1$  to  $G_2$ **Definition.** The edit distance between two graphs  $G_1$  and  $G_2$  is the (cost

 $\delta(G_1, G_2) = \min \left\{ \gamma_C(f) \mid f : G_1 \to G_2 \right\}$ 

Definition. A distance function  $\delta$  over graphs is a metric if it satisfies

- δ is positive definite;
- $\delta(G_1, G_2) \ge 0$
- $\delta(G_1,G_2)=0$  if and only if  $G_1\cong G_2$
- $\delta$  is symmetric:  $\delta(G_1, G_2) = \delta(G_2, G_1)$
- $\delta$  is triangular:  $\delta(G_1, G_2) \leq \delta(G_1, G_3) + \delta(G_3, G_2)$

A distance metric is useful for searching in a metric space

**GRAPH EDIT DISTANCE** 

**INSTANCE** Graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ , positive integer K

QUESTION Is  $\delta(G_1, G_2) \leq K$ ?

MINIMUM GRAPH TRANSFORMATION

**INSTANCE** Graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ 

**SOLUTION** A transformation that makes  $G_1$  isomorphic to  $G_2$ 

**MEASURE** Number of edges removed from  $E_1$  and added to  $E_2$ 

#### COMPLEXITY

#### AND

#### **APPROXIMATION**

RESULTS

- Subgraph isomorphism is NP-complete [GT48]
- Restriction to planar graphs remains NP-complete
- Fixed planar subgraph isomorphism is in P
- Maximum common subgraph is NP-complete [GT49]
- Approximation is APX-hard [GT46]
- Restriction to graphs of bounded degree is in APX
- Graph edit distance is NP-complete
- Approximation is APX-hard [GT49]

- Quest for practical algorithms
- Combinatorial methods
- Continuous optimization methods
- Relaxation labeling
- Mean-field annealing
- Probabilistic relaxation
- Discrete optimization methods
- Discrete relaxation
- Simulated annealing
- Genetic search

#### RECENT RESULTS

#### 0N

#### EXACT

#### SOLUTIONS

- Bunke, H. (1997). On a relation between graph edit distance and maximum common subgraph. Pattern Recogn. Lett. 18, 8, 689-694.
- The graph edit distance coincides with

$$\delta(G_1,G_2) = |V_1| + |V_2| - 2|\hat{V}_{12}|$$

if the cost function is such that

$$c_{\nu d} = c_{\nu i} = 1$$

- Bunke, H. and Schearer, K. (1998). A graph distance metric based 255-259. on the maximal common subgraph. Pattern Recogn. Lett. 19, 3–4,
- The graph distance measure given by

$$\delta(G_1,G_2) = 1 - rac{|\hat{V}_{12}|}{\max(|V_1|,|V_2|)}$$

is a metric

- Bunke, H. (1999). Error-correcting graph matching: On the influence of the underlying cost function. IEEE T. Pattern Anal. 21, 9, 917–922.
- The graph edit distance is a metric if the cost function is such that

$$c_{\nu d} + c_{\nu i} \leqslant c_{\nu s}$$

- Messmer, B. T. and Bunke, H. (1999). A decision tree approach to graph and subgraph isomorphism detection. Pattern Recogn. 32, 12, 1979–1998.
- Fixed graph and subgraph isomorphism is dealt with by storing all permutation matrices of the fixed graph in a decision tree
- Computational complexity is time  $\Theta(n_1^3)$  and space  $\Theta(3^{n_2})$  after preprocessing time  $\Theta(n_2^{n_2})$

- Eppstein, D. (1999). Subgraph isomorphism in planar graphs and 1-27. related problems. Journal of Graph Algorithms and Applications 3, 3,
- Fixed planar subgraph isomorphism is dealt with by partitioning the programming within each piece planar graph into pieces of small tree width, and applying dynamic
- Computational complexity is  $\Theta(n_2)$

- Cortadella, J. and Valiente, G. (2000). A relational view of subgraph Computer Science (Québec, Canada, 2000), pp. 45-54. isomorphism. In Proc. 5th Int. Seminar on Relational Methods in
- An explicit representation of the relation containing all and only all subgraph isomorphisms is built by intersection of binary relations
- Space efficiency is achieved by using symbolic techniques

