

RYZEN PROCESSOR MICROARCHITECTURE

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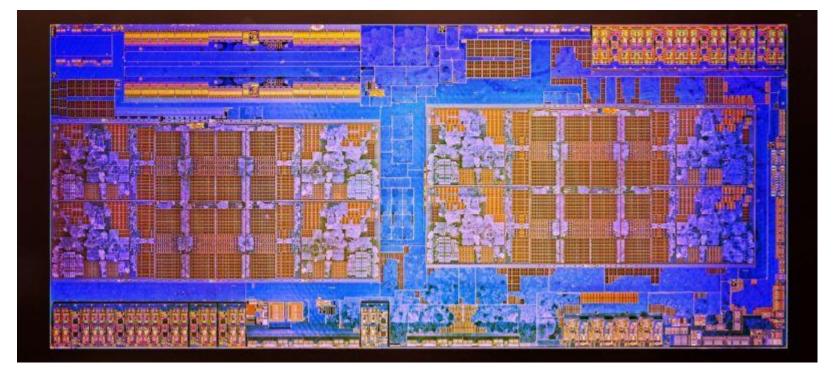
AGENDA

- Road to RYZEN AMD's new X86 CPU
- ▲ High level architecture of Zen
 - Core engine
 - Floating point unit
 - Cache hierarchy
 - Simultaneous Multithreading
- Zen based SoCs
- Design Challenges
- Contributions from India Design Center
- Next step

