CS698F Advanced Data Management

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- Consider graph as a list of edges $\langle \text{S,P,O} \rangle$
- Various options of distribution
 - S as the key, P as the key, O as the key
 - $\langle \text{S,O}\rangle$ as the key, $\langle \text{S,P}\rangle$ as the key... so on
 - What does this remind you of?
- Distributed join method changes slightly for each type of distribution.

- S as the key
 - All patterns that have join on S alone can be done without any "shipping" of tuples/edges.



Explanation shown on the board



- S as the key
 - Any combination of S-O join or a join on P requires "shipping".



Explanation shown on the board



- Then how to decide distribution "key"?
 - S, O, P?
 - Or \langle S,O \rangle , \langle O,P \rangle , \langle S,P \rangle ?
- What do you join on always?
 - Vertices of the graph
- Do all joins happen only on S, or only on O, or S-O?
 - Most joins happen on S, O, and S-O too.

- A join on S means you need "outgoing" edges of vertices.
- A join on O means you need "incoming" edges of vertices.
- A join on S-O means you need "incoming" AND "outgoing" edges of vertices.
- Then how to distribute?
 - For every edge of the graph, distribute $\langle hash(S), \, out\text{-edge-list}\rangle, \, and \, \langle hash(O), \, inc\text{-edge-list}\rangle$
- This ensures that first level join in any pattern can always be performed locally.

Back to our example





Join plans

- Can we perform this join without "shipping" of tuples/edges?
 - If we have used only (hash(S), out-edges) as the distribution policy?
 - If we have used only (hash(O), inc-edges) as the distribution policy?
 - If we have used **both** of the above?





- AFTER the first join, we have generated following partial results.
 - \langle :Jerry, :hasFriend, var-bind(?friend), :actedIn var-bind(?sitcom) \rangle
 - We need to join this partial result with the bindings for (?sitcom :location :NYC).
 - Which join is this?
 - S-S, O-O, or S-O?
 - Can we perform this join without "shipping"?
 - If yes, how?
 - If not, why not?

- At the first join
 - What was the join variable ?friend
 - So all the graph edges will be distributed according to the variable bindings (values) of *?friend*
 - All the same bindings/values of *?friend* will be on the same compute-node.
 - That will not help while joining on ?sitcom
 - Why not?

- The values/bindings of ?sitcom in (:Jerry, :hasFriend, varbind(?friend), :actedIn var-bind(?sitcom) are located on the compute-node according to hash(var-bind(?friend))
- For a join on ?sitcom, we need to have graph edges located on compute node according to hash(var-bind(? sitcom))
 - This needs to be taken care of by shipping!
 - How do you decide what to ship, and where to ship?

Options to ship

- Two options:
 - Ship (:Jerry, :hasFriend, var-bind(?friend), :actedIn var-bind(?sitcom))
 by hash(var-bind(?sitcom)) from each result.
 - Ship each graph edge matching (?sitcom :location :NYC) to ALL the compute-nodes?
 - Why? Why not hash(var-bind(?sitcom)) from each edge matching (?sitcom :location :NYC)?
- Which of these would be cheaper?
 - Any intuition?







- When can we do these joins simultaneously?
 - If we distribute using only S as the key?
 - If we distribute using only O as the key?
 - If we distribute using both?
- How do we combine the independent join results?

- After the first two simultaneous joins partial results of
 - <:Jerry, :hasFriend, var-bind(?friend), :actedIn, var-bind(?sitcom)>
 - Each result located as per hash(var-bind(?friend))
 - <var-bind(?friend), :actedIn, var-bind(?sitcom), :location, :NYC>
 - Each result located as per hash(var-bind(?sitcom))
- Now what to ship where?
 - Hint: what would you join these patial results on?

Two options to ship

• Two options:

- Ship <: Jerry, :hasFriend, var-bind(?friend), :actedIn var-bind(?sitcom)> by hash(var-bind(?sitcom)) from each result, leaving other results at their location.
- Ship <var-bind(?friend), :actedIn, var-bind(?sitcom), :location, :NYC> by hash(var-bind(?friend)) from each result, leaving other results at their location.
- Which of these would be cheaper?
 - How do you decide?
 - Hint: recall our join result size estimation!



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