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M. Atre

Announcements

Recap

Reachabilit Queries

Query Processing

Advanced Data Management

Medha Atre

Office: KD-219 atrem@cse.iitk.ac.in

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Reachability Queries

Query Processing Assignment-3 will be posted in the next couple of days. It will be a programming assignment requiring use of Hadoop.

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How it is all tied together

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Recap

Reachability Queries

Query Processing

Data Management

- Single compute node.
- Cluster.
- For single compute node query optimization, data compression, indexing.
- For a cluster all the aspects above plus cost of data shipping, data distribution strategies etc.
- Special topics in graph data reachability queries, regular path queries, top-k queries, keyword search.
- Theoretical aspects of query optimization and data distribution.

Reachability defined

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Recap

Reachability Queries

Query Processing Given a graph G(V, E) with V as the set of nodes and E as the set of edges, a reachability query asks – does there exists any path between nodes x and y.

In case of directed graph, the path is directed and other for undirected graphs the path is undirected.

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The graphs may have cycles too, giving rise to "strongly connected components".

Query Processing Challenges

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Recap

Reachability Queries

Query Processing

- Identification of strongly connected components and merging them.
- Efficient traversal over large graphs, e.g., number of vertices and edges being several millions.
- High computational complexity O(V³) for standard algorithms like all-pair shortest path, which also gives us *transitive clousure* for answering reachability.
- High storage space O(V²) if the entire transitive clousure is meant to be stored.

Some specific approaches

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Some trivial brute-force approaches

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Query Processing

- Repeated adjancency matrix multiplication.
 R = A + A² + A³... + Aⁿ Computational complexity is O(n⁴), note that at a time we do only one matrix multiplication and we do this n times.
- Do a DFS walk over the DAG of the graph (after coalescing SCCs) and construct compressed bit-vectors [SIGMOD 2010].
- Do repeated walk over the DAG of the graph, and construct hyper-dimensional label covers [Grail, VLDB Conf. 2010].