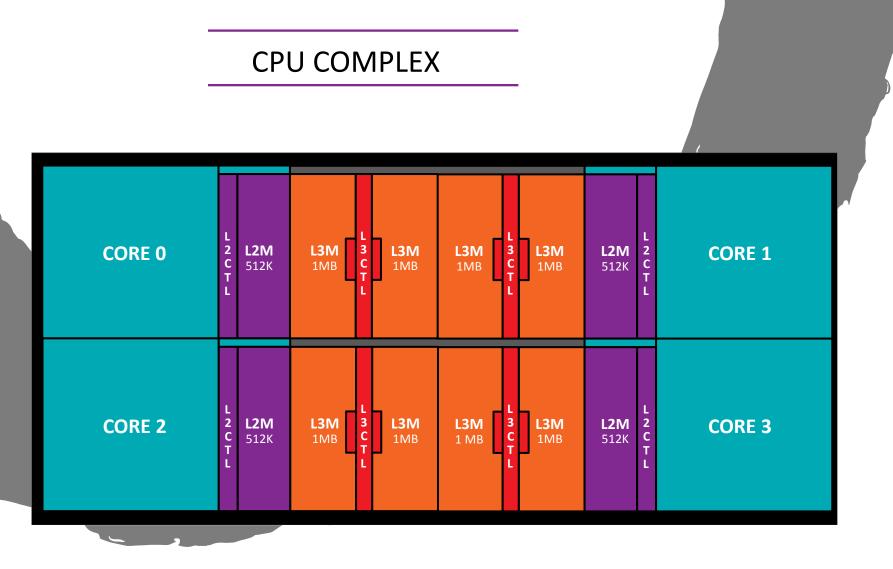




RYZEN CPU

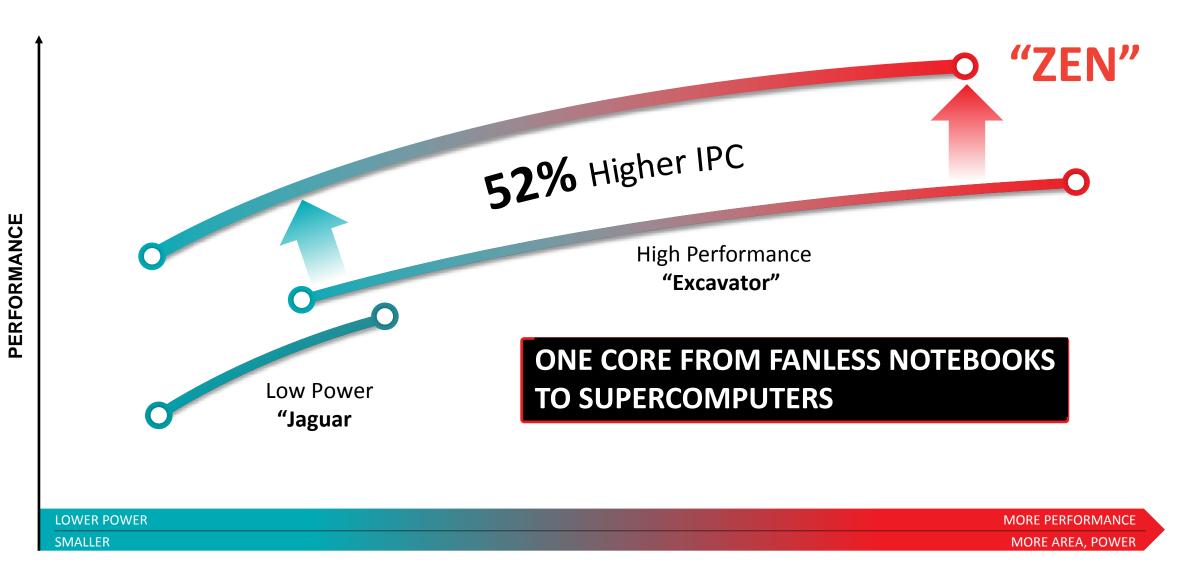


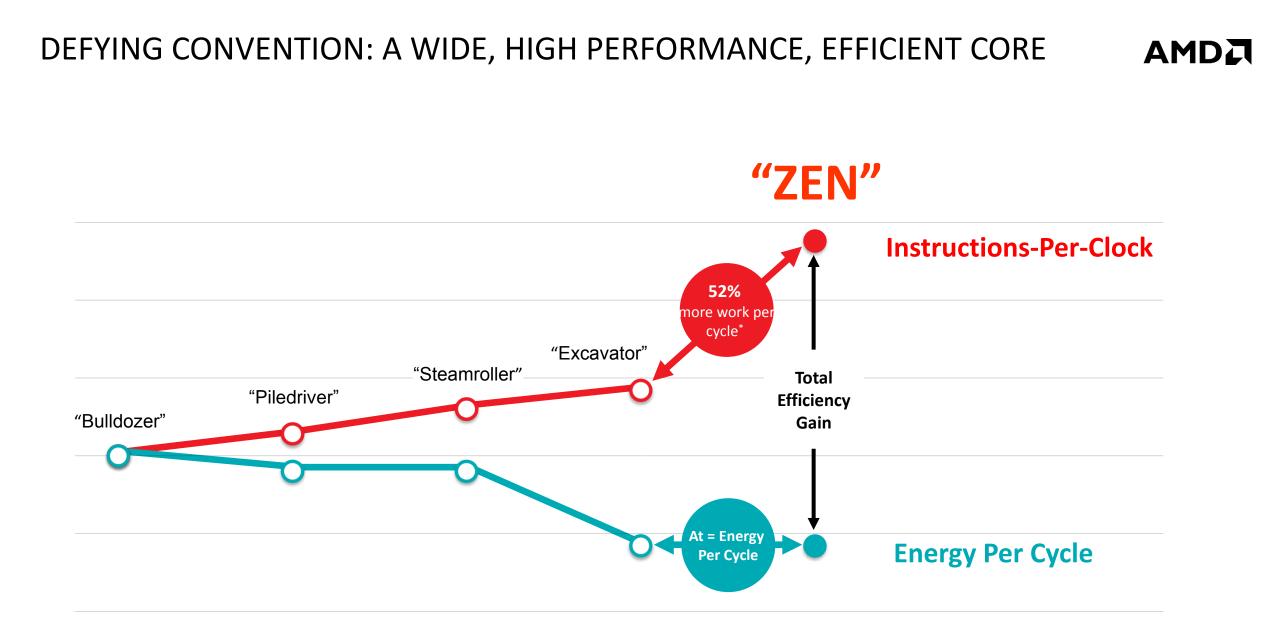
- ► High performance CPU
- ▶ 8 cores 16 threads
 - **Zen** core
- Over 4.8 billion transistors
- Over 2000 meters of wiring
- 14 nm process technology



- A CPU complex (CCX) is four cores connected to an L3 Cache.
- The L3 Cache is 16-way associative, 8MB, mostly exclusive of L2.
- The L3 Cache is made of 4 slices, by low-order address interleave.
- Every core can access every cache with same average latency
- Ryzen CPU in Thread Ripper has 2 CCXs
- Epyc Server has 8 CCXs
- ▲ Ryzen mobile has 1 CCX

