# Lecture-2 (Performance Evaluation) CS422-Spring 2020





#### Let's start with performance





## Performance

- <u>Latency</u> (execution/response time): time to finish one task. It is additive (Performance = 1/latency)
- *Throughput* (bandwidth): number of tasks/unit time. It is not additive
- Example: move people from A to B, **10** miles
  - Car: capacity = 5, speed = 60 miles/hour
  - Bus: capacity = 60, speed = 20 miles/hour
  - Latency: car = 10 min, bus = 30 min
  - Throughput: car = 15 PPH (w/ return trip), bus = 60 PPH

 Latency lags bandwidth, Bandwidth hurts latency, Read: https://cacm.acm.org/magazines/2004/10/6401-latency-lagsbandwith/fulltext

#### Latency vs Bandwidth

Latency vs Bandwidth, How they affect each other?

Latency helps bandwidth but not vice versa.



#### **Bandwidth usually hurts latency**

Queues -Bandwidth

#### Some More on Latency and Bandwidth

Bandwidth problems can be cured with money. Latency problems are harder because the speed of light is fixed – you can't bribe God

https://www.youtube.com/watch?list=PL2LuePcZTMh\_MzNHqZWNdvWdAnAThHCKK&v=lfqgpuH10uc&feature=emb\_logo https://www.youtube.com/watch?v=GNK-67JUH7M&list=PL2LuePcZTMh\_MzNHqZWNdvWdAnAThHCKK&index=2 https://www.youtube.com/watch?v=5CxpoGwCxKU&list=PL2LuePcZTMh\_MzNHqZWNdvWdAnAThHCKK&index=3

#### Energy and Power

Energy: Measure of using power for some time

Power: Instantaneous rate of energy transfer







# Amdahl's Law Source: The Guardian $ExTime_{new} = ExTime_{old} \times \left[ (1 - Fraction_{enhanced}) + \frac{Fraction_{enhanced}}{Speedup_{enhanced}} \right]$ $Speedup_{overall} = \frac{ExTime_{old}}{ExTime_{new}} = \frac{1}{(1 - Fraction_{enhanced})} + \frac{Fraction_{enhanced}}{Speedup_{enhanced}}$

#### Amdahl's Law

Which one will provide better overall speedup?

A. Small speedup on the large fraction of execution time.

- B. Large speedup on the small fraction of execution time.
- C. Does not matter.

Depends on the difference between small and large. Mostly it is A.

Amdahl's law for parallel processing

#### Evaluation

- To Compare Processor A with Processor X by running programs
- How many programs?



- The programs that you care.
- What if I want to build a new one (processor, caches, DRAM) ?



## World of Benchmarks

• SPEC CPU 2017 (<u>https://www.spec.org/cpu2017/</u>)

The **SPEC CPU® 2017** benchmark package contains SPEC's nextgeneration, industry-standardized, CPU intensive suites for measuring and comparing compute intensive performance, stressing a system's processor, memory subsystem and compiler.

SPECspeed: used for comparing time for a computer to complete single tasks

**SPECrate:** measure the throughput or work per unit of time.

## World of Benchmarks

CloudSuite (<u>https://www.cloudsuite.ch/</u>)

CloudSuite is a benchmark suite for cloud services. The benchmarks are based on real-world software stacks and represent real-world setups.

#### PARSEC (<u>https://parsec.cs.princeton.edu/</u>)

Benchmark suite composed of multithreaded programs. The suite focuses on emerging workloads and was designed to be representative of next-generation shared-memory programs for chip-multiprocessors.

#### World of Benchmarks

• MobileBench (<u>https://mobilebench.engineering.asu.edu/</u>) comprising a selection of representative smart phone applications.

## Pitfalls of Benchmarks

- Benchmark not representative
  - Your workload is I/O bound  $\rightarrow$  SPECCPU is useless
- Benchmark is too old
  - Benchmarks age poorly (SPEC CPU 2006 and then CPU 2017)
  - Benchmarketing pressure causes vendors to optimize compiler/hardware/software to benchmarks
  - $\rightarrow$  Need to be periodically refreshed

#### Non-Benchmarks

- Application kernels: A small code fragment or part of the program
- Synthetic benchmark : Not part of any real program!!
- Micro-benchmark

• OK! So, I will create a chip and then evaluate these benchmarks

#### World of Simulators

- Functional Simulator: Used to verify the correct execution of the program. Can not be used for performance evaluation.
- Performance simulators:
- (i) Trace-driven: ChampSim(<u>https://github.com/ChampSim/ChampSim</u>)
- (ii) Execution-driven: gem5, Multi2sim

Functional simulator is part of the performance simulators.

#### Evaluation Contd..

Pick a *relevant* benchmark suite

Measure IPC of each program

Summarize the performance using:

Arithmetic Mean (AM)

Geometric Mean (GM)

Harmonic Mean (HM)

Which one to choose?



#### Which machine performs better over IMTEL and why?

Contd.	ABM	AND	
App. one	2	3	
App. two	1.5	2	
App. three	e 1.3	0.3	
A.M.	1.60	1.76	
G.M.	1.57	1.21	
H.M.	1.54	0.72	

#### AM on Ratios Y Х Mine? 100 App. 1 1 1000 10 App. 2 Normalized to X X Y 100 App. 1 1 Y is 50 times faster than X App. 2 0.01 1 AM 50.005 1 ° Mine? Normalized to Y Х γ 0.01 1 App. 1 App. 2 100 1 AM 50.005 X is 50 times faster than Y 1

#### When to use what?

Edgar H. Sibley Panel Editor Using the arithmetic mean to summarize normalized benchmark results leads to mistaken conclusions that can be avoided by using the preferred method: the geometric mean.

Do not use A.M. on normalized numbers

#### HOW NOT TO LIE WITH STATISTICS: THE CORRECT WAY TO SUMMARIZE BENCHMARK RESULTS

Use G.M. for normalized numbers



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# What about Multi-core Systems?

Application *i* running on an N-core system

Throughput =  $\sum$  IPC (i)

Individual Slowdown (i) = CPI-together (i) / CPI-alone (i)

Weighted Speedup =  $\sum$  (IPC-together(i) / IPC-alone (i))

Harmonic Mean of Speedups =  $N/\Sigma$  (IPC-alone(i)/IPC-together (i))

Unfairness =

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Max-Slowdown/Min-Slowdown =
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max(Individual slowdowns)/min(individual slowdowns)

#### Todos

- Clock cycle, machine cycle, tick, .....
- AM and GM on ratios
- Reading assignment: Latency lags bandwidth and YouTube links

From Lecture-1 (Deadline: January 12):

Compare and contrast KD vs RM building [Put your perspective on Piazza]

Assignment 0: <u>https://docs.google.com/forms/d/e/1FAIpQLSeofKOH0a9oktzGdf7zW-</u> <u>9GpMmcVfnxqsbgLuspFGeJ-MfRIg/viewform</u>