- Larrosa, J. and Valiente, G. (2000). Graph pattern matching using 196 Workshop on Graph Transformation Systems (Berlin, 2000), pp. 189constraint satisfaction. In Proc. Joint APPLIGRAPH and GETGRATS
- Neighborhood constraints are exploited for domain filtering
- The new algorithm never visits more nodes than really full lookstructure constraints ahead and than forward checking using degree constraints and
- A benchmark for subgraph isomorphism is proposed

- Fernández, M.-L. and Valiente, G. (2001). A graph distance measure supergraph. Pattern Recogn. Lett. 22, 6–7, 753–758. combining maximum common subgraph and minimum common
- The graph distance measure given by

$$\delta(G_1,G_2) = |\check{G}_{12}| - |\hat{G}_{12}|$$

is a metric, where |G| = |V| + |E|

#### RECENT

#### RESULTS

#### 0 N

#### APPROXIMATE

#### SOLUTIONS

- Wilson, R. C., Evans, A. N., and Hancock, E. R. (1995). Relational 411-421. matching by discrete relaxation. Image and Vision Computing 13, 5,
- Approximate subgraph isomorphism is dealt with as a nonlinear optimization problem for a global measure of relational consistency
- A Bayesian measure of relational consistency is based on the sum of the matching probabilities over  $\Gamma(v)$  for all  $v \in V_1$

- Gold, S. and Rangarajan, A. (1996). algorithm for graph matching. IEEE T. Pattern Anal. 18, 4, 377-388. A graduated assignment
- Approximate subgraph isomorphism is dealt with as a nonlinear optimization problem
- The algorithm uses a "continuation method" to transform the discrete assignment problem into a continuous problem, in order to avoid poor local minima
- Computational complexity is  $O(m_1m_2)$

- Cross, A. D. J., Wilson, R. C., and Hancock, E. R. (1997). Inexact 970. graph matching using genetic search. Pattern Recogn. 30, 6, 953-
- Approximate subgraph isomorphism is dealt with as a nonlinear optimization problem for a global Bayesian measure of relational consistency
- The crossover process is realized at the level of subgraphs, rather than using string-based or random crossover
- Empirical results show
- Polynomial convergence time
- Convergence rate more rapid than simulated annealing

- Messmer, B. T. and Bunke, H. (1998). A new algorithm for error-493-504 tolerant subgraph isomorphism detection. IEEE T. Pattern Anal. 20, 5,
- Fixed approximate subgraph isomorphism is dealt with by storing a recursive decomposition of a set of fixed graphs
- Common subgraphs of different fixed graphs are represented only once
- The method is only sublinearly dependent on the number of fixed graphs

- El-Sonbaty, Y. and Ismail, M. A. (1998). A new algorithm for subgraph optimal isomorphism. Pattern Recogn. 31, 2, 205–218.
- Approximate subgraph isomorphism is dealt with as minimum weighted bipartite matching of decomposed subgraphs
- Graph G<sub>1</sub> is decomposed into n<sub>1</sub> subgraphs
- Graph  $G_2$  is decomposed into  $n_2$  subgraphs
- Computational complexity is average case  $\Theta(n_1^2 n_2^2)$ , worst case  $\Theta(n_1^2 n_2^2 \min(n_1, n_2))$
- Hidden weight of structure preservation

OPEN

PROBLEMS

#### Special cases

- Restriction to planar graphs remains NP-complete
- Planar clique is in P
- Planar Hamiltonian circuit is NP-complete
- 0 Restriction to graphs of bounded degree
- Bounded degree clique is in P
- Bounded degree graph isomorphism is in P
- Restriction to interval graphs
- Interval graph isomorphism is in P
- Restriction to chordal graphs
- Chordal clique is in P
- Chordal graph isomorphism is isomorphism-complete

- Approximation algorithms
- Most algorithms for approximate subgraph isomorphism and related problems are not approximation algorithms
- Approximate solutions are empirically shown to be close to the optimum, only for particular problem instances
- Theoretical analysis of existing algorithms for approximate subgraph isomorphism and related problems
- Polynomial-time approximation algorithms with bounded absolute or relative error (for special cases)
- Polynomial-time approximation algorithms with bounded inputdependent relative error