AMD CPU DESIGN OPTIMIZATION POINTS





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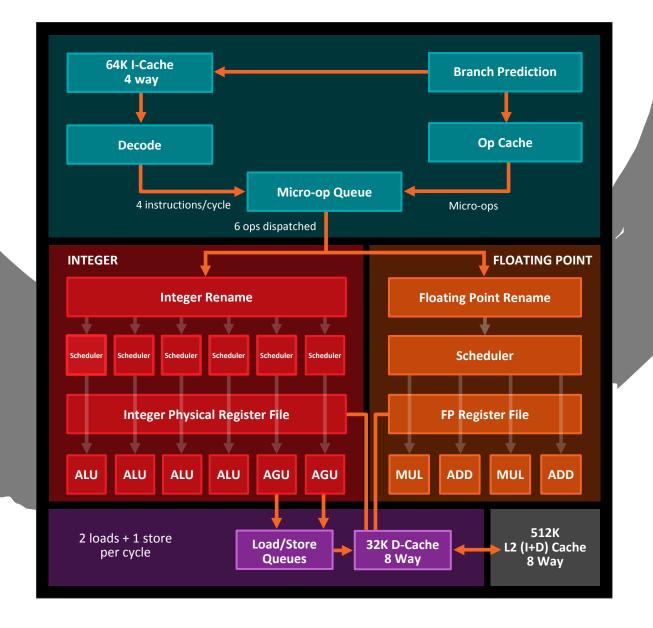
"ZEN"

Totally New High-performance Core Design PERFORMANCE AND POWER

DESIGNED FROM THE GROUND UP FOR OPTIMAL BALANCE OF

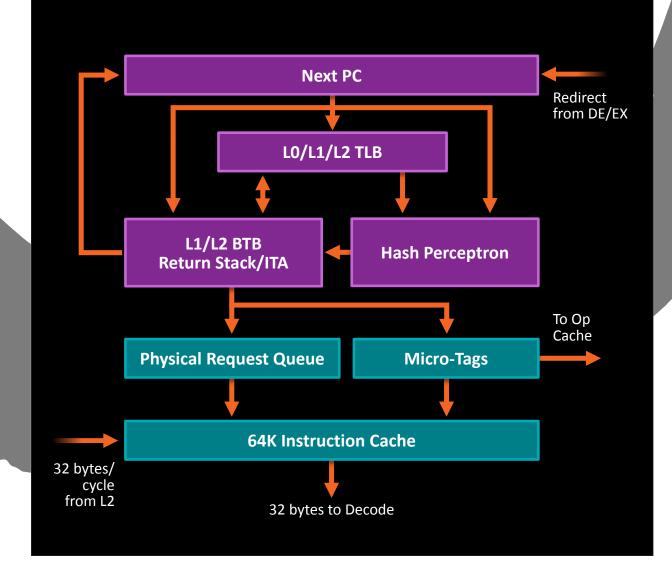
New High-Bandwidth, Low Latency Cache System

Simultaneous Multithreading (SMT) for High Throughput Energy-efficient FinFET Design Scales from Enterprise to Client Products



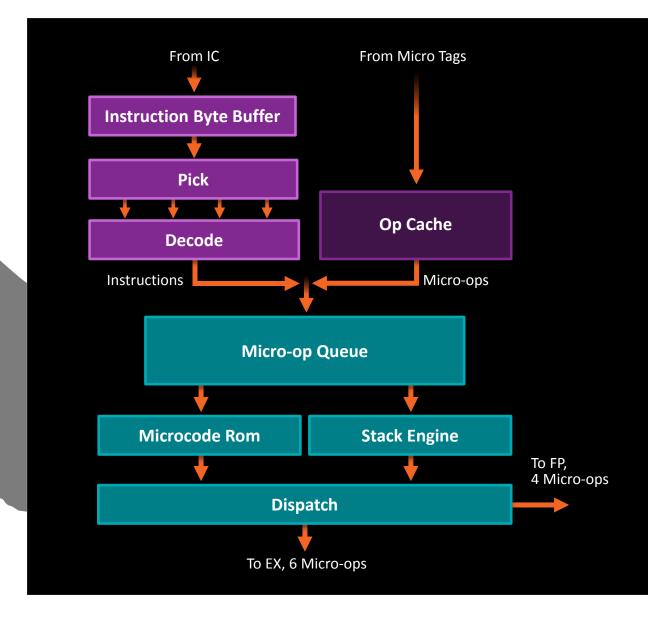
ZEN MICROARCHITECTURE

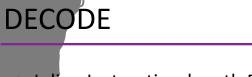
- Fetch Four x86 instructions
- Op Cache instructions
 - Integer units
 - Large rename space 168 Registers
 - 192 instructions in flight/8 wide retire
- ▲ Load/Store units
 - 72 Out-of-Order Loads supported
- ▲ Floating Point units
 - built as 4 pipes, 2 Fadd, 2 Fmul
- ▲ I-Cache 64K, 4-way
- ▲ D-Cache 32K, 8-way
- ▲ L2 Cache 512K, 8-way
- ▲ Large shared L3 cache 8M 16 way
- 2 threads per core



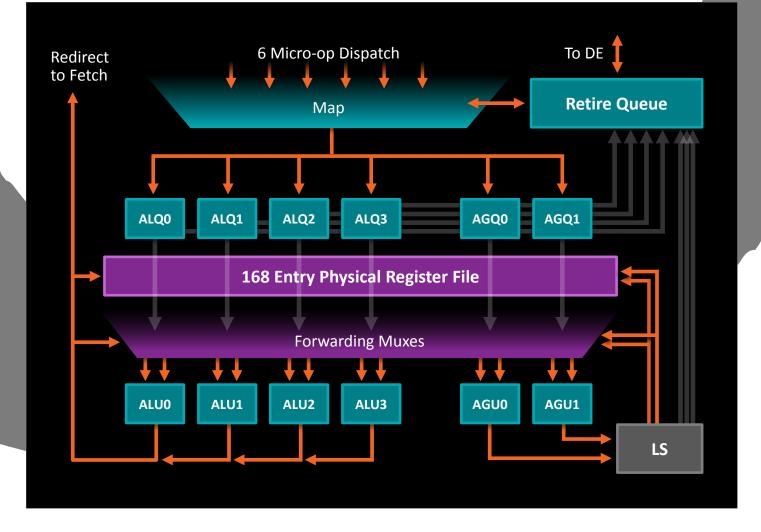
FETCH

- Decoupled Branch Prediction
- ▲ **T**LB in the BP pipe
 - 8 entry L0 TLB, all page sizes
 - 64 entry L1 TLB, all page sizes
 - 512 entry L2 TLB, no 1G pages
- ▲ 2 branches per BTB entry
- ▲ Large L1 / L2 BTB
- ▲ 32 entry return stack
- ▲ Indirect Target Array (ITA)
- ▲ 64K, 4-way Instruction cache
- ▲ Micro-tags for IC & Op cache
- ▲ 32 byte fetch



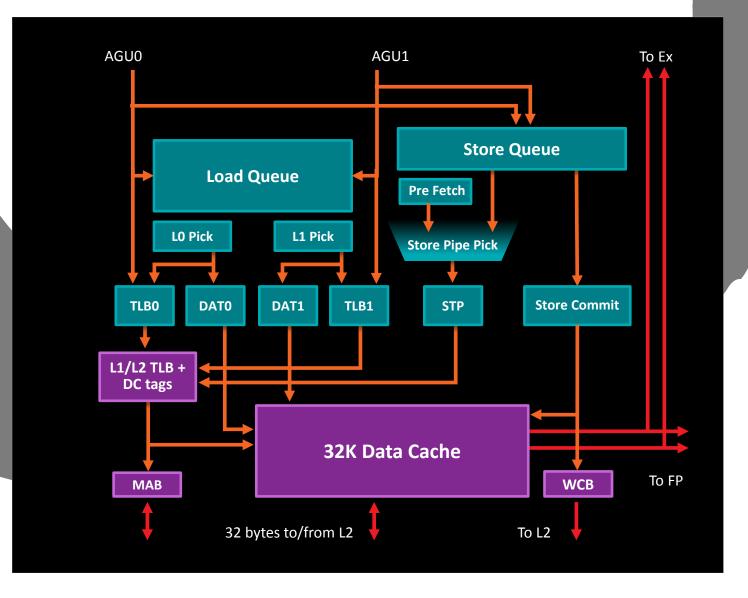


- Inline Instruction-length Decoder
- Decode 4 x86 instructions
 - Op cache
- Micro-op Queue
- ▲ Stack Engine
- Branch Fusion
- Memory File for Store to Load Forwarding



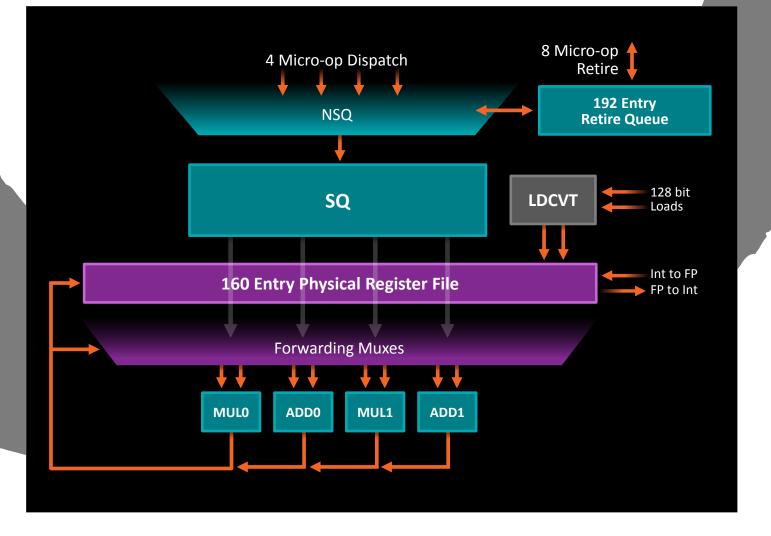
EXECÚTE

- ✓ 6x14 entry Scheduling Queues
- ▲ 168 entry Physical Register File
 - 6 issue per cycle
 - 4 ALU's, 2 AGU's
- ▲ 192 entry Retire Queue
- ▲ 2 Branches per cycle
- Move Elimination
- ▲ 8-Wide Retire



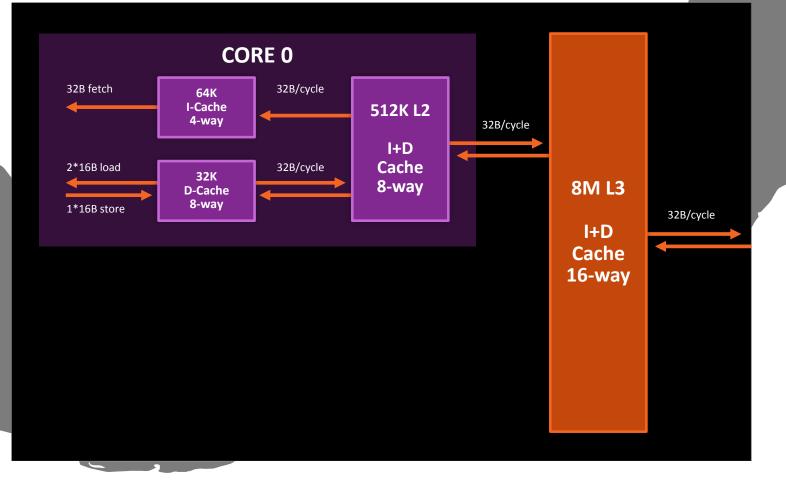
LOAD/STORE AND L2

- 72 Out of Order Loads
- ▲ 44 entry Store Queue
 - Split TLB/Data Pipe, store pipe
- 64 entry L1 TLB, all page sizes
- ▲ 1.5K entry L2 TLB, no 1G pages
- 32K, 8 way Data Cache
 Supports two 128-bit accesses
- Optimized L1 and L2 Prefetchers
- ▲ 512K, private (2 threads), inclusive



FLOATING POINT

- 2 Level Scheduling Queue
- 160 entry Physical Register File8 Wide Retire
- 1 pipe for 1x128b store
- Accelerated Recovery on Flushes
- SSE, AVX1, AVX2, AES and legacy mmx/x87 compliant
- ▲ 2 AES units



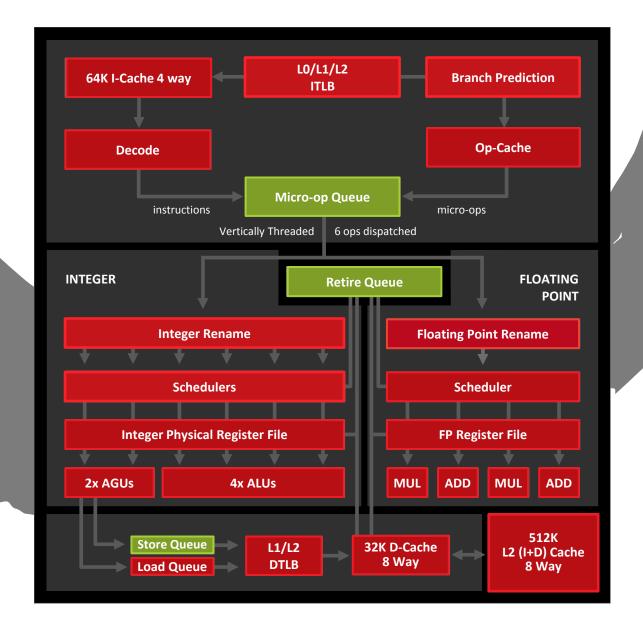
ZEN CACHE HIERARCHY

- Fast private 512K L2 cache
- ▲ Fast shared L3 cache
 - High bandwidth enables prefetch improvements
- ▲ L3 is filled from L2 victims
- ▲ Fast cache-to-cache transfers
- Large Queues for Handling L1 and L2 misses



CPU COMPLEX

- A CPU complex (CCX) is four cores connected to an L3 Cache.
- The L3 Cache is 16-way associative, 8MB, mostly exclusive of L2.
- The L3 Cache is made of 4 slices, by low-order address interleave.
- Every core can access every cache with same average latency
- Ryzen has 2 CCXs
- Naples Server has 4 CCXs



SMT OVERVIÉW

- ▲ All structures fully available in 1T mode
- Front End Queues are round robin with priority overrides
- Increased throughput from SMT



Competitively shared structures

Statically Partitioned

ZEN PERFORMANCE & POWER IMPROVEMENTS

BETTER CORE ENGINE

- SMT two threads per core
- Better branch prediction
- Large Op Cache
- Wider micro-op dispatch 6 vs. 4
- Larger Instruction Schedulers
 Integer: 84 vs. 48 | FP: 96 vs. 60
- Larger retire 8 ops vs. 4 ops
- 4 issue FPU
- Larger Retire Queue 192 vs. 128
- Larger Load Queue 72 vs. 44
- Larger Store Queue

BETTER CACHE

- Write back L1 cache
- Faster L2 cache
- Faster L3 cache
- Faster Load to FPU: 7 vs. 9 cycles
- Better L1 and L2 data prefetcher
- Close to 2x the L1 and L2 bandwidth
- Total L3 bandwidth up 5x

LOWER POWER

- Aggressive clock gating with multi-level regions
- Write back L1 cache
- Large Op Cache
- Stack Engine
- Move elimination
- Power focus from project inception

52% IPC PERFORMANCE UPLIFT

RYZEN BASED SOCS

EPYC SERVER

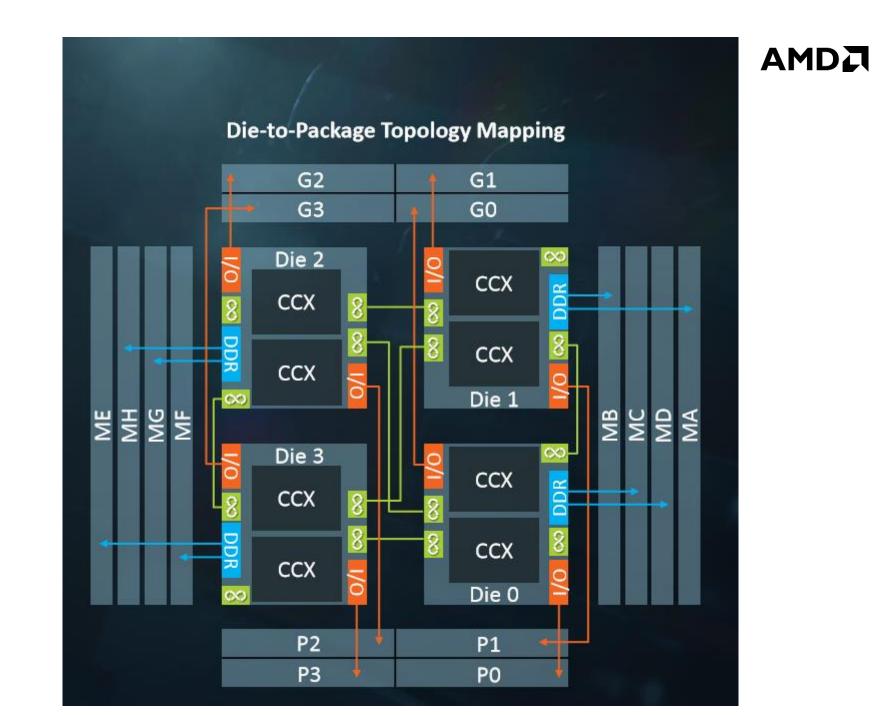
- 8 CCX, 4 Dies in each socket
- Connected by fast Infinity fabric
- Support for virtualization, encryption, ECC
- Targeted for Data center

▲ THREAD RIPPER DESKTOP

- High Frequency
- Gaming performance

RAVEN RIDGE LOW POWER APU (CPU + GPU)

- 4 M L3
- Power optimization



DESIGN CHALLENGES

- ▲ Long and expensive processor development project
 - More than 2 million engineering hours for Zen over four years
- Need to make key design decisions early in the project which has long lasting impact
 - Dependence on unknowns like process readiness in the fab, yield
 - Area, power
 - Availability of technology
- Conflicting needs of various SoCs
 - Server : Bigger LLC, MT performance more important than 1T performance
 - Desktop: High frequency, 1T is more important (aggressive prefetching)
 - Client: Area and power sensitive, smaller LLC
 - Semicustom: Specific ISA supports
 - Ultra low power
- Workloads to consider
- Competitiveness of the parts

CONTRIBUTIONS FROM INDIA DESIGN CENTER

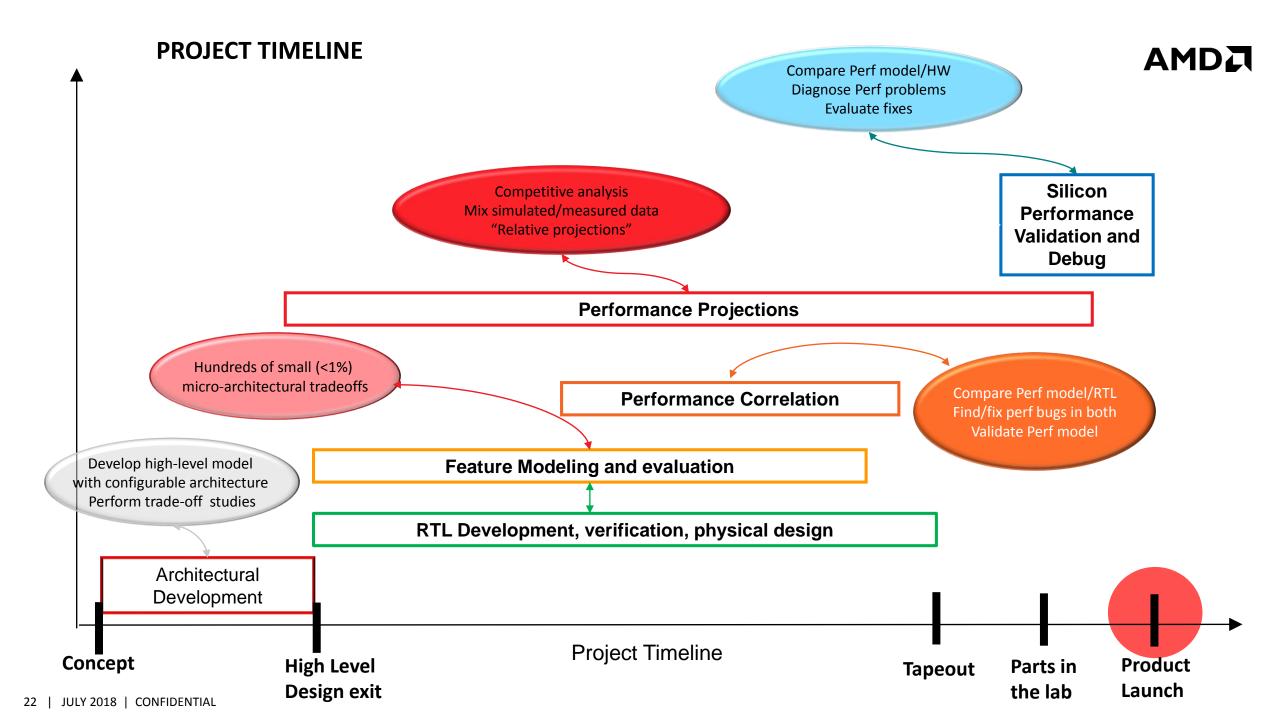


- Performance modeling, micro-architectural feature exploration, performance verification
 Development of cache hierarchy performance model (L2, L3, prefetcher) and architecture
- ▲ Traffic optimization between cores and memory
- ▲ All (Server, Client and Cores) workloads analysis, characterization and collection
- ▲ Core, Data Fabric and SoC level verification
- ▲ L3 RTL development for Zen based client part
- Zen based client SoC architecture development
- Compiler optimizations for Zen based server (Epyc)

WORK IN PROGRESS

- Zen2 hardware is in the lab
 - 7 nm core with significant performance uplift over Zen
- Microarchitecture development, performance modeling and correlation for core following Zen2
- Exploring micro-architectural ideas for future processors
- Improving performance analysis and modeling methodologies
 - Performance bottleneck analysis
 - Very high thread count simulation, impact of sharing and locks
 - Warming-up large caches
 - Execution driven vs trace driven simulation
 - SoC level modeling CPU/GPU/accelerators interaction
- Workloads analysis
 - Cloud, virtualized workloads, machine learning, HPC workloads

Lots of challenging problem to work on!!



PERFORMANCE ANALYSIS AND MODELING

- **Develop and explore micro-architectural ideas for performance improvement**
- Variety of tools and technologies deployed for performance analysis & architectural innovation
 Cycle accurate simulator
- Constant state of evolution, due to changing workloads, and architecture
 - More IPC, MP, SMT, large caches
 - SW and HW more tightly linked than ever
- ▲ Correct performance projection is critical for the success of the project
- Workloads analysis, characterization and tracing
- Server workloads SPEC CPU 2006, SPEC CPU2017, Legacy Enterprise workloads, Spec JBB, Cloud, NoSql database
- Client workloads Games, Cinebench, Productivity suite, various use case scenarios

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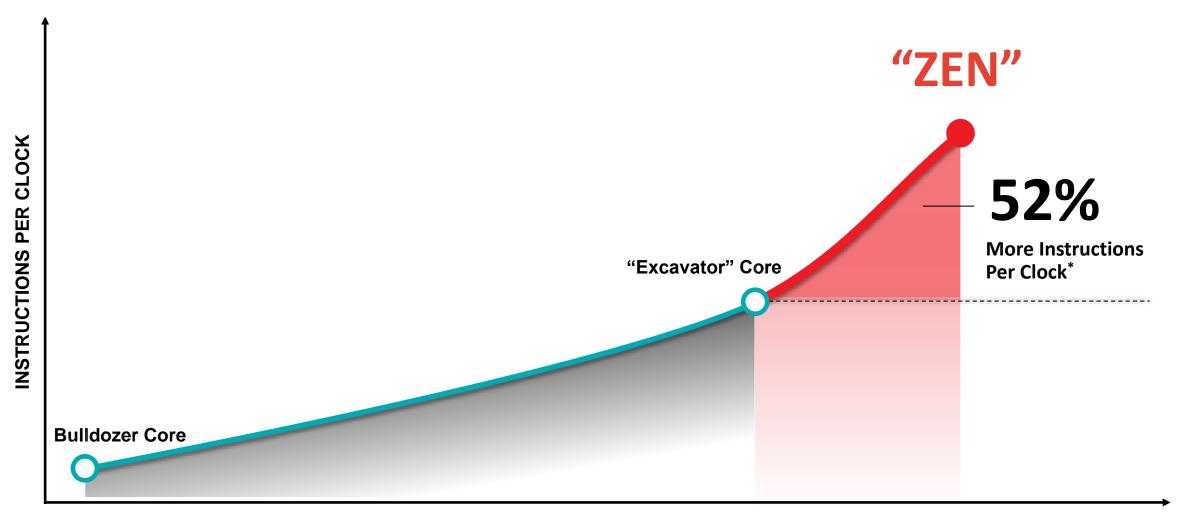
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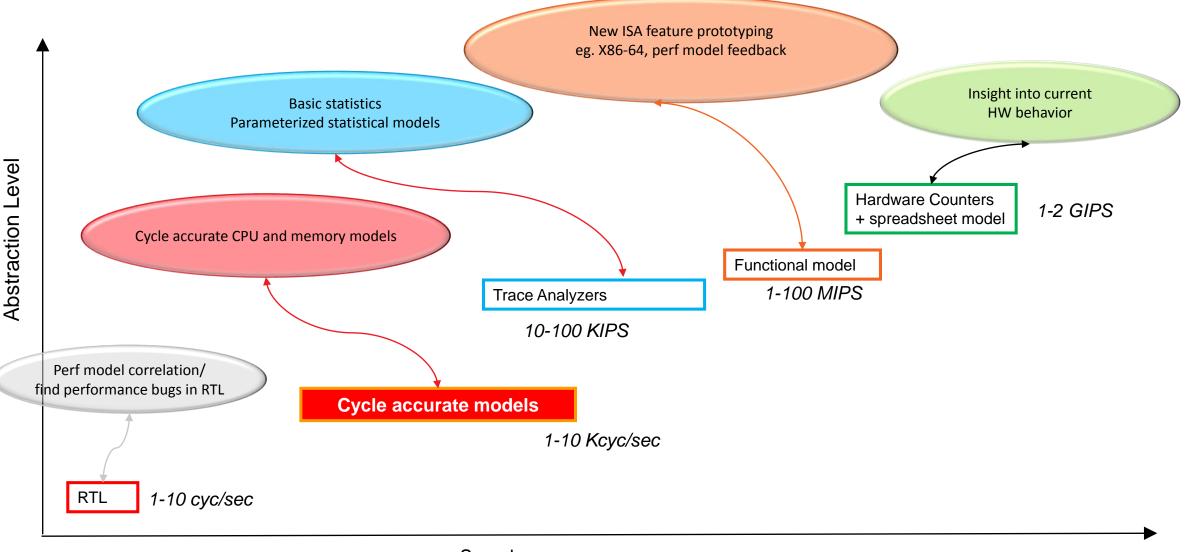
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AMD X86 CORES: DRIVING COMPETITIVE PERFORMANCE



*Based on internal AMD estimates for "Zen" x86 CPU core compared to "Excavator" x86 CPU core.

PERFORMANCE ANALYSIS TOOLS



